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ROCHESTER INSTITUTE OF TECHNOLOGY
SCHOOL OF COMPUTER SCIENCE AND TECHNOLOGY

IMPLEMENTATION OF AN
LALR(1) PARSER GENERATOR

A Thesis Submitted in Partial Fulfillment of
Master of Science in Computer Science Degree Program

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DATE: March 1980

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LALR(1) Parser Generator

by

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Implementation of an LALR(1) Parser Generator

Introduction

The purpose of this thesis was to implement an LALR(1) parser generator using the algorithm and methods presented in Aho and Ullman (1977). Generally speaking this meant to input the definition of a LR(1) grammar and output tables that could be used by a parser to decide whether or not arbitrary sentences from the grammar are syntactically correct.

This paper is organized into the following sections:

- I. Background. A brief summary of languages, grammar, parsing, and parsing techniques is presented.
- II. LALR(1) Parser Generator. A review of the method suggested by Aho and Ullman (1977).
- III. LALR(1) Parser Generator and Parser Overview. The Overview shows the interaction between the parser generator and the parser.
- IV. LALR(1) Parser Generator. A detailed description of the procedures used to implement the parser generator.
- V. Shift-Reduce Parser. A detailed description of the parser implemented to test the parser generator tables.
- VI. Sample runs of selected grammars through the parser generator and parser are presented and explained.
- VII. Conclusion. Discussion of implementation problems encountered, possible extensions and concluding remarks.

I. Background

Languages

According to Cleaveland and Uzgalis (1977) a "language" is a defined set of strings. "Strings" are a sequence of symbols using the alphabet of the language.

Grammars

Grammars define languages. "Phrase structured" grammars are defined by Barrett and Couch (1979) as a four-tuple (Σ, N, P, S) where:

- Σ Is the set of all terminal symbols within the grammar.
- N Is the set of all nonterminal symbols within the grammar where Σ and N are disjoint.
- P Is the set of all productions of the form $y \rightarrow x$ where y and x are in $(N\cup\Sigma)^*$ and y contains at least one element in N, and
- S Is a designated start symbol in N.

Chomsky (1965) defined four classes of phrase structured grammars:

- Type 3 which defines Regular languages,
- Type 2 which defines Context Free languages,
- Type 1 which defines Context Sensitive languages, and
- Type 0 which defines Recursively Enumerable languages.

Each type of grammar defines a set of languages which is a proper subset of the set defined by any lower numbered grammar type. For example, Type 0 grammars can define all of the languages definable by Type 1, Type 2 or Type 3.

The difference between the four classes of grammar are explained below using the four-tuple mentioned above.

Type 3 grammars have productions of the form:

- $A \rightarrow xB$ or
- $A \rightarrow x$ where A and B are in N and x is in Σ^* .

Type 2 grammars have productions of the form:

- $x \rightarrow y$ where x is a member of N and y is any string in $(N\cup\Sigma)^*$

Type 1 grammars have productions of the form:

$x \rightarrow y$ where x and y are members of $(N \cup \Sigma)^*$, x contains at least one member of N , and $|x| \leq |y|$.

Type 0 grammars have productions of the form:

$x \rightarrow y$ where x is a member of $(N \cup \Sigma)^+$ and y is a member of $(N \cup \Sigma)^*$.

It can be seen by examining the types of grammars that as we proceed from Type 3 to Type 0 the restrictions on the productions decrease thereby allowing greater freedom in the resultant languages and also increasing the complexity of the parsing problem.

Context free languages are the most practical because they strike a happy medium between generality and reasonable cost. In other words these languages can be general enough to look like plain english (i.e. PL/1 and ALGOL) yet restrictive enough so that practical compilers can be written for them. PL/1 and ALGOL are examples of context free languages although there are context sensitive elements in both these languages that are handled by other means than through the language definition.

One very important part of the compiler is the parser. It is the parser's job to check incoming sentences from the language for correct syntax. Does the string fulfill the requirements of the grammar?

As an example of the parsing problem consider the following grammar:

G: ($\{+, *, (,), a\}$, $\{E, T, F\}$, P, E) where P is the set of productions:

1. $E \rightarrow E+T$
2. $E \rightarrow T$
3. $T \rightarrow T * F$
4. $T \rightarrow F$
5. $F \rightarrow (E)$
6. $F \rightarrow a$

$a + a$ is a valid sentence in this language because it can be derived from the following productions:

$E \rightarrow E+T$	rule 1
$\rightarrow E+F$	rule 4
$\rightarrow E+a$	rule 6
$\rightarrow T+a$	rule 2
$\rightarrow F+a$	rule 4
$\rightarrow a+a$	rule 6

However, the sentence $a+$ is invalid because we can see by visual inspection that the terminal symbol $+$ cannot be the last terminal in the sentence.

The parsing problem is to take the input sentences and determine if, by applying the correct sequence of productions, the sentence is valid for this grammar.

Three possible approaches to the parsing problem are:

1. Backtracking whereby we start with the first production and then resolve a nonterminal on the right part of the production arbitrarily by another production and keep on picking other productions until it is determined that an error has been made. At that point the parser must "backtrack" through the parse tree to try another path until either all paths have been exhausted or the sentence is accepted. Obviously this approach would be very time consuming.
2. Deterministic Top-Down Parser where the parser starts with the first production and by looking at the next input symbol can positively determine the next production to apply. This is also called a LL(1) parser where LL means Left to right scan of input sentence with Leftmost nonterminal resolved first.
3. Deterministic Bottom-up Parser where the parser starts with the input string and works "backward" through the productions to obtain the correct parse. One of the broad classes of bottom-up parsers is the LR(k) parser where LR means Left to Right scan of input with rightmost nonterminal resolved first. The (k) means that the parser will look at the next k symbols to decide which production to use. The LR(1) parser is a special case where the parser looks ahead one symbol to determine which production to use.

Another name for this type of parsing is "shift-reduce" parsing and in this name lies the key to this technique. As the input string of terminals is examined the parser looks for any substring that matches the right part of any production (also called a "handle") in the grammar. Until a match is found the input terminals are shifted onto a stack. When the match is found then the stacked symbols are replaced with the left side of the production. The replacement process is called a reduction, hence the name shift-reduce parsing. The shifting and reducing continues until either: (a) the final reduction gives the start symbol in which case the input string is valid for this grammar or (b) no more reductions can be made and the start symbol has not been reached in which case the string is invalid for this grammar.

Later on in this paper the complete algorithm for shift-reduce parsing is presented.

As summarized by Barrett and Couch (1979) the technique that is most general yet is able to detect a syntax error at the very earliest point is the last technique described above; the LR(1) parser.

This technique is very similar to the one outlined in Aho and Ullman (1977) except Aho and Ullman go one step beyond the LR(1) parser to the LALR(1) parser where LA stands for "lookahead". As it turns out LR(1) parsers are very expensive when it comes to space. The LALR(1) parser is a significant improvement over the LR(1) parser in this respect.

The original work done on LR(k) parsers is Knuth (1965). The k indicates the general case where k symbols are looked at before deciding which production to apply next.

II. LALR(1) Parser Generator

One method described in Aho and Ullman (1977) of obtaining the LALR(1) parsing tables is to construct the LR(1) parsing tables and then by examining the generated items reduce the tables to LALR(1) tables.

The key element of the LR(1) parsing table construction is to generate the LR(1) sets-of-items. According to Barrett and Couch (1979) an item is a production carrying a position marker and a lookahead symbol. Intuitively everything to the left of the marker has been recognized and the lookahead symbol represents a possible terminal grammar symbol that could appear next on the input stream for this production. The LR(1) sets-of-items are a series of states that define the possible parser for the given grammar. As an input string is analyzed the parser moves through the various states to determine if the sentence is correct. As an example of an item consider the following:

$A \rightarrow B.CD, +$

This is an item that references production:

$A \rightarrow BCD$

and that the first symbol in this production, B has been recognized and that the lookahead symbol is a '+'. The lookahead can be any terminal symbol in the grammar or a '\$' signifying that the end of the string has been reached.

The following procedures, taken from Aho and Ullman (1977), generate the LR(1) sets-of-items.

Procedure Main;

```
begin
  C := { CLOSURE( {S' -> 'S,$' } ) };
  repeat
    for each set of items I in C and each grammar
      symbol X such that GOTO(I,X) is not empty
      and not already in C do
        add GOTO(I,X) to C
    until no more sets of items can be added to C
  end
  procedure CLOSURE(I);
  begin
    repeat
      for each item  $[A \rightarrow \alpha \cdot B\beta, a]$  in I, each
        production  $B \rightarrow y$ , and each terminal b in FIRST( $\beta a$ )
        such that  $[B \rightarrow \cdot y, b]$  is not in I do
          add  $[B \rightarrow \cdot y, b]$  to I;
    until no more items can be added to I;
    return
  end;
  procedure GOTO(I,X);
  begin
    let J be the set of items  $[A \rightarrow \alpha X\beta, a]$ , such that
       $[A \rightarrow \alpha \cdot X\beta, a]$  is in I;
    return CLOSURE(J)
  end;
```

The function FIRST (B) returns the set of terminals that begin strings derived from B.

Procedure CLOSURE calculates the transitive completion mentioned in Knuth (1965) and means intuitively that if item $A \rightarrow \cdot B\beta$ is in CLOSURE (I) then we would expect, at some point in the parsing process, to see a string derivable from $B\beta$. Therefore CLOSURE adds items to this set of the form $B \rightarrow \cdot \gamma$.

The GOTO procedure calculates, for a given state and grammar symbol, the next state the parser should go to.

The following algorithm, extracted from Aho and Ullman (1977) (Algorithm 6.3 in their book) gives the method for constructing the LR(1) parsing tables GOTO and ACTION.

Construction of LR(1) parsing tables.

Input. A grammar G augmented by production $S' \rightarrow S$.

Output. If possible, the canonical LR parsing action function ACTION and goto function GOTO.

Method.

1. Construct $C = \{I_0, I_1, \dots, I_n\}$, the collection of sets of LR(1) items for G.
2. State i of the parser is constructed from I_i . The parsing actions for state i are determined as follows:
 - a) If $[A \rightarrow \alpha \cdot a\beta, b]$ is in I_i and $\text{GOTO}(I_i, a) = I_j$, then set $\text{ACTION}[i, a]$ to "shift j ."
 - b) If $[A \rightarrow \alpha \cdot, a]$ is in I_i , then set $\text{ACTION}[i, a]$ to "reduce $A \rightarrow \alpha$."
 - c) If $[S' \rightarrow \cdot S, \$]$ is in I_i , then set $\text{ACTION}[i, \$]$ to "accept."

If a conflict results from the above rules, the grammar is said not to be LR(1) and the algorithm is said to fail.

3. The goto transitions for state i are determined as follows:
If $\text{GOTO}(I_i, A) = I_j$, then $\text{GOTO}[i, A] = j$
4. All entries not defined by rules (2) through (3) are made "error."
5. The initial state of the parser is the one constructed from the set containing item $[S' \rightarrow \cdot S, \$]$.

The action function ACTION and the goto function GOTO are two dimensional arrays where the first dimensions represents the states and the second dimension represents the grammar symbols. The elements of the goto function are states whereas the elements of the action function indicate the action that the parser is to take: shift, reduce by a production number, accept sentence, or error meaning the sentence cannot be parsed.

To obtain the LALR(1) parsing tables Aho and Ullman (1977) present the following algorithm (Algorithm 6.4 in their text).

Input. A grammar G augmented by production $S' \rightarrow S$.

Output. The LALR parsing tables ACTION and GOTO.

Method.

1. Construct $C = \{I_0, I_1, \dots, I_n\}$, the collection of sets of LR(1) items.
2. For each core present among the sets of LR(1) items, find all sets having that core, and replace these sets by their union.
3. Let $C' = \{J_0, J_1, \dots, J_m\}$ be the resulting sets of LR(1) items. The parsing actions^m for state i are constructed from J_i in the same manner as in Algorithm 6.3. If there is a parsingⁱ action conflict, the algorithm fails to produce a parser, and the grammar is said not to be LALR(1).
4. The GOTO table is constructed as follows: If J is the union of one or more sets of LR(1) items, i.e., $J = I_1 \cup I_2 \cup \dots \cup I_m$, then the cores of $\text{GOTO}(I_1, X)$, $\text{GOTO}(I_2, X)$, ..., $\text{GOTO}(I_m, X)$ are the same, since I_1, I_2, \dots, I_m all have the same core. Let K be the union of all sets² of items having the same core as $\text{GOTO}(I_1, X)$. Then $\text{GOTO}(J, X) = K$.

III. LALR(1) Parser Generator and Parser Overview

General Flowchart (see figure 1.)

The project, as developed, ended up with two distinct parts: The parser generator and a parser. The parser generator basically inputs a grammar definition and output five files:

1. A report file which describes the actions taken by the parser generator plus generates the action and goto tables.
2. Symbol table which defines the symbols in the grammar and the associated shorthand notations.
3. Production table which contains the internal representation of the productions.
4. The LALR size tables. This file contains the dimensions of the resultant LALR action and GOTO tables.
5. The LALR action and GOTO tables.

The parser was needed to test various grammar strings to see if they could be parsed correctly. Besides the four files generated from the Parser Generator (see figure 1) the parser has an input file of grammar strings to be tested and an output report file which gives the parse of the input strings or the partial parse if the string is not part of the grammar.

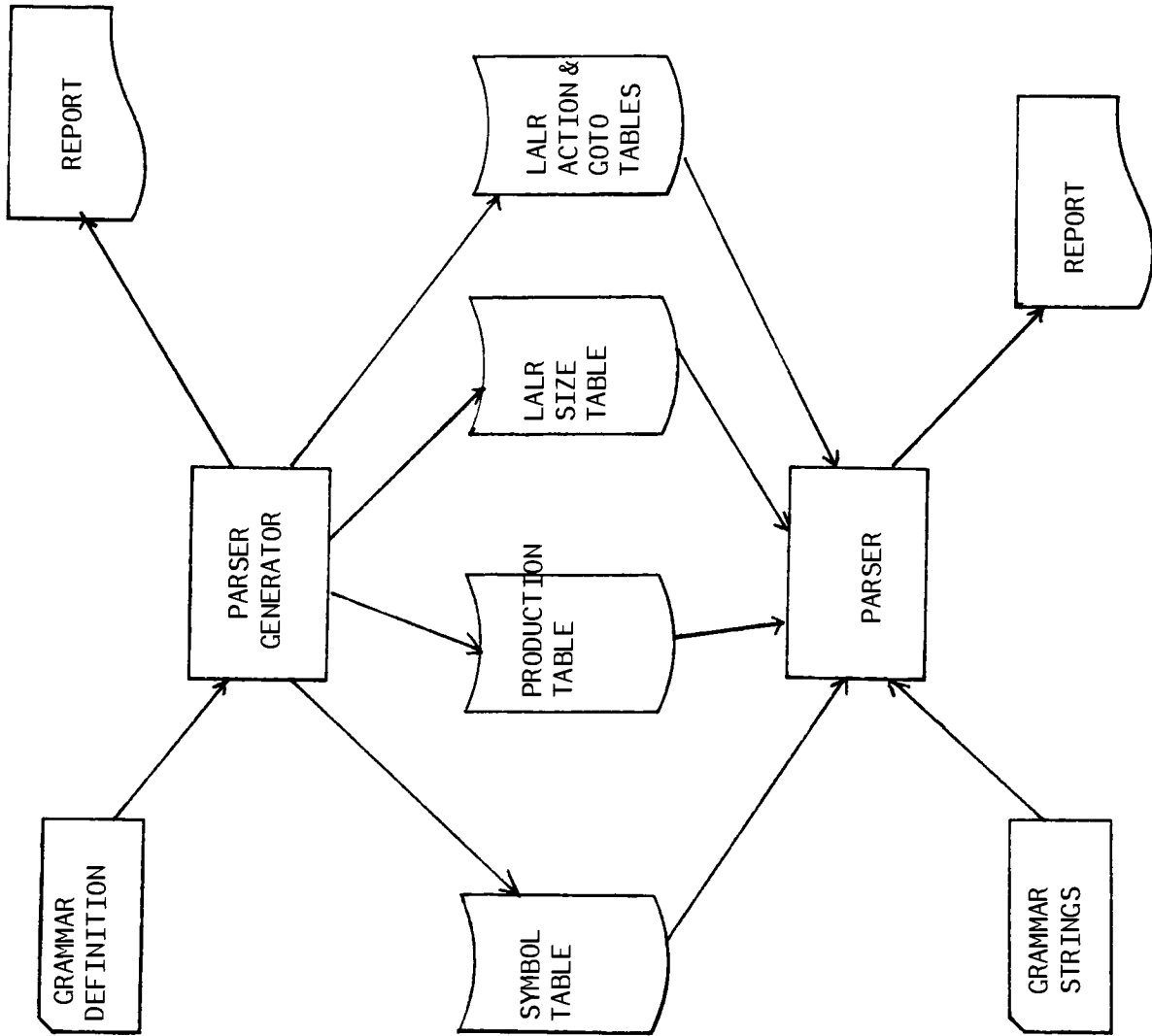


Figure 1

IV. LALR(1) Parser Generator

Important Data Structures

A. SYMBOLS (Appendix D p.2)

An array where each element has the following structure:

1. 1 character that is internal representation of symbol.
2. 1 bit that indicates whether symbol is a terminal ('0') or a non-terminal ('1').
3. Fixed binary number that gives the length of the external representation of symbol.
4. Complete identifier (up to 30 characters) of external identifier.

B. PROD (Appendix D p.2)

An array where each element is a production in the grammar and has the following structure:

1. The number of this production. Numbering starts from one and increases by one for each production defined in the grammar. The next element after the last production has this field set to zero as a delimiter.
2. One character (internal representation) that is the left part of the production.
3. A fixed binary number that is the number of symbols in the right part of the production.
4. Up to six characters of the right part of the production. Each character is the internal representation of a grammar symbol.

C. ITEMS (Appendix D p.2)

This array stores the generated items. Each element of the array has the following structure:

1. STATE is a fixed binary number that is the state number of the generated item.
2. NEXT which is a fixed binary number that is the element number in the ITEMS array of the next item in this state. A -1 indicates that this is the last item in this state.
3. PROD is a fixed binary number that references the production (in PROD array) that this item has been generated from.
4. POSIT is a fixed binary number that indicates the recognized part of this item.
5. LA is one character that is the look ahead symbol.

D. FIRSTS (Appendix D p.2)

This array contains, for each grammar symbol, the possible terminal symbols that can derive the object symbol. The symbols in this table are in the internal format. Each element of this array has the following structure.

1. NO_FIRSTS which is a binary number that is the number of first symbols contained in the next element.
2. FIRSTS which is in itself an array where each element is a possible FIRST symbol of the symbol in question.

Notice that this array does not have the object symbol itself stored. That is because the relative position of each symbol in array SYMBOLS (see above) is the same for array FIRSTS.

E. STATE_ITEMS (Appendix D p.3)

This array contains the number of generated items in each state. This array speeds up the process of determining identically equal states. By checking this array two states can be checked for equality in number of elements before checking for identically equal states.

F. GOTOS (Appendix D p.3)

This array has two dimensions; one for the grammar symbols and one for the parser states. When the parser does a shift action it then references this array to determine the next state to go to.

G. ACTION (Appendix D p.3)

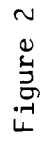
This array has two dimensions; one for the grammar symbols and one for the parser states. The elements of the array indicate whether the parser is to shift, reduce or accept the input grammar string. There are also invalid actions which could be the result of a non LR(1) grammar.

H. LALR_GOTO (Appendix D p.3)

Same as the GOTOS array except this goto array is for the LALR parsing table whereas the GOTOS array was for the LR(1) parsing table.

I. LALR_ACT (Appendix D p.3)

Same as the ACTION array except this action array is for the LALR parsing table whereas the ACTION array was for the LR(1) parsing table.



A. Main Program (Appendix D pp 6-9)

The main processing breaks down into the following steps:

1. Call GRAMMAR

This procedure (see full description below) reads in the grammar definition and sets up the terminals, non-terminals, production and internal grammar symbols.

2. The grammar in its internal format is printed.

3. Call COMP_FIRSTS computes the FIRST function for each grammar symbol and stores it in the array FIRSTS.

4. Call LR_ITEMS generates the LR(1) items for this grammar.

5. Call SR_TAB generates the LR(1) shift-reduce tables which are contained in arrays ACTION and GOTOS.

6. The LR(1) ACTION and GOTOS arrays are printed.

7. Call LALR generates the LALR(1) shift-reduce tables.

8. The LALR(1) shift-reduce tables are printed.

9. Finally the files needed by the parser, namely the symbol table, the production table, the LALR size table and the action and goto tables are written out to disk.

B. LR_ITEMS (Appendix D p.10)

The LR_ITEMS procedure causes the set of LR(1) items to be generated primarily by calling the procedures GOTO and CLOSURE. LR_ITEMS sets up the first item, adds it to the ITEMS array by calling ADD_ITEM and then calls CLOSURE to perform closure on that set of items. The result of this first call to CLOSURE is that state zero items have been created.

The procedure then makes repeated calls to GOTO (for the goto function) until all possible sets of items have been added to the ITEMS array.

Lastly the item table is printed.

C. CLOSURE (Appendix D p.11)

This procedure performs closure on the set of items contained in array ITEMS_SET where N_ITST is the number of items in the set. Each element of the array ITEMS_SET is an integer that references an item in the array ITEMS.

The main do loop executed is performed until all items in the set are exhausted. For each item in the set to be closed procedure FND_PROD and procedure MFIRST are called. FND_PROD returns a set of applicable productions and MFIRST returns a set of terminals that satisfy the lookahead function. New items are then created by taking all possible combinations of the productions and the lookahead symbols. These new items are added to the ITEMS array by procedure ADD_ITEM.

D. FND_PROD (Appendix D p.12)

This procedure is passed an item of the format $[A \Rightarrow \alpha.B\beta, a]$ and returns a set of productions in the grammar that are of the form: $B \Rightarrow \gamma$

In other words it returns all productions in the grammar whose left part is the first symbol of the unrecognized part of the item passed it.

The production numbers; of the returned productions, are stored in array PROD_# and the number of productions returned is stored in N_PROD.

E. MFIRST (Appendix D p.13)

This procedure is passed an item and it examines the second symbol of the unrecognized portion of the right part of the item. It then returns either:

1. The set of first symbols (by calling procedure FIRST) for that examined symbol, or
2. The lookahead symbol for the input items if the examined symbol is null.

F. FIRST (Appendix D p.13)

This procedure is passed a grammar symbol and it returns the set of first symbols for the input symbol. The set of first symbols for each grammar symbol has already been computed by procedure COMP_FIRSTS early in the program therefore this procedure is rather trivial.

G. PTR_SYM (Appendix D p.14)

This is a utility procedure that returns the index number to the passed grammar symbol. This procedure is used by procedures FIRST and COMP_FIRSTS because they have the grammar symbol and need the index to that symbol for use in array SYMBOLS.

H. ADD_TRM (Appendix D p.14)

This procedure is used by procedure COMP_FIRSTS to add symbols to set of first symbols in array FIRSTS. The input symbol is not added to the object symbols' list if it is already there.

I. TRM (Appendix D p.14)

This is a utility procedure that given a grammar symbol it returns a bit that indicates whether the passed symbol is a terminal ('0') or non-terminal ('1').

J. COMP_FIRSTS (Appendix D p.15)

This procedure computes the set of first symbols for each grammar symbol. The SYMBOLS array is setup such that all non-terminals appear before all terminal symbols. COMP_FIRSTS takes advantage of that structure by starting with the last entry in the SYMBOLS array and working towards the first entry. Each terminal symbol has as its set of first symbols itself. When the scan encounters a non-terminal

it looks for a production whose left part is that non-terminal then it gets the first symbol of the right part and obtains the set of first symbols for it to use for the object non-terminal.

By using this procedure to compute the set of first symbols for each grammar symbol at the beginning of the program, a considerable amount of processing is saved by not doing recursive calls later on.

K. ADD_ITEM (Appendix D p.16)

This procedure adds items to the ITEMS array. It is passed the following information:

1. The state that the item is to be added to.
2. The production number that it is based on.
3. A number that represents the number of symbols recognized in the right part of the production.
4. The lookahead symbol.

The procedure scans the ITEMS array to find the last item of appropriate state and chains it to the new item to be added and then it adds the new item.

A count is kept for the number of items in each state in array STATE_ITEMS; this count is updated.

L. GOTO (Appendix D pp. 17-18)

The GOTO procedure computes the goto function for construction of the LR(1) sets of items. It is provided upon entry the state it is to perform the goto function for and the symbol it is to look for. The third value passed is the relative position of the symbol within the symbol table.

The procedure first finds the first item, within the ITEMS array, of the passed state. Secondly, the procedure copies out all items within that state whose next symbol, beyond the recognized part, is the passed symbol. The items that are copied are placed in two places.

1. in the ITEMS array under a new state and
2. in ITEMS_SET so that closure can be done on this new state.

A test is made on variable N_ITST, which is the number of items to close, and if its zero the procedure is exited. If it is not zero then there are items to perform closure on and procedure CLOSURE is called.

Because of the iterative nature of the LR(1) item generation procedure it is possible to generate duplicate GOTO states therefore this procedure then does a check to see if the newly created set of items has indeed already been added to the universal set. Since the number of generated items in each state is saved in array STATE_ITEMS a quick check of that array will tell which of the states have equal number of items as compared to the new goto state. For those states that have an equal number of items as the new state they must be checked for identically equal by doing an item by item comparison.

If the new state is unique, then it is left as created and the element selected by the input state and symbol of the two dimensional GOTOS array is updated with the number of the newly created state. However, if the new state had already existed then that state number is placed in the GOTOS element as referenced by the input state and symbol.

M. SR_TAB (Appendix D p.19)

This procedure creates the canonical LR parsing table for the input grammar. The technique used is outlined in algorithm 6.3 of Aho and Ullman (1977).

Each item is examined in the ITEMS array and depending upon the item it will result in an array element of ACTION to be filled in with a code that signifies Accept, Reduce or Shift. Accept is the case if the item is the completely recognized first production with no more look ahead characters. Reduce is the case if the item is a completely recognized production. Shift is the case for any item that is a partially recognized production whose next symbol is a terminal.

Before the appropriate action is stored in array ACTION the selected element is examined to see if it has already been filled in with a different action. If it has been then we have an action conflict which means that this grammar is not an LR(1) grammar.

N. GRAMMAR (Appendix D pp. 20-21)

The purpose of this procedure is to read in the grammar definition and set up the internal one character grammar symbols.

The rules for the grammar definition are as follows:

1. Each record must contain a valid 4 character transaction code of:

NONT for Nonterminal
TERM for Terminal, or
PROD for Production.

2. The first non-terminal defined is considered to be the start symbol.
3. Terminals and non-terminals are separated by at least one blank and have a maximum length of 30 characters.
4. Only one production allowed per record and the left part is separated from the right part by a colon.

For transaction codes NONT and TERM the records are scanned picking out each grammar symbol and saved in the SYMBOLS array. Array REP_SYM is used to sequentially assign a single character internal grammar symbol to each input terminal and non-terminal.

Each production definition is scanned for its grammar symbols and these are converted into the internal format and the whole production stored into the PROD array.

O. INT (Appendix D p.22)

This is a utility procedure which converts a full grammar symbol to its internal representation. The conversion is done by scanning the SYMBOLS array.

P. ALTER (Appendix D p.22)

This procedure is used in the conversion process from the LR(1) parsing tables to the LALR parsing tables. In the transformation process when two states are deemed equal then this is called to merge the action entries of the two states into one.

The procedure is passed two state numbers and it merges the action entries from the second state into the first only if the action entry of the first was undefined.

Q. LALR (Appendix D pp.23-24)

This procedure does the conversion from the LR(1) parsing tables to the LALR parsing tables by examining the canonical set of LR(1) items. If two sets of items, in the LR(1) items, are equal without respect to their lookahead symbols then they can be combined.

This procedure does that processing in the following steps:

1. Looks at all combinations of the states taken two at a time to determine equality. Equality of any two states is determined by procedure IDENT. If equality is determined then procedure ALTER is called and array TRANSFORM is updated to reflect the combination of the two states.
2. By looking at array TRANSFORM a list of deleted states is created.
3. The end result is to cause a decrease in the number of states and in the process of state deletion there are valid states that are beyond the state limits of the LALR tables. This portion reassigns the valid states to the places previously occupied by the deleted states. This necessitates moving the appropriate vectors of arrays GOTOS and ACTION to the deleted places plus transforming any reference to a deleted state to a valid state.

4. Arrays LALR_GOTO and LALR_ACT are dynamically allocated with the new sizes.
5. The newly allocated LALR action and goto tables entries are filled in from the transformed LR(1) action and goto tables.

R. IDENT (Appendix D p.25)

This procedure is called by procedure LALR to determine if two states are equal so that they can be reduced to one state for the LALR parsing tables. Equality is determined by checking that each item in one state has an equal item in the second state (meaning that they both reference the same production and have the same number of recognized symbols) and vice versa.

The procedure returns a bit value of one if the states are equal and zero if not.

V. Shift-Reduce Parser

The purpose of this program is to take strings from the grammar defined to the parser generator and to parse them. If they do not parse correctly then we can assume that the strings are not from the grammar.

Besides inputting the parsing tables and the production tables, which are necessary to perform the basic function of this program, the symbol table from the parser generator is also input. By doing this it was not necessary to write a lexical analyzer for each grammar tested. All the information necessary for lexical analysis is contained within the symbol table therefore we have in effect a general lexical analysis within this program.

Important Data Structures

SYMBOLS, PROD, GOTOS, and ACTION are the important data structures to this program and they are the same ones as described under the parser generator program.

The GOTOS and ACTION arrays can be either LR(1) or LALR. It makes no difference to this program since it is only concerned with using the actions described within each to perform its function.

Program Description

A. Main Procedure (Appendix E pp.4-6)

1. Declarations, initializations, and program housekeeping are done first.
2. After all the files are opened, SYMTAB, PROD and LALRSZ arrays are initialized by reading the appropriate files.
3. The LALRSZ array gives the dimension of the GOTOS and ACTION arrays so that they can be dynamically allocated. Those arrays are initialized by reading file LALRTAB.
4. The first symbol is obtained by calling procedure GETSYM.
5. Based upon the current state on top of the stack and the input symbol one of the following actions, based upon the entry in the ACTION array, takes place:
 - a. Invalid action which results in an invalid string message being produced and the string being flushed.
 - b. Accept action which means the string was correctly parsed.
 - c. Shift which causes the input symbol to be pushed onto the stack plus the appropriate goto state and another symbol obtained.
 - d. Reduce which causes reduction via a grammar production. The top 2n objects on the stack are popped off and replaced by the left part of the production and a goto state, where n is the number of symbols in the right part of the production.

Parser Procedure Relationships

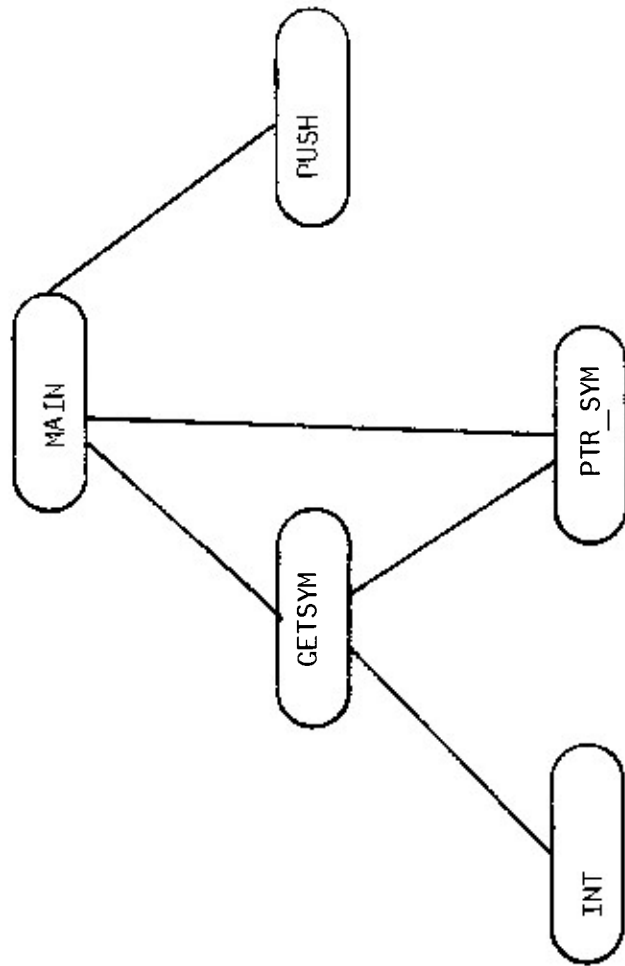


Figure 3

B. PUSH (Appendix E p.7)

This is a utility procedure that takes a value and puts it on the stack.

C. PTR_SYM (Appendix E p.7)

This utility procedure returns an integer index to the passed grammar symbol. The grammar symbol needs to be converted to a relative number so that it can be used to index into GOTOS and ACTION arrays.

D. GETSYM (Appendix E p.8)

This procedure reads in a grammar string, converts the input grammar symbols to their internal representation and returns the next symbol to the calling procedure. When it runs out of symbols in the current string it obtains another record to process.

E. INT (Appendix E p.9)

This is a utility procedure which converts a full grammar symbol to its internal representation. The conversion is done by scanning the SYMBOLS array.

VI. Example Explanation

Three examples of the automatic parser output and symtactical analysis of selected input strings are presented in appendices A, B, and C.

Specifically, Appendix A uses the same grammar example that Aho and Ullman (1977) use in their example 6.10. Appendix B uses the grammar for standard algebraic expressions and Appendix C uses a subset of the EASY language defined by Wetherell (1978) for IF THEN ELSE statements.

Each appendix presents a different LR(1) grammar and is organized as follows:

1. The input grammar definition specifying the grammar nonterminals, terminals and production rules.
2. The internal grammar definition used by the parser generator. The purpose of having an internal grammar is so that grammar symbols can be reduced to single characters which facilitates the processing.
3. The computed first symbols for each grammar symbol.
4. The generated LR(1) items.
5. The LR(1) goto and action tables.
6. The LR(1) to LALR(1) transformation which shows which LR(1) states were deleted and the final state transformation from LR(1) to LALR(1).
7. The LALR(1) goto and action tables.
8. Examples of the parsing action taken on selected strings from the grammar.

In order to analyze input grammar strings the parser uses the method presented by Aho and Ullman (1972) in their algorithm 5.7 which follows:

ALGORITHM 5.7

LR(k) parsing algorithm.

Input. A set T of LR(k) tables for an LR(k) grammar $G=(N, \Sigma, P, S,)$, with $T_0 \in T$ designated as the initial table, and an input string $z \in \Sigma^*$, which is to be parsed.

Output. If $z \in L(G)$, the right parse of G . Otherwise, an error indication.

Method. Perform steps (1) and (2) until acceptance occurs or an error is encountered. If acceptance occurs, the string in the output buffer is the right parse of z .

- (1) The lookahead string u , consisting of the next k input symbols, is determined.
- (2) The parsing action function f of the table on top of the pushdown list is applied to the lookahead string u .
 - (a) If $f(u)=\text{shift}$, then the next input symbol, say a , is removed from the input and shifted onto the pushdown list. The goto function g of the table on top of the pushdown list is applied to a to determine the new table to be placed on top of the pushdown list. We then return to step (1). If there is no next input symbol or $g(a)$ is undefined, halt and declare error.
 - (b) If $f(u)=\text{reduce } i$ and production i is $A \rightarrow \alpha$, then $2|\alpha|$ symbols* are removed from the top of the pushdown list, and production number i is placed in the input buffer. A new table T' is then exposed as the top table of the pushdown list, and the goto function of T' is applied to A to determine the next table to be placed on top of the pushdown list. We place A and this new table on top of the pushdown list and return to step (1).
 - (c) If $f(u)=\text{error}$, we halt parsing (and, in practice, transfer to an error recovery routine).
 - (d) If $f(u)=\text{accept}$, we halt and declare the string in the output buffer to be the right parse of the original input string.

*If $\alpha = X_m \dots X_r$, at this point the top of the pushdown list will be of the form $T_0 X_1 T_1 X_2 T_2 \dots X_r T_r$. Removing $2|\alpha|$ symbols removes the handle from the top of the pushdown list along with any intervening LR tables.

To illustrate how this technique works let's consider the example on p. A5 considering the following:

1. Use algorithm 5.7 in Aho and Ullman (1972) where $k=1$.
2. In the example on p. A5:
 $D = E$ means that D is the external grammar symbol and E is the internal grammar symbol.
3. Use the LALR(1) goto and action tables on p. A4.
4. The initial state is zero.
5. $\$$ represents the end of the input string.

STACK	INPUT	COMMENT
0	EE\$	
OE4	E\$	Stack input character and state 4
OC2	E\$	Reduce using Production number 4
OC2E4	\$	Stack input character and state 4
OC2C5	\$	Reduce using production number 4
OB1	\$	Reduce using production number 2

With state 1 on top of the stack and no more input the LALR(1) Action function says accept.

To obtain the rightmost parse of a particular grammar sentence one just has to apply the productions in reverse order from any example. To illustrate let's apply the production numbers in reverse order for the example on p. B10.

Production

1	GRAMMAR →	EXPRESSION
3	→	TERM
4	→	TERM * FACTOR
7	→	TERM * A
5	→	FACTOR * A
6	→	(EXPRESSION) * A
2	→	(EXPRESSION + TERM) * A
5	→	(EXPRESSION + FACTOR) * A
7	→	(EXPRESSION + A) * A
3	→	(TERM + A) * A
5	→	(FACTOR + A) * A
7	→	(A + A) * A

Which does indeed yield the expected grammar string.

All the other examples presented are valid input strings for their particular grammars except for the following ones:

The input string shown on p. A6 fails because the grammar defined generates strings in the format of B^mDB^nD where m does not have to equal n and either m or n can equal zero. Therefore, the string BD must not be from this grammar.

The input string shown on p. B9 fails because the operators $*$ and $+$ cannot be next to each other.

The input string shown on p. C6 fails because in the defined grammar the IF THEN ELSE statement must be terminated with FI.

VII. Conclusion

Problems Encountered

During the implementation there were several notable problem areas. The first problem was that the end result (the parser generator) was quite large taking approximately 800 PL/1 statements. The algorithms and techniques presented in pseudo code by Aho and Ullman expanded significantly when actually set down in a real language.

The second problem encountered had to do with the computation of FIRST for each grammar symbol. The first attempt called for the FIRST calculation to be done in the CLOSURE procedure each time CLOSURE was called. This technique proved to be wasteful and overly complicated. Using the recursive technique was academically appealing but unnecessary. A far simpler and direct approach was to examine the grammar productions and calculate FIRST for each grammar symbol prior to generating the items.

The third problem encountered was dealing with duplicate generated states. Due to the iterative nature of generating the LR(1) items it is possible to generate a new state that is identically equal to a previous state. Therefore, when each new state is generated it must be checked against each previously generated state to ensure its uniqueness. Although some steps were taken to optimize this process it could not be eliminated entirely.

The last problem encountered was during array compression after the LALR(1) array dimensions had been established. In order to achieve the true benefit of the LALR(1) parsing tables, the arrays must contain only valid states. Generally there are more LR(1) states than the resultant LALR(1) states and the deleted states are, unfortunately, not the last column vectors in the arrays. Therefore, it becomes necessary to move valid states, that are beyond the LALR(1) state limits, to deleted LR(1) states. This process includes moving the valid states into the deleted states and changing any reference to the moved state to its new state number.

Possible Extensions

There are two extensions that can be made to this implementation. The first one deals with generating the LALR(1) tables directly instead of going through the intermediate process of generating the LR(1) tables. Using the intermediate step of producing the LR(1) tables is instructive on the technique but in reality takes up much more space than is necessary. Generating the LALR(1) tables directly would be much more efficient.

The second extension that could be investigated is in the area of error processing. In the current implementation when an error is detected in the parsing of a sentence from the grammar the sentence is flushed and the next sentence is processed. Instead of flushing the sentence an error processing function could be invoked to attempt to use the partial parse and perhaps "correct" the sentence so that it is syntactically correct.

Concluding Remarks

There are four benefits to the parser generator described in this paper. The first benefit is that the LALR(1) parser generator can accept a broad class of grammars. Secondly it detects syntax errors at the earliest possible time in the parsing process. Thirdly the LALR(1) parsing tables are space efficient. Lastly, the generated output of the parser generator greatly simplifies the parser. The parser's actions are determined in a straight forward manner based upon the table entries. In a real implementation the parsing tables could be generated ahead of time and the parser could be directed to use the appropriate set of tables. The end result is that one parser could parse sentences from several different grammars.

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```
NONT G S C
TERM B D
PROD G : S
PROD S : C C
PROD C : B C
PROD C : D
```

INTERNAL GRAMMAR:

NONTERMINALS: A B C

TERMINALS: D E

START SYMBOL: A

PRODUCTIONS:

```
( 1 ) A -> B
( 2 ) B -> CC
( 3 ) C -> DC
( 4 ) C -> E
```

COMPUTED FIRSTS:

```
A ED
B ED
C ED
D D
E E
```

GENERATED ITEMS

NO.	STATE	ITEMS
1	0	A -> .B \$
2	00	AB -> .CCC \$
3	000	CCC -> .DDC \$
4	0000	DDC -> .EEC \$
5	00000	EEC -> .FEE \$
6	000000	FFF -> .GFF \$
7	0000001	AG -> .C \$
8	00000011	AC -> .CC \$
9	000000112	ACC -> .CCC \$
10	0000001122	CCCC -> .EEEE \$
11	00000011223	CCCC -> .EEC \$
12	000000112233	CCCC -> .EEC \$
13	0000001122333	CCCC -> .EEC \$
14	00000011223333	CCCC -> .EEC \$
15	000000112233333	CCCC -> .EEC \$
16	0000001122333333	CCCC -> .EEC \$
17	00000011223333334	CCCC -> .EEC \$
18	000000112233333344	CCCC -> .EEC \$
19	0000001122333333445	CCCC -> .EEC \$
20	00000011223333334455	CCCC -> .EEC \$
21	000000112233333344556	CCCC -> .EEC \$
22	0000001122333333445567	CCCC -> .EEC \$
23	00000011223333334455678	CCCC -> .EEC \$
24	000000112233333344556789	CCCC -> .EEC \$
25	0000001122333333445567890	CCCC -> .EEC \$
26	00000011223333334455678901	CCCC -> .EEC \$

LR(1) GOTO FUNCTION:

	0	1	2	3	4	5	6	7	8	9
A										
B	1									
C	2		5	8			9			
D	3		6	3			6			
E	4		7	4			7			

LR(1) ACTION FUNCTION:

	0	1	2	3	4	5	6	7	8	9
A										
B										
C										
D	S		S	S	4		S		3	
E	S		S	S	4		S		3	
\$		A				2		4		3

LALR TRANSFORMATION

LR(1) STATES TO DELETE:

6 7 9

LR TO LALR STATE TRANSFORM:

LR	0	1	2	3	4	5	6	7	8	9
LALR	0	1	2	3	4	5	3	4	6	6

LALR(1) GOTO FUNCTION:

	0	1	2	3	4	5	6
--	---	---	---	---	---	---	---

A	1						
B	2		5	6			
C	3		3	3			
D	3		3	3			
E	4		4	4			

LALR(1) ACTION FUNCTION:

	0	1	2	3	4	5	6
--	---	---	---	---	---	---	---

A							
B							
C							
D	S		S	S	4		3
E	S		S	S	4		3
\$		A			4	2	3

INPUT STRING:

D D

```
D => E
STACK INPUT CHARACTER AND STATE 4
D => E
REDUCE USING PRODUCTION NO. 4
STACK INPUT CHARACTER AND STATE 4
=> $
REDUCE USING PRODUCTION NO. 4
REDUCE USING PRODUCTION NO. 2
REDUCE USING PRODUCTION NO. 1
VALID INPUT STRING
```

— INPUT STRING:
B D
B => D
STACK INPUT CHARACTER AND STATE 3
D => E
STACK INPUT CHARACTER AND STATE 4
=> \$
REDUCE USING PRODUCTION NO. 4
REDUCE USING PRODUCTION NO. 3
INVALID STRING FOR THIS GRAMMAR
INPUT STRING FLUSHED

INPUT STRING:

B B D B D

```

B => D
    STACK INPUT CHARACTER AND STATE    3
B => D
    STACK INPUT CHARACTER AND STATE    3
D => E
    STACK INPUT CHARACTER AND STATE    4
B => D
    REDUCE USING PRODUCTION NO.      4
    REDUCE USING PRODUCTION NO.      3
    REDUCE USING PRODUCTION NO.      3
    STACK INPUT CHARACTER AND STATE    3
D => E
    STACK INPUT CHARACTER AND STATE    4
=> $
    REDUCE USING PRODUCTION NO.      4
    REDUCE USING PRODUCTION NO.      3
    REDUCE USING PRODUCTION NO.      2
    REDUCE USING PRODUCTION NO.      1
VALID INPUT STRING

```



```

NONT GRAMMAR EXPRESSION TERM FACTOR
TERM + * ( ) A
PROD GRAMMAR : EXPRESSION
PROD EXPRESSION : EXPRESSION + TERM
PROD EXPRESSION : TERM
PROD TERM : TERM * FACTOR
PROD TERM : FACTOR
PROD FACTOR : ( EXPRESSION )
PROD FACTOR : A

```

```

INTERNAL GRAMMAR:
NONTERMINALS: A B C D
TERMINALS: E F G H I
START SYMBOL: A
PRODUCTIONS:
( 1) A -> B
( 2) B -> BEC
( 3) B -> C
( 4) C -> CFD
( 5) C -> D
( 6) D -> GBH
( 7) D -> I

```

COMPUTED FIRSTS:

```

A IG
B IG
C IG
D IG
E E
F F
G G
H H
I I

```

GENERATED ITEMS

NO.	STATE	ITEMS
1	0	A
2	0	V
3	0	V
4	0	V
5	0	V
6	0	V
7	0	V
8	0	V
9	0	V
10	0	V
11	0	V
12	0	V
13	0	V
14	0	V
15	0	V
16	0	V
17	0	V
18	1	A
19	1	V
20	1	V
21	2	V
22	2	V
23	2	V
24	2	V
25	2	V
26	3	V
27	3	V
28	3	V
29	4	V
30	4	V
31	4	V
32	4	V
33	4	V
34	4	V
35	4	V
36	4	V
37	4	V
38	4	V
39	4	V
40	4	V
41	4	V
42	4	V
43	4	V
44	4	V
45	4	V
46	4	V
47	4	V
48	5	V
49	5	V
50	5	V
51	6	V
52	6	V
53	6	V
54	6	V
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57	6	V

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[illegible]

LR(1) GOTO FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21
A																						
B	1				8																	
C	2				9		13															
D	3				10		3	14														
E		6							15													
F			7							17												
G	4				11		4	4														
H									16													
I	5				12		5	5														

LR(1) ACTION FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	
A																							
B																							
C																							
D		S	3	5		7			S	3	5		7	2	4		6		S	2	4	6	
E			S	5		7				S	5		7	S	4		6			S	4	6	
F	S				S		S	S				S				S	6	S					
G									S	3	5		7					S		S	2	4	6
H	S				S		S	S				S				S			S				
I		A	3	5		7								2	4		6						

LALR TRANSFORMATION

LR(1) STATES TO DELETE:

9 10 11 12 15 17 18 19 20 21

LR TO LALR STATE TRANSFORM:

LR	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21
LALR	0	1	2	3	4	5	6	7	8	2	3	4	5	9	10	6	11	7	8	9	10	11

LALR(1) GOTO FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11
--	---	---	---	---	---	---	---	---	---	---	----	----

A												
B	1				8							
C	2				2		9					
D	3				3		3	10				
E		6							6			
F			7							7		
G	4				4		4	4				
H											11	
I	5				5		5	5				

LALR(1) ACTION FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11
--	---	---	---	---	---	---	---	---	---	---	----	----

A												
B												
C												
D												
E		S	3	5		7			S	2	4	6
F			S	5		7			S	S	4	6
G	S				S		S	S				
H			3	5		7			S	2	4	6
I	S				S		S	S				
\$		A	3	5		7				2	4	6

INPUT STRING:

A + A + A * A

```

A => I
    STACK INPUT CHARACTER AND STATE      5
+ => E
    REDUCE USING PRODUCTION NO.      7
    REDUCE USING PRODUCTION NO.      5
    REDUCE USING PRODUCTION NO.      3
    STACK INPUT CHARACTER AND STATE      6
A => I
    STACK INPUT CHARACTER AND STATE      5
+ => E
    REDUCE USING PRODUCTION NO.      7
    REDUCE USING PRODUCTION NO.      5
    REDUCE USING PRODUCTION NO.      2
    STACK INPUT CHARACTER AND STATE      6
A => I
    STACK INPUT CHARACTER AND STATE      5
* => F
    REDUCE USING PRODUCTION NO.      7
    REDUCE USING PRODUCTION NO.      5
    STACK INPUT CHARACTER AND STATE      7
A => I
    STACK INPUT CHARACTER AND STATE      5
=> $
    REDUCE USING PRODUCTION NO.      7
    REDUCE USING PRODUCTION NO.      4
    REDUCE USING PRODUCTION NO.      2
    REDUCE USING PRODUCTION NO.      1
VALID INPUT STRING

```

INPUT STRING:

A * A + A + A

```

A => I
    STACK INPUT CHARACTER AND STATE    5
* => F
    REDUCE USING PRODUCTION NO.    7
    REDUCE USING PRODUCTION NO.    5
    STACK INPUT CHARACTER AND STATE    7
A => I
    STACK INPUT CHARACTER AND STATE    5
+ => E
    REDUCE USING PRODUCTION NO.    7
    REDUCE USING PRODUCTION NO.    4
    REDUCE USING PRODUCTION NO.    3
    STACK INPUT CHARACTER AND STATE    6
A => I
    STACK INPUT CHARACTER AND STATE    5
+ => E
    REDUCE USING PRODUCTION NO.    7
    REDUCE USING PRODUCTION NO.    5
    REDUCE USING PRODUCTION NO.    2
    STACK INPUT CHARACTER AND STATE    6
A => I
    STACK INPUT CHARACTER AND STATE    5
=> $
    REDUCE USING PRODUCTION NO.    7
    REDUCE USING PRODUCTION NO.    5
    REDUCE USING PRODUCTION NO.    2
    REDUCE USING PRODUCTION NO.    1
VALID INPUT STRING

```


INPUT STRING:

A * A + A * + A

```

A => I
  STACK INPUT CHARACTER AND STATE    5
* => F
  REDUCE USING PRODUCTION NO.    7
  REDUCE USING PRODUCTION NO.    5
  STACK INPUT CHARACTER AND STATE    7
A => I
  STACK INPUT CHARACTER AND STATE    5
+ => E
  REDUCE USING PRODUCTION NO.    7
  REDUCE USING PRODUCTION NO.    4
  REDUCE USING PRODUCTION NO.    3
  STACK INPUT CHARACTER AND STATE    6
A => I
  STACK INPUT CHARACTER AND STATE    5
* => F
  REDUCE USING PRODUCTION NO.    7
  REDUCE USING PRODUCTION NO.    5
  STACK INPUT CHARACTER AND STATE    7
+ => E
  INVALID STRING FOR THIS GRAMMAR
  INPUT STRING FLUSHED
A => I
=> $

```

INPUT STRING:

(A + A) * A

```
( => G
  STACK INPUT CHARACTER AND STATE 4
A => I
  STACK INPUT CHARACTER AND STATE 5
+ => E
  REDUCE USING PRODUCTION NO. 7
  REDUCE USING PRODUCTION NO. 5
  REDUCE USING PRODUCTION NO. 3
  STACK INPUT CHARACTER AND STATE 6
A => I
  STACK INPUT CHARACTER AND STATE 5
) => H
  REDUCE USING PRODUCTION NO. 7
  REDUCE USING PRODUCTION NO. 5
  REDUCE USING PRODUCTION NO. 2
  STACK INPUT CHARACTER AND STATE 11
* => F
  REDUCE USING PRODUCTION NO. 6
  REDUCE USING PRODUCTION NO. 5
  STACK INPUT CHARACTER AND STATE 7
A => I
  STACK INPUT CHARACTER AND STATE 5
=> $
  REDUCE USING PRODUCTION NO. 7
  REDUCE USING PRODUCTION NO. 4
  REDUCE USING PRODUCTION NO. 3
  REDUCE USING PRODUCTION NO. 1
VALID INPUT STRING
```

INPUT STRING:

A + A * A

```

A => I
  STACK INPUT CHARACTER AND STATE    5
+ => E
  REDUCE USING PRODUCTION NO.    7
  REDUCE USING PRODUCTION NO.    5
  REDUCE USING PRODUCTION NO.    3
  STACK INPUT CHARACTER AND STATE    6
A => I
  STACK INPUT CHARACTER AND STATE    5
* => F
  REDUCE USING PRODUCTION NO.    7
  REDUCE USING PRODUCTION NO.    5
  STACK INPUT CHARACTER AND STATE    7
A => I
  STACK INPUT CHARACTER AND STATE    5
=> $
  REDUCE USING PRODUCTION NO.    7
  REDUCE USING PRODUCTION NO.    4
  REDUCE USING PRODUCTION NO.    2
  REDUCE USING PRODUCTION NO.    1
VALID INPUT STRING

```

```

NONT <COND_STMT> <SCS> <COND_CL>
NONT <TRUE> <FALSE> <COND_BODY> <ELSE>
TERM SEG_BODY FI IF EXPR THEN ELSE
PROD <COND_STMT> : <SCS>
PROD <SCS> : <COND_CL> <TRUE> FI
PROD <SCS> : <COND_CL> <TRUE> <FALSE> FI
PROD <COND_CL> : IF EXPR
PROD <TRUE> : THEN <COND_BODY>
PROD <FALSE> : <ELSE> <COND_BODY>
PROD <ELSE> : ELSE
PROD <COND_BODY> : SEG_BODY

```

INTERNAL GRAMMAR:

NONTERMINALS: A B C D E F G

TERMINALS: H I J K L M

START SYMBOL: A

PRODUCTIONS:

- (1) A -> B
- (2) B -> CDI
- (3) B -> CDEI
- (4) C -> JK
- (5) D -> LF
- (6) E -> GF
- (7) G -> M
- (8) F -> H

COMPUTED FIRSTS:

```

A J
B J
C J
D E J
E F H
G H
H I
I J
J K
K L
L M
M

```

GENERATED ITEMS

NO.	STATE	ITEMS
1	0	A -> .B,
2	0	B -> .COI,
3	0	E -> .COI,
4	0	O -> .K, \$
5	1	A -> .COI,
6	2	B -> .COI,
7	2	B -> .COI,
8	2	O -> .COI,
9	2	O -> .COI,
10	3	O -> .COI,
11	4	B -> .COI,
12	4	B -> .COI,
13	4	B -> .COI,
14	4	G -> .COI,
15	5	O -> .COI,
16	5	O -> .COI,
17	5	O -> .COI,
18	5	O -> .COI,
19	6	O -> .COI,
20	7	O -> .COI,
21	8	O -> .COI,
22	8	O -> .COI,
23	9	O -> .COI,
24	10	O -> .COI,
25	11	O -> .COI,
26	11	O -> .COI,
27	12	O -> .COI,
28	12	O -> .COI,
29	13	O -> .COI,
30	14	O -> .COI,
31	15	O -> .COI,

LR(1) GOTO FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
A																
B																
C	1															
D	2															
E			4													
F					7											
G					8	11			14							
H					9	12			15							
I								13								
J	3															
K																
L			5	6												
M					10											

LR(1) ACTION FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
A																
B																
C																
D																
E																
F																
G																
H																
I																
J	S				S	S		S	S		7	5	8		6	8
K																
L			S	S												
M					S		4									
\$		A								2		5	8		3	

- LALR TRANSFORMATION

LR(1) STATES TO DELETE:

15

LR TO LALR STATE TRANSFORM:

LR	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
LALR	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	12

LALR(1) GOTO FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
A															
B	1														
C	2														
D			4												
E					7										
F					11			14							
G					8										
H					12			12							
I					9		13								
J	3														
K					6										
L			5												
M					10										

LALR(1) ACTION FUNCTION:

	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
A															
B															
C															
D															
E															
F															
G															
H															
I															
J	S				S	S		S		7	5	8		6	
K															
L															
M															
\$		A			S		4			2	5	8		3	

```

INPUT STRING:
IF EXPR THEN SEG_BODY FI

IF => J
  STACK INPUT CHARACTER AND STATE 3
EXPR => K
  STACK INPUT CHARACTER AND STATE 6
THEN => L
  REDUCE USING PRODUCTION NO. 4
  STACK INPUT CHARACTER AND STATE 5
SEG_BODY => H
  STACK INPUT CHARACTER AND STATE 12
FI => I
  REDUCE USING PRODUCTION NO. 6
  REDUCE USING PRODUCTION NO. 5
  STACK INPUT CHARACTER AND STATE 9
=> $
  REDUCE USING PRODUCTION NO. 2
  REDUCE USING PRODUCTION NO. 1
  VALID INPUT STRING

```



```
INPUT STRING:
IF EXPR THEN SEG_BODY

IF => J
  STACK INPUT CHARACTER AND STATE 3
EXPR => K
  STACK INPUT CHARACTER AND STATE 6
THEN => L
  REDUCE USING PRODUCTION NO. 4
  STACK INPUT CHARACTER AND STATE 5
SEG_BODY => H
  STACK INPUT CHARACTER AND STATE 12
=> $
  INVALID STRING FOR THIS GRAMMAR
  INPUT STRING FLUSHED
```

```

INPUT STRING:
IF EXPR THEN SEG_BODY ELSE SEG_BODY FI

IF => J
  STACK INPUT CHARACTER AND STATE    3
EXPR => K
  STACK INPUT CHARACTER AND STATE    6
THEN => L
  REDUCE USING PRODUCTION NO.        4
  STACK INPUT CHARACTER AND STATE    5
SEG_BODY => H
  STACK INPUT CHARACTER AND STATE   12
ELSE => M
  REDUCE USING PRODUCTION NO.        8
  REDUCE USING PRODUCTION NO.        5
  STACK INPUT CHARACTER AND STATE   10
SEG_BODY => H
  REDUCE USING PRODUCTION NO.        7
  STACK INPUT CHARACTER AND STATE   12
FI => I
  REDUCE USING PRODUCTION NO.        8
  REDUCE USING PRODUCTION NO.        6
  STACK INPUT CHARACTER AND STATE   13
=> $
  REDUCE USING PRODUCTION NO.        3
  REDUCE USING PRODUCTION NO.        1
  VALID INPUT STRING

```

PL/I OPTIMIZING COMPILER

OPTIONS SPECIFIED

MAR(2 72 1) NUM GN TERM;

OPTIONS USED

AGGREGATE
COMPILE
GONUMBER
INSOURCE
LMESSAGE
MAP
NEST
NUMBER
OBJECT
OFFSET
OPTIONS
SOURCE
STORAGE

ADCCOUNT
MODECK
NOESU
NOELUM
NOGOSTMT
NOIMPRECISE
NOINCLUDE
NOINTERRUPT
NULLST
NUSYN
XREF(SHORT)
TERMINAL INUAGGREGATE,
NOATTRIBUTES,
NOESU,
NOINSOURCE,
NOLIST,
NOMAP,
NOOFFSET,
NOOPTIONS,
NOSOURCE,
NOSTORAGE,
NOXREF)

VERSION 1 RELEASE 3.0 P/F 08

TIME: 07.55.10

DATE: 28 FEB 80

PAGE 1

/* PARSER GENERATOR */

PL/I OPTIMIZING COMPILER

NUMBER LEV NT

```

450 1 0 DCL ITEMS_SET(100) FIXED BIN(15); /* SET OF ITEMS USED BY CLOSURE */
460 1 0 DCL N_ITEM FIXED BIN(15); /* NO. OF ITEMS IN ITEMS_SET */
470 1 0 DCL PROO_# (10) FIXED BIN(15); /* USED BY FNO_PROO */
480 1 0 DCL N_PROO FIXED BIN(15); /* NO. OF PRODUCTIONS IN PROO_# */
490 1 0 DCL STATE_ITEMS(10:29) FIXED BIN(15) STATIC INIT((30)0);
      /* NO. OF ITEMS IN EACH STATE */
510 1 0 DCL GOTOS(15,0:29) FIXED BIN(15) STATIC INIT((450)-1);
      /* (M,N) WHERE M = # GRAMMAR SYMBOLS AND
      /* N = # STATES.
540 1 0 DCL ACTION(15,0:29) FIXED BIN(15) STATIC INIT((450)-1);
      /* (M,N) WHERE M = # GRAMMAR SYMBOLS AND
      /* N = # STATES.
570 1 0 DCL 1 LALR_SZ, /* LALR GOTO & ACTION ARRAYS */
      2 M FIXED BIN(15); /* NO. OF GRAMMAR SYMBOLS */
      2 N FIXED BIN(15); /* NO. OF STATES */
600 1 0 DCL LALR_GOTO(LALR_SZ,M,0:LALR_SZ,N) FIXED BIN(15) CTL;
610 1 0 DCL LALR_ACT(LALR_SZ,M,0:LALR_SZ,N) FIXED BIN(15) CTL;
ZEF00450
ZEF00460
ZEF00470
ZEF00480
ZEF00490
ZEF00500
ZEF00510
ZEF00520
ZEF00530
ZEF00540
ZEF00550
ZEF00560
ZEF00570
ZEF00580
ZEF00590
ZEF00600
ZEF00610

```

NUMBER LEV NT

```

030      1 0 0      /* MISC DECLARATIONS */
040      1 1 0      UCL (I,J,K,L) FIXED BIN(15);
050      1 1 0      DCL RSND FIXED BIN(15) STATIC INIT(1);
060      1 1 0      UCL (M,N,P) FIXED DEC(3,0);
070      1 1 0      UCL (INTERM,TERM) CHAR(100) VAR;
080      1 1 0      UCL (ADUR,SUBSTR) BUILTIN;
090      1 1 0      DCL HIGH STATE FIXED BIN(15) STATIC INIT(0);
100      1 1 0      DCL NEW STATE FIXED BIN(15);
110      1 1 0      DCL CHAR CHAR(1);
120      1 1 0      DCL CHAR2 CHAR(2);
130      1 1 0      DCL D_CHAR2 FIXED DEC(3,0) BASED(P_CHAR2);
140      1 1 0      DCL P_CHAR2 PTR;
150      1 1 0      DCL P_CHAR2 = ADUR(CHAR2);
          DCL MORE BIT(1) STATIC INIT(.1'b);
ZEF00620
ZEF00630
ZEF00640
ZEF00650
ZEF00660
ZEF00670
ZEF00680
ZEF00690
ZEF00700
ZEF00710
ZEF00720
ZEF00730
ZEF00740
ZEF00750

```

NUMBER LEV NT

770	1	0	/* ON UNITS AND INITIALIZATION */	ZEF00760
780	1	0	ON SUBRG SIGNAL ERROR;	ZEF00770
			ON STKG SIGNAL ERROR;	ZEF00780
790	1	0	ON ERROR SNAP	ZEF00790
			BEGIN;	ZEF00800
810	2	0	CLOSE FILE (SYSPRINT);	ZEF00810
830	2	0	END;	ZEF00820
			END;	ZEF00830
840	1	0	ON ENDFILE(IN) MORE = '0'B;	ZEF00840
850	1	0	UPEN FILE (SYSPRINT);	ZEF00850
			FILE (IN);	ZEF00860

NUMBER LEV NT

```

      880 1 0      /* MAIN PROCESSING */
              CALL GRAMMAK: /* SET UP SYMBOL & PRODUCTION TABLES */
              /*
              PRINT GRAMMAK (START SYMBOL,NONTERMINALS,TERMINALS,
              AND PRODUCTIONS */
      910 1 0      DO I = 1 TO NO_SYM WHILE (SYMBOLS(I).SYM ^= ' ');
      920 1 1      IF SYMBOLS(I).TYPE
      940 1 1      THEN NONTERM = NONTERM || SYMBOLS(I).SYM || ' ';
      950 1 1      ELSE TERM = TERM || SYMBOLS(I).SYM || ' ';
              END;
      960 1 0      PUT EDIT ('INTERNAL GRAMMAR:',NONTERMINALS: ',NONTERM,
              TERMINALS: ',TERM,START SYMBOL: ',
              SYMBOLS(1).SYM,PRODUCTIONS: ')
              ( SKIP(2),CUL(1),A,
              SKIP,CUL(4),A,SKIP(0),CUL(18),A,
              SKIP,CUL(4),A,SKIP(0),CUL(12),A,
              SKIP,CUL(4),A,SKIP(0),CUL(18),A,
              SKIP,CUL(4),A);
      1040 1 0      DO I = 1 TO NO_PROD WHILE (PROD(I).NO ^= 0);
      1050 1 1      N = PROD(I).NO;
      1060 1 1      PUT EDIT ('( ',N,') ',PROD(I).LP,' -> ',PROD(I).RP )
      1080 1 1      ( COL(8),A,P(29),4 A);
              END;
      1090 1 0      CALL COMP_FIRSTs: /* COMPUTE FIRSTS FOR EACH NONTERMINAL */
      1100 1 0      PUT EDIT ('GENERATED ITEMS') (A) PAGE;
      1110 1 0      PUT EDIT ('NO ','STATE','ITEMS')
              (SKIP(2),COL(1),A,CUL(5),A,COL(15),A);
      1130 1 0      CALL LR_ITEMS: /* GENERATE LR(1) ITEMS */
      1140 1 0      CALL SK_TAB: /* GENERATE SHIFT-REDUCE TABLE */

```

ZEF00870

ZEF00880

ZEF00890

ZEF00900

ZEF00910

ZEF00920

ZEF00930

ZEF00940

ZEF00950

ZEF00960

ZEF00970

ZEF00980

ZEF00990

ZEF01000

ZEF01010

ZEF01020

ZEF01030

ZEF01040

ZEF01050

ZEF01060

ZEF01070

ZEF01080

ZEF01090

ZEF01100

ZEF01110

ZEF01120

ZEF01130

ZEF01140

PL/I OPTIMIZING COMPILER /* PARSER GENERATOR */

NUMBER LEV NT

```

1160      0      /* OUTPUT GUTUS ARRAY */
1170      1 0 PUT EDIT ('LR(1) GOTO FUNCTION:') (A) PAGE;
1180      1 0 PUT EDIT (' ') (SKIP,A);
1190      1 0 UO J = 0 TO NEWSTATE;
1200      1 1 M = A;
1210      1 1 PUT EDIT (' ',M) (A,P*29);
1220      1 1 END;
1230      1 0 DO I = 1 BY 1 UNTIL(SYMBOLS(I).SYM = ' ');
1240      1 1 PUT EDIT (SYMBOLS(I).SYM) (SKIP,A);
1250      1 1 UO J = 0 TO 9;
1260      1 1 IF GOTOS(I,J) = -1
1270      1 2 THEN M = 0;
1280      1 2 ELSE M = GOTOS(I,J);
1290      1 2 PUT EDIT (' ',M) (A,P*22);
1300      1 1 END;

1320      /* OUTPUT ACTION FUNCTION */
1330      1 0 PUT EDIT ('LR(1) ACTION FUNCTION:') (SKIP(2),A);
1340      1 0 PUT EDIT (' ') (SKIP,A);
1350      1 1 DO J = 0 TO NEWSTATE;
1360      1 1 M = J;
1370      1 1 PUT EDIT (' ',M) (A,P*29);
1380      1 1 DO I = 1 BY 1 UNTIL(SYMBOLS(I).SYM = ' ');
1390      1 1 IF SYMBOLS(I).SYM = 'S';
1400      1 1 THEN CHAR = 'S';
1410      1 1 ELSE CHAR = SYMBOLS(I).SYM;
1420      1 1 PUT EDIT (CHAR) (CUL(1),A);
1430      1 1 DO J = 0 TO NEWSTATE;
1440      1 1 IF CHAR2 = ACTION(I,J); /* ASSUME REDUCE */
1450      1 2 IF ACTION(I,J) = -1; /* ERROR CONDITION */
1460      1 2 THEN CHAR2 = ' ';
1470      1 2 IF ACTION(I,J) = 0; /* ACCEPT */
1480      1 2 THEN CHAR2 = 'A';
1490      1 2 IF ACTION(I,J) = -2; /* SHIFT */
1500      1 2 THEN CHAR2 = 'S';
1510      1 2 IF ACTION(I,J) > 0;
1520      1 2 THEN PUT EDIT (' ',CHAR2) (A,P*25);
1530      1 2 ELSE PUT EDIT (' ',CHAR2) (2 A);
1540      1 2 END;
1550      1 1 END;

1560      CALL LALR; /* CREATE LALR TABLES */

```

ZEF01150
 ZEF01160
 ZEF01170
 ZEF01180
 ZEF01190
 ZEF01200
 ZEF01210
 ZEF01220
 ZEF01230
 ZEF01240
 ZEF01250
 ZEF01260
 ZEF01270
 ZEF01280
 ZEF01290
 ZEF01300
 ZEF01310
 ZEF01320
 ZEF01330
 ZEF01340
 ZEF01350
 ZEF01360
 ZEF01370
 ZEF01380
 ZEF01390
 ZEF01400
 ZEF01410
 ZEF01420
 ZEF01430
 ZEF01440
 ZEF01450
 ZEF01460
 ZEF01470
 ZEF01480
 ZEF01490
 ZEF01500
 ZEF01510
 ZEF01520
 ZEF01530
 ZEF01540
 ZEF01550
 ZEF01560

PL/I OPTIMIZING COMPILER

/* PARSEK GENERATOR */

NUMBER LEV NT

```

1580      0      /* OUTPUT LALR GOTOS ARRAY */
1590      1 0 PUT EDIT ('LALR(1) GOTO FUNCTION: ') (SKIP(2),A);
1600      1 0 PUT EDIT (' ') (SKIP,A);
1610      1 0 DO J = 0 TO LALR_SZ-1;
1620      1 1 M = J;
1630      1 1 PUT EDIT (' ',M) (A,P,Z9);
1640      1 0 DO I = 1 BY 1 WHILE (SYMBOLS(I).SYM ^= ' ');
1650      1 1 PUT EDIT (SYMBOLS(I).SYM) (SKIP,A);
1660      1 1 DO J = 0 TO LALR_SZ-1;
1670      1 2 IF LALR_GOTO(I,J) = -1
1680      1 2 THEN M = 0;
1690      1 2 ELSE M = LALR_GOTO(I,J);
1700      1 2 PUT EDIT (' ',M) (A,P,ZZ);
1710      1 2 END;
1720      1 1 END;

1740      0      /* OUTPUT LALR ACTION FUNCTION */
1750      1 0 PUT EDIT ('LALR(1) ACTION FUNCTION: ') (SKIP(2),A);
1760      1 0 PUT EDIT (' ') (SKIP,A);
1770      1 0 DO J = 0 TO LALR_SZ-1;
1780      1 1 M = J;
1790      1 1 PUT EDIT (' ',M) (A,P,Z9);
1800      1 1 DO I = 1 BY 1 UNTIL (SYMBOLS(I).SYM = ' ');
1810      1 1 IF SYMBOLS(I).SYM = ' ';
1820      1 1 THEN CHAR = 'S';
1830      1 1 ELSE CHAR = SYMBOLS(I).SYM;
1840      1 1 PUT EDIT (CHAR) (CUL(1),A);
1850      1 1 DO J = 0 TO LALR_SZ-1;
1860      1 2 U-CHAR2 = LALR_ACT(I,J); /* ASSUME REDUCE */
1870      1 2 IF LALR_ACT(1,J) = -1 /* ERROR CONDITION */
1880      1 2 THEN CHAR2 = ' ';
1890      1 2 IF LALR_ACT(1,J) = 0 /* ACCEPT */
1900      1 2 THEN CHAR2 = 'A';
1910      1 2 IF LALR_ACT(1,J) = -2 /* SHIFT */
1920      1 2 THEN CHAR2 = 'S';
1930      1 2 IF LALR_ACT(1,J) > 0
1940      1 2 THEN PUT EDIT (' ',CHAR2) (A,P,Z9);
1950      1 2 ELSE PUT EDIT (' ',CHAR2) (2 A);
1960      1 2 END;
1970      1 1 END;

```

ZEF01570
 ZEF01580
 ZEF01590
 ZEF01600
 ZEF01610
 ZEF01620
 ZEF01630
 ZEF01640
 ZEF01650
 ZEF01660
 ZEF01670
 ZEF01680
 ZEF01690
 ZEF01700
 ZEF01710
 ZEF01720

 ZEF01730
 ZEF01740
 ZEF01750
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 ZEF01770
 ZEF01780
 ZEF01790
 ZEF01800
 ZEF01810
 ZEF01820
 ZEF01830
 ZEF01840
 ZEF01850
 ZEF01860
 ZEF01870
 ZEF01880
 ZEF01890
 ZEF01900
 ZEF01910
 ZEF01920
 ZEF01930
 ZEF01940
 ZEF01950
 ZEF01960
 ZEF01970

NUMBER LEV NT

```

2010 1 0
/*-----*
/* OUTPUT KEY TABLES
/*-----*/
OPEN FILE(SYMTAB),
FILE(PRODTAB),
FILE(LALRTAB),
FILE(LALRSZ);

2050 1 0 WRITE FILE(SYMTAB) FROM(SYMBOLS);
2060 1 0 WRITE FILE(PRODTAB) FROM(PROD);
2070 1 0 WRITE FILE(LALRSZ) FROM(LALR_SZ);

2080 1 0 TBL(2) FAXED IN(15);
2090 1 0 DCL C_TBL CHAR(4) BASED (P_TBL);
2100 1 0 DCL P_TBL PTR;
2110 1 0 P_TBL = ADDR (TBL(1));

2120 1 0 DO I = 1 TO LALR_SZ-M;
2130 1 1 DO J = 0 TO LALR_SZ-N;
2140 1 2 TBL(1) = LALR_GUTU (I,J);
2150 1 2 TBL(2) = LALR_ACT (I,J);
2160 1 2 WRITE FILE (LALRTAB) FROM (C_TBL);
2170 1 2 END;
2180 1 1 END;

2190 1 0 CLOSE FILE(SYMTAB),
FILE(PRODTAB),
FILE(LALRTAB),
FILE(LALRSZ);

```

ZEF01960
ZEF01990
ZEF02000
ZEF02010
ZEF02020
ZEF02030
ZEF02040

ZEF02050
ZEF02060
ZEF02070

ZEF02080
ZEF02090
ZEF02100
ZEF02110

ZEF02120
ZEF02130
ZEF02140
ZEF02150
ZEF02160
ZEF02170
ZEF02180

ZEF02190
ZEF02200
ZEF02210
ZEF02220

PL/1 OPTIMIZING COMPILER

/* PARSER GENERATOR */

NUMBER LEV NT

```

2230 1 0 LR_ITEMS: PROC;
2240 2 0 DCL(I,J,K) FIXED BIN(15);
2250 2 0 DCL(P,N) FIXED DEC(3,0);

      /* SET UP FIRST ITEM */
2270 2 0 NEXT_I = 1;
2280 2 0 CALL ADD_ITEM(0,1,0,'$');
2290 2 0 ITEMS SET(I) = 1;
2300 2 0 N_ITEMS = 1;

2310 2 0 CALL CLOSURE;

2320 2 0 NEWSTATE = 0;
2330 2 0 STATE = 0;
2340 2 0 AGAIN: DO I = 1 TO NO SYN WHILE(SYMBOLS(I).SYN ^= ' ');
2350 2 1 NEWSTATE = NEWSTATE + 1;
2360 2 1 STATE_ITEMS(NEWSTATE) = 0;
2370 2 1 CALL GOTU(STATE,SYMBOLS(I).SYM,1);
2380 2 1 ENDO;
2390 2 1 STATE = STATE + 1;
2400 2 1 DO I = 1 TO NEXT_I-1;
2410 2 1 IF ITEMS(I).STATE = STATE
      THEN GOTU AGAIN;
2430 2 1 ENDO;

      /* PRINT ITEM TABLE */
2450 2 0 DO I = 1 TO NEXT_I-1;
2460 2 1 P = 1;
2470 2 1 N = ITEMS(I).STATE;
2480 2 1 J = ITEMS(I).PROD;
2490 2 1 K = ITEMS(I).POSIT;
2500 2 1 IF K <= 0 THEN NEW_RP = ' ' || PROD(J).RP;
2510 2 1 IF K >= 0 THEN NEW_RP = SUBSTR(PROD(J).RP,1,K) || PROD(J).NO_RP || ' ';
      ELSE NEW_RP = SUBSTR(NEW_RP,1,NO_RP) || SUBSTR(PROD(J).RP,1,K) || PROD(J).NO_RP || ' ';
2530 2 1 NEW_RP = SUBSTR(NEW_RP,1,NO_RP) || SUBSTR(PROD(J).RP,1,K) || PROD(J).NO_RP || ' ';
2550 2 1 PUT EDIT (P,N,PROD(J).LA,
2570 2 1 (P,N,PROD(J).LP,
2590 2 1 (COL(1),P,ZZ9,COL(6),P,ZZ9,COL(15),3 A);
      ENDO;

2600 2 0 ENDO; /* LR_ITEMS */

```

ZEF02230
 ZEF02240
 ZEF02250
 ZEF02260
 ZEF02270
 ZEF02280
 ZEF02290
 ZEF02300
 ZEF02310
 ZEF02320
 ZEF02330
 ZEF02340
 ZEF02350
 ZEF02360
 ZEF02370
 ZEF02380
 ZEF02390
 ZEF02400
 ZEF02410
 ZEF02420
 ZEF02430
 ZEF02440
 ZEF02450
 ZEF02460
 ZEF02470
 ZEF02480
 ZEF02490
 ZEF02500
 ZEF02510
 ZEF02520
 ZEF02530
 ZEF02540
 ZEF02550
 ZEF02560
 ZEF02570
 ZEF02580
 ZEF02590
 ZEF02600

PL/1 OPTIMIZING COMPILER /* PARSER GENERATOR */

NUMBER LEV NT

```

2610 1 0 CLOSURE: PROC; /* SET OF ITEMS TO BE CLOSED IS DEFINED BY ARRAY
      ITEMS_SET AND THE NUMBER OF ITEMS IN THE SET
      IS CONTAINED IN N_ITST */
2650 2 0 DCL (I,J,K,L,M) FIXED BIN(15);
2660 2 0 DO I = 1 BY 1 UNTIL (I >= N_ITST); /* UNTIL NO MORE ITEMS CAN
      BE ADDED */
2680 2 1 M = ITEMS_SET(I); /* POINTS TO ITEM IN ITEM TABLE */
2690 2 1 CALL END_PROD(M);
2700 2 1 IF N_PROD = 0
      THEN GOTU CLOEND; /* NEXT SYMBOL IN ITEM MUST BE
      NONTERMINAL WITH PRODUCTIONS */
2730 2 1 CALL MFIRST(M);
2740 2 1 DO J = 1 TO N_PROD;
2750 2 2 DO K = 1 TO N_FRST;
2760 2 3 CALL ADD_ITEM (ITEMS(M).STATE,PROD_#(J),0,FRST(K));
2770 2 3 END; /* END K LOOP */
2780 2 2 END; /* END J LOOP */
2790 2 1 CLOEND: END; /* END I LOOP */
2810 2 0 END; /* CLOSURE */
      ZEF02610
      ZEF02620
      ZEF02630
      ZEF02640
      ZEF02650
      ZEF02660
      ZEF02670
      ZEF02680
      ZEF02690
      ZEF02700
      ZEF02710
      ZEF02720
      ZEF02730
      ZEF02740
      ZEF02750
      ZEF02760
      ZEF02770
      ZEF02780
      ZEF02790
      ZEF02800
      ZEF02810

```

PL/I OPTIMIZING COMPILER /* PARSER GENERATOR */

NUMBER LEV NT

```

2820 1 0 FND_PROD: PROC(ITEM);
/* POINTS TO ITEM IN ITEM ARRAY. */
/* RETURNS LIST OF PRODUCTIONS. */
2850 2 0 DCL ITEM FIXED BIN(15);
2860 2 0 DCL I FIXED BIN(15);
2870 2 0 DCL SYMBOL CHAR(1);
2880 2 0 N_PROD = 0;
2890 2 0 SYMBOL = SUBSTR(PROD(ITEM),PROD).RP,ITEMS(ITEM).POSIT+1,1);
2900 2 0 IF (ITEM(SYMBOL)) (SYMBOL = 1);
2920 2 0 THEN RETURN; /* SYMBOL EITHER TERMINAL OR BLANK */
2930 2 1 DO I = 1 BY 1 WHILE (PROD(I).NO = 0);
IF SYMBOL = PROD(I).LP
THEN DO;
N_PROD = N_PROD + 1;
PROD_#(N_PROD) = PROD(I).NO;
END;
END;
2950 2 2 N_PROD = N_PROD + 1;
2960 2 2 PROD_#(N_PROD) = PROD(I).NO;
2970 2 2 END;
2980 2 1
2990 2 0 ENU; /* FND_PROD */
ZEF02820
ZEF02830
ZEF02840
ZEF02850
ZEF02860
ZEF02870
ZEF02880
ZEF02890
ZEF02900
ZEF02910
ZEF02920
ZEF02930
ZEF02940
ZEF02950
ZEF02960
ZEF02970
ZEF02980
ZEF02990

```

/* PARSER GENERATOR */

PL/I OPTIMIZING COMPILER

NUMBER LEV NT

```

3000 1 0 MFIRST: PROC (ITEM);
3010 2 0 DCL ITEM FIXED BIN(15);
3020 2 0 N_FIRST = 0;
3030 2 0 SYMBOL = SUBSTR(PKUD(ITEMS(ITEM).PKUD).KP,ITEMS(ITEM).POSIT+2,1);ZEF03030
                                           ZEF03000
3040 2 0 IF SYMBOL = ' '
      THEN;
3050 2 0 ELSE CALL FIRST(SYMBOL);
3060 2 0 IF N_FIRST = 0
      THEN DO;
3070 2 0   N_FIRST = 1;
3080 2 1   FKST(1) = ITEMS(ITEM).LA;
3090 2 1   END;
3100 2 1   END;
3110 2 1   END;
3120 2 0 END; /* MFIRST */

3130 1 0 FIRST: PROC (SYM);
3140 2 0 DCL SYM CHAR(1);
3150 2 0 DCL (1,J) FIXED BIN(15);
3160 2 0 J = FIRST_SYM(SYM);
3170 2 0 N_FIRST = FIRSTS(J).NO_FIRSTS;
3180 2 0 DO I = 1 TO FIRSTS(J).MULTFIRSTS;
3190 2 1   FIRST(I) = FIRSTS(J).FIRSTS(I);
3200 2 1   END;
3210 2 0 END; /* FIRST */
3220 2 0

```

ZEF03000
 ZEF03010
 ZEF03020
 ZEF03030
 ZEF03040
 ZEF03050
 ZEF03060
 ZEF03070
 ZEF03080
 ZEF03090
 ZEF03100
 ZEF03110
 ZEF03120
 ZEF03130
 ZEF03140
 ZEF03150
 ZEF03160
 ZEF03170
 ZEF03180
 ZEF03190
 ZEF03200
 ZEF03210

NUMBER LEV NT

```

3220 1 0 PTR_SYM: PROC(SYM);
3230 2 0 UCL SYM_CHAR(1);
3240 2 0 UCL 1 FIXED BIN(15);
3250 2 0 DO 1 = 1 TO NO_SYM UNTIL(SYMBOLS(1).SYM = SYM); END;
3260 2 0 RETURN(1);
3270 2 0 END; /* PTR_SYM */

3280 1 0 ADD_TRM: PROC(SYM,1);
/*-----
ADD SYM TO FIRSTS OF SYMBOL POINTED TO BY 1
IF SYMBOL IS NOT ALREADY THERE.
-----*/
3330 2 0 UCL SYM_CHAR(1);
3340 2 0 DCL (I,J,K) FIXED BIN(15);
3350 2 0 IF K = FIRSTS(1).NO_FIRSTS;
3360 2 0 IF K > 0 THEN
3370 2 1 DO 1 = 1 TO FIRSTS(1).NO_FIRSTS;
3380 2 1 IF SYM = FIRSTS(1).FIRSTS(J)
3390 2 1 THEN RETURN;
3400 2 1 END;
3410 2 0 FIRSTS(1).NO_FIRSTS = K + 1;
3420 2 0 FIRSTS(1).FIRSTS(K+1) = SYM;
3430 2 0 TOT_FIRSTS = TOT_FIRSTS + 1;
3440 2 0 END; /* ADD_TRM */

3450 1 0 TRM: PROC (SYM) RETURNS ( BIT );
3460 2 0 UCL SYM_CHAR(1);
3470 2 0 UCL 1 FIXED BIN(15);
3480 2 0 DO 1 = 1 TO NO_SYM UNTIL (SYM = SYMBOLS(1).SYM);
3490 2 1 END;
3500 2 0 RETURN (SYMBOLS(1).TYPE);
3510 2 0 END; /* TRM */

```

ZEF03220

ZEF03230

ZEF03240

ZEF03250

ZEF03260

ZEF03270

ZEF03280

ZEF03290

ZEF03300

ZEF03310

ZEF03320

ZEF03330

ZEF03340

ZEF03350

ZEF03360

ZEF03370

ZEF03380

ZEF03390

ZEF03400

ZEF03410

ZEF03420

ZEF03430

ZEF03440

ZEF03450

ZEF03460

ZEF03470

ZEF03480

ZEF03490

ZEF03500

ZEF03510

NUMBER LEV NT

```

3520 1 0 COMP_FIRSTS: PROC;
3530 2 0 DCL (I,J,K,L,M,N,P) FIXED BIN(15);
3540 2 0 DCL SYM CHAR(1);
3550 2 0 DO J = 1 TO NO_PROD UNTIL(PROD(J).NO = 0); END;
3560 2 0 IF J > 1 THEN N = J - 1; ELSE SIGNAL ERROR;
3570 2 0 DO I = 1 TO NO_SYM UNTIL(SYMBOLS(I).SYM = ' '); END;
3580 2 0 IF I > 1 THEN M = I - 1; ELSE SIGNAL ERROR;
3590 2 0 C_F1: P = TOT_FIRSTS;
3600 2 0 DO I = M TO 1 BY -1; /* GO THRU SYMBOLS */
3610 2 0 IF TRM(SYMBOLS(I).SYM) /* IF NONTERMINAL */
3620 2 1 THEN
3630 2 2 DO J = N TO 1 BY -1; /* GO THRU PRODUCTIONS */
3640 2 2 IF SYMBOLS(I).SYM = PROD(J).LP
3650 2 2 THEN
3660 2 3 DO;
3670 2 3 SYM = SUBSTR(PROD(J).RP,1,1);
3680 2 3 IF TRM(SYM) /* IF NONTERMINAL */
3690 2 3 THEN
3700 2 3 DO;
3710 2 3 K = PTR_SYM(SYM);
3720 2 4 IF FIRSTS(K).NO_FIRSTS = 0
3730 2 4 THEN
3740 2 4 ELSE DO L = 1 TO FIRSTS(K).NO_FIRSTS;
3750 2 4 CALL ADD_TRM
3760 2 4 (FIRSTS(K).FIRSTS(L),I);
3770 2 4 END;
3780 2 4 END;
3790 2 4 END; ELSE CALL ADD_TRM(SYM,I); /* ADD TERMINAL */
3800 2 3 ENDO;
3810 2 3 ENDO;
3820 2 3 ENDO; ELSE CALL ADD_TRM(SYMBOLS(I).SYM,I); /* TERMINAL */
3830 2 2 ENDO;
3840 2 2 IF TOT_FIRSTS = P THEN ; ELSE GOTO C_F1;
3850 2 0
/*-----*/
/* PRINT COMPUTED FIRSTS. */
/*-----*/
3890 2 0 PUT EDIT ('COMPUTED FIRSTS: ', ' ') (A,SKIP,A) SKIP(2);
3900 2 0 DO I = 1 BY 1 WHILE(SYMBOLS(I).SYM = ' ');
3910 2 1 PUT EDIT (SYMBOLS(I).SYM, ' ') (2 A);
3920 2 1 DO J = 1 TO FIRSTS(I).NO_FIRSTS;
3930 2 2 PUT EDIT (FIRSTS(I).FIRSTS(J)) (A);
3940 2 2 ENDO;
3950 2 2 PUT EDIT (' ') (A) SKIP;
3960 2 1 ENDO;
3970 2 0 ENDO; /* COMP_FIRSTS */

```

ZEF03520
 ZEF03530
 ZEF03540
 ZEF03550
 ZEF03560
 ZEF03570
 ZEF03580
 ZEF03590
 ZEF03600
 ZEF03610
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 ZEF03630
 ZEF03640
 ZEF03650
 ZEF03660
 ZEF03670
 ZEF03680
 ZEF03690
 ZEF03700
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 ZEF03820
 ZEF03830
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 ZEF03870
 ZEF03880
 ZEF03890
 ZEF03900
 ZEF03910
 ZEF03920
 ZEF03930
 ZEF03940
 ZEF03950
 ZEF03960
 ZEF03970

NUMBER LEV NT

```

3980 1 0 ADD_ITEM: PROC (STATE,PRUD,POSIF,LA);
3990 2 0 DCL (STATE,PRUD,PUSIT,L,J,K) FIXED BIN(15);
4000 2 0 DCL (P,N) FIXED DEC(3,0);
4010 2 0 DCL LA CHAR(1);
      /* ADD ITEMS TO ITEM TABLE */
4030 2 0 DO L = 1 TO NEXT_I-1 UNTIL ((ITEMS(L).STATE = STATE)
      /* FIND LAST ENTRY FOR STATE */
4060 2 1 IF (ITEMS(L).STATE = STATE) &
      (ITEMS(L).PRUD = PRUD) &
      (ITEMS(L).POSIT = POSIT) &
      (ITEMS(L).LA = LA)
      THEN RETURN; /* ITEM ALREADY EXISTS */
4110 2 1 END;
4120 2 0 IF L > NEXT_I-1
      THEN
      ELSE ITEMS(L).NEXT = NEXT_I;
4140 2 0 ITEMS(NEXT_I).STATE = STATE;
4150 2 0 ITEMS(NEXT_I).NEXT = -1;
4160 2 0 ITEMS(NEXT_I).PRUD = PRUD;
4170 2 0 ITEMS(NEXT_I).POSIT = POSIT;
4180 2 0 ITEMS(NEXT_I).LA = LA;
4190 2 0 N_LIST = N_LIST + 1;
4200 2 0 ITEMS_SET(N_LIST) = NEXT_I;
4210 2 0 NEXT_I = NEXT_I + 1;
4220 2 0 STATE_ITEMS(STATE) = STATE_ITEMS(STATE) + 1;
4230 2 0 IF STATE > HIGH_STATE
4240 2 0 THEN HIGH_STATE = STATE; /* SAVE HIGHEST STATE NO. WITH
      VALID ITEMS IN IT */
4270 2 0 END; /* ADD_ITEM */
ZEF03960
ZEF03990
ZEF04000
ZEF04010
ZEF04020
ZEF04030
ZEF04040
ZEF04050
ZEF04060
ZEF04070
ZEF04080
ZEF04090
ZEF04100
ZEF04110
ZEF04120
ZEF04130
ZEF04140
ZEF04150
ZEF04160
ZEF04170
ZEF04180
ZEF04190
ZEF04200
ZEF04210
ZEF04220
ZEF04230
ZEF04240
ZEF04250
ZEF04260
ZEF04270

```

PL/1 OPTIMIZING COMPILER

/* PARSER GENERATOR */

NUMBER LEV NT

```

4280 1 0 GOTO: PKOC (STATE,X,SYM);
      /* STATE = CURRENT STATE TO PERFORM GOTO ON.
      X = SYMBOL TO LOOK FOR.
      SYM = POSITION OF X IN SYMBOL TABLE. */
4320 2 0 DCL (STATE,SYM,I,J,K,L,S,N1) FIXED BIN(15);
4330 2 0 DCL EQUAL BIT(1);
4340 2 0 DCL X CHAR(1);

4350 2 0 DO I = 1 TO NEXT_I-1 UNTIL (ITEMS(I).STATC = STATE);
4370 2 1 END;
4380 2 0 S_N1 = NEXT_I; /* SAVE NEXT_I VALUE */
4390 2 0 N_ITST = 0;

4400 2 0 DO WHILE (I <= -1); /* DO FOR ALL ITEMS IN STATE */
4410 2 1 J = ITEMS(I).PRUO;
4420 2 1 K = ITEMS(I).POSIT + 1;
4430 2 1 IF SUBSTR (PRUO(J),RP,K,1) = X
      THEN CALL ADD_ITEM (NEWSTATE,J,K,ITEMS(I).LA);
      I = ITEMS(I).NEXT;
4450 2 1 END;
4460 2 1 IF N_ITST = 0
      THEN DO;
4470 2 0 NEWSTATE = HIGH_STATE; /* NEEDED BECAUSE NEWSTATE
4480 2 1 RETURN; /* BLINDLY INCREMENTED BY
4490 2 1 END;
4500 2 1
4510 2 1
4520 2 0 CALL CLOSURE;

```

```

ZEF04280
ZEF04290
ZEF04300
ZEF04310
ZEF04320
ZEF04330
ZEF04340
ZEF04350
ZEF04360
ZEF04370
ZEF04380
ZEF04390
ZEF04400
ZEF04410
ZEF04420
ZEF04430
ZEF04440
ZEF04450
ZEF04460
ZEF04470
ZEF04480
ZEF04490
ZEF04500
ZEF04510
ZEF04520

```

PL/I OPTIMIZING COMPILER

/* PARSER GENERATOR */

NUMBER LEV NT

```

4540      2 0      /* CHECK FOR DUPLICATE GOTO STATE */
4550      2 0      EQUAL = '0'B;
4560      2 1      DO I = 1 TO NEWSTATE-1;
                     IF STATE_ITEMS(I) = STATE_ITEMS(NEWSTATE)
                     THEN
                     DO;
                         /* CHECK FOR STATE IDENTITY */
                         DO J = 1 TO NEXT I-1 UNTIL (ITEMS(J).STATE = I); END;
                         DO WHILE (J = -1);
                             K = S_NI;
                             DO WHILE (K = -1);
                                 IF (ITEMS(J).PROD = ITEMS(K).PROD) &
                                    (ITEMS(J).POSIT = ITEMS(K).POSIT) &
                                    (ITEMS(J).LA = ITEMS(K).LA)
                                 THEN GOTO GOTO2; /* ITEMS = GET NEXT ONE */
                                 K = ITEMS(K).NEXT;
                             END;
                         END;
                         GOTO GOTO3; /* STATES (I & NEWSTATE) = */
                     END;
                     J = ITEMS(J).NEXT;
                     ENDO;
                     EQUAL = '1'B; /* STATES (I & NEWSTATE) ARE EQUAL */
                     GOTO GOTO4;
                     ENDO;
4740      2 1      GOTO3: ENDO;
4750      2 1      GOTO4:
4770      2 0      GOTO4: IF EQUAL
                     THEN DO;
                         GOTOS(SYM.STATE) = I; /* BACKOUT NEWSTATE ITEMS */
                         NEXT I = S_NI;
                         NEWSTATE = NEWSTATE - 1;
                         HIGH_STATE = NEWSTATE;
                     END;
                     ELSE GOTOS(SYM.STATE) = NEWSTATE;
4840      2 0      ENDO; /* GOTO */
4850      2 0      ENDO;
4860      2 0      ENDO;

```

ZEF04530
 ZEF04540
 ZEF04550
 ZEF04560
 ZEF04570
 ZEF04580
 ZEF04590
 ZEF04600
 ZEF04610
 ZEF04620
 ZEF04630
 ZEF04640
 ZEF04650
 ZEF04660
 ZEF04670
 ZEF04680
 ZEF04690
 ZEF04700
 ZEF04710
 ZEF04720
 ZEF04730
 ZEF04740
 ZEF04750
 ZEF04760
 ZEF04770
 ZEF04780
 ZEF04790
 ZEF04800
 ZEF04810
 ZEF04820
 ZEF04830
 ZEF04840
 ZEF04850
 ZEF04860

NUMBER LEV NT

```

4870 1 0 SK_TAB: PRUD;
4880 2 0 DCL CHAR CHAR(1);
4890 2 0 DCL (I,J,VALUE) FIXED BIN(15);
4900 2 0 DCL (M,N) FIXED BIN(3,0);
4910 2 0 DO I = 1 TO NEXT I-1;
4920 2 1 IF (ITEMS(I).PRUD = 1) &
      (ITEMS(I).POSIT = PRUD(ITEMS(I).PRUD).NO_RP) &
      (ITEMS(I).LA = '$') /* ACCEPT */
      THEN DO;
          CHAR = '$';
          VALUE = 0;
          GOTO ACTION_FILL;
        END;
5000 2 1 IF ITEMS(I).POSIT = PRUD(ITEMS(I).PRUD).NO_RP
      THEN DO;
          CHAR = ITEMS(I).LA;
          VALUE = ITEMS(I).PRUD; /* REDUCE PRUD. NO. */
          IF ITEMS(I).LA = '$'
          THEN GOTO ACTION_FILL;
          IF ~TRM(ITEMS(I).LA) /* MUST BE TERMINAL */
          THEN GOTO ACTION_FILL;
        END;
5090 2 1 CHAR = SUBSTR(PRUD(ITEMS(I).PRUD).RP,ITEMS(I).POSIT+1,1);
5100 2 1 IF ~TRM(CHAR) /* IF CHAR IS A TERMINAL */
      THEN DO;
          VALUE = -2; /* SHIFT */
          GOTO ACTION_FILL;
        END;
5120 2 2
5130 2 2
5140 2 2 GOTO END_DO;
5150 2 1 GOTO END_DO;
5160 2 1 ACTION_FILL:
5180 2 2 DO J = 1 BY 1 UNTIL ((SYMBOLS(J).SYM = CHAR) |
      (SYMBOLS(J).SYM = '$'));
5200 2 1 /* IF SYM = '$' THEN CHAR MUST = '$' */
      (ACTION(J,ITEMS(I).STATE) = -1) |
      (ACTION(J,ITEMS(I).STATE) = VALUE)
      THEN ACTION(J,ITEMS(I).STATE) = VALUE;
5230 2 1 ELSE DO;
5240 2 2 M = ITEMS(I).STATE;
5250 2 2 N = I;
5260 2 2 PUT EDIT ('ACTION CONFLICT: STATE = *M,
      *CHARACTER = *CHAR, ITEM = *N)
      (SKIP,A,P,Z9,3 A,P,Z9);
5290 2 2 END;
5300 2 1 END_DO;
5320 2 0 END; /* SK_TAB */

```

```

ZEF04870
ZEF04880
ZEF04890
ZEF04900
ZEF04910
ZEF04920
ZEF04930
ZEF04940
ZEF04950
ZEF04960
ZEF04970
ZEF04980
ZEF04990
ZEF05000
ZEF05010
ZEF05020
ZEF05030
ZEF05040
ZEF05050
ZEF05060
ZEF05070
ZEF05080
ZEF05090
ZEF05100
ZEF05110
ZEF05120
ZEF05130
ZEF05140
ZEF05150
ZEF05160
ZEF05170
ZEF05180
ZEF05190
ZEF05200
ZEF05210
ZEF05220
ZEF05230
ZEF05240
ZEF05250
ZEF05260
ZEF05270
ZEF05280
ZEF05290
ZEF05300
ZEF05310
ZEF05320

```

PL/I OPTIMIZING COMPILER /* PARSER GENERATOR */

NUMBER LEV NT

```

5330 1 0 GRAMMAR: PROC;
/*-----
GRAMMAR DEFINITION:
VALID IC'S ARE:
NONT FOR NUNTERMINALS SEPARATED BY BLANKS
FIRST NUNTERMINAL IS START SYMBOL
TERMINALS SEPARATED BY BLANKS
TERM FOR PRODUCTIONS - 1 PER RECORD
PRD FOR <LEFT PART> : <RIGHT PART>
-----*/

5430 UCL CARU CHAR(60);
5440 UCL A_CARD(80) CHAR(1) BASED(P_CARD);
5450 DCL P_CARD PTR;
5460 P_CARD = ADDR(CARD);
5470 UCL (1,J,K,L,M) FIXED BIN(15);
5480 UCL REP_SYM CHAR(40) STATIC INIT
      ('ABCOEFGHIJKLMNOPQRSTUVWXYZ0123456789')(*C);
5500 UCL REPSYM(40) CHAR(1) BASED (P_RS);
5510 UCL P_RS PTR;
5520 P_RS = ADDR (REP_SYM);

5530 READ FILE (IN) INTO (CARD);
5540 I = 0;
5550 J = 0;
5560 DO WHILE (MORE);
5570 PUT EDIT (CARD) (SKIP,A);
5580 SELECT (SUBSTR(CARD,1,5));

      WHEN ('*NUNT','*TERM') /* NUNTERMINAL & TERMINALS */
      DO; L = 6;

          UNONT:
          DO K = L BY 1 UNTIL(A_CARD(K) ~= ' '); END;
          IF K > 80
            THEN LEAVE; /* DONE WITH SCAN */
          DO L = K BY 1 UNTIL(A_CARD(L) = ' '); ENO;
          IF L > 80
            THEN L = 80; /* ASSUME MAX VALUE */
          I = I + 1;
          SUBSTR(SYMBOLS(I),FULL,1,L-K) = SUBSTR(CARD,K,L-K);
          SYMBOLS(I):SYM = REPSYM(RSND);
          SYMBOLS(I+1):SYM = ' ';
          RSND = RSND + 1;
          IF SUBSTR(CARD,I,5) = 'NONT'
            THEN SYMBOLS(I):TYPE = '1'b;
            ELSE SYMBOLS(I):TYPE = '0'b;
          SYMBOLS(I):L_FULL = L-K;
          GOTO GNONT;
        ENO;

      WHEN ('*PROD') /* PRODUCTION */
      DO;

```

ZEF05330
 ZEF05340
 ZEF05350
 ZEF05360
 ZEF05370
 ZEF05380
 ZEF05390
 ZEF05400
 ZEF05410
 ZEF05420

 ZEF05430
 ZEF05440
 ZEF05450
 ZEF05460
 ZEF05470
 ZEF05480
 ZEF05490
 ZEF05500
 ZEF05510
 ZEF05520

 ZEF05530
 ZEF05540
 ZEF05550
 ZEF05560
 ZEF05570
 ZEF05580

 ZEF05590
 ZEF05600
 ZEF05610
 ZEF05620
 ZEF05630
 ZEF05640
 ZEF05650
 ZEF05660
 ZEF05670
 ZEF05680
 ZEF05690
 ZEF05700
 ZEF05710
 ZEF05720
 ZEF05730
 ZEF05740
 ZEF05750
 ZEF05760
 ZEF05770
 ZEF05780
 ZEF05790

 ZEF05800
 ZEF05810

NUMBER LEV NT

```

5820      J = J + 1;
5830      PROD(J).NO = J;
5840      PROD(J).RP = 0;
5850      PROD(J).RP = 0;
5860      DO K = 6 TO 80 UNTIL(A_CARD(K) = ' '); END;
5870      DO L = 6 TO 80 UNTIL(A_CARD(L) = ' '); END;
5880      IF (K = 80) | (L = 80)
5890      THEN DO;
5900          PUT EDIT ('INVALID PRODUCTION.') (A) SKIP;
5910          GOTO GPEND;
5920      END;
5930      DO L = L BY -1 UNTIL(A_CARD(L) = ' '); END;
5940      PROD(J).LP = INT(SUBSTR(CARD,K,L-K+1));
5950      L = L + 1;
5960      GPROD:
5970      DO L = L BY 1 UNTIL ((A_CARD(L) = ' ') &
5980          (A_CARD(L) = ' ')); END;
5990      IF L >= 80 THEN LEAVE;
6000      DO K = L BY 1 UNTIL(A_CARD(K) = ' '); END;
6010      M = M + 1;
6020      SUBSTR(PROD(J).RP,M,1) = INT(SUBSTR(CARD,L,K-L));
6030      L = K;
6040      PROD(J).NO_MP = M;
6050      GOTO GPROD;
6060      GPEND:
6070      END;

6090      OTHERWISE PUT EDIT ('UNRECOGNIZED IC.') (SKIP,A);
6100      END; /* END SELECT */
6110      READ FILE (IN) INTO (CARD);
6120      END;
6130      END; /* GRAMMAR */

```

5820 ZEF05820
5830 ZEF05830
5840 ZEF05840
5850 ZEF05850
5860 ZEF05860
5870 ZEF05870
5880 ZEF05880
5890 ZEF05890
5900 ZEF05900
5910 ZEF05910
5920 ZEF05920
5930 ZEF05930
5940 ZEF05940
5950 ZEF05950
5960 ZEF05960
5970 ZEF05970
5980 ZEF05980
5990 ZEF05990
6000 ZEF06000
6010 ZEF06010
6020 ZEF06020
6030 ZEF06030
6040 ZEF06040
6050 ZEF06050
6060 ZEF06060
6070 ZEF06070
6080 ZEF06080
6090 ZEF06090
6100 ZEF06100
6110 ZEF06110
6120 ZEF06120
6130 ZEF06130

NUMBER LEV NT

```

6140 1 0 INT: PROC(GSYM) RETURNS(CHAR(1));
/*-----
   GIVEN FULL GRAMMAR SYMBOL THIS PROCEDURE RETURNS INTERNAL
   GRAMMAR SYMBOL.
   -----*/
6190 2 0 DCL GSYM CHAR(30);
6200 2 0 DCL SYM CHAR(1);
6210 2 0 DCL (I,J) FIXED BIN(15);
6220 2 0 DO I = 1 TO NO_SYM;
6230 2 1 J = SYMBOLS(I).L-FULL;
6240 2 1 IF SUBSTR(SYMBOLS(I).FULL,I,J) = SUBSTR(GSYM,1,J)
        THEN DO;
6260 2 2 SYM = SYMBOLS(I).SYM;
6270 2 2 RETURN (SYM);
6280 2 2 END;
6290 2 1 PUT EDIT ('UNDEFINED GRAMMAR SYMBOL: ',GSYM)
        (2 A) SKIP;
6300 2 0 SYM = 'X';
6320 2 0 RETURN (SYM);
6330 2 0 ENDD; /* INT */
6340 2 0 ENDD; /* INT */

6350 1 0 ALTER: PROC (ONE,TWO); AND STATE TWO ARE EQUAL.
/* STATE ONE
   MERGE ACTION ENTRIES FOR STATE TWO INTO STATE ONE. */
6380 2 0 DCL (ONE,TWO,I) FIXED BIN(15);
6390 2 0 DO I = 1 BY 1 UNTIL (SYMBOLS(I).SYM = ' ');
6400 2 1 IF ACTION (I,ONE) = -1 /* UNDEFINED SLOT */
        THEN ACTION (I,ONE) = ACTION (I,TWO);
6420 2 1 ENDD;
6430 2 0 ENDD; /* ALTER */
ZEF06140
ZEF06150
ZEF06160
ZEF06170
ZEF06180
ZEF06190
ZEF06200
ZEF06210
ZEF06220
ZEF06230
ZEF06240
ZEF06250
ZEF06260
ZEF06270
ZEF06280
ZEF06290
ZEF06300
ZEF06310
ZEF06320
ZEF06330
ZEF06340
ZEF06350
ZEF06360
ZEF06370
ZEF06380
ZEF06390
ZEF06400
ZEF06410
ZEF06420
ZEF06430

```


NUMBER LEV NT

```

6440 1 0 LALR: PROC;
6450 2 0 DCL TRANSFORM(10:25) FIXED BIN(15) STATIC INIT(
      0,1,2,3,4,5,6,7,8,9,10,11,12,13,14,15,16,17,18,19,20,
      21,22,23,24,25);
6480 2 0 DCL LALR_ST FIXED BIN(15);
6490 2 0 DCL DELETE(15) FIXED BIN(15) STATIC INIT ((15) 32767);

6500 2 0 UCL (1,J,K,L,M,N) FIXED BIN(15);
6510 2 0 PUT EDIT ('LALR TRANSFORMATION:') (A) PAGE;
6520 2 0 LALR_ST = STATE; /* GET CURRENT TOTAL NO. OF STATES */
6530 2 0 DO I = 0 BY 1 WHILE (I < STATE);
6540 2 1 DO J = 1 BY 1 WHILE (J < STATE);
6550 2 2 IF (TRANSFORM(I) = J)
      THEN IF IDENT(I,J) /* CHECK FOR STATE IDENTITY */
      THEN DO;
          TRANSFORM(J) = I;
          CALL ALTER(I,J); /* ALTER ACTION */
          LALR_ST = LALR_ST + 1;
        END;
      END;

6590 2 3 /* DETERMINE DELETED STATES */
6600 2 0 K = 1;
6610 2 0 DO I = 0 TO STATE;
6680 2 1 IF TRANSFORM(I) = I
      THEN;
      ELSE /* I -> EQUIVALENT STATE */
        DO;
          DELETE(K) = I;
          K = K + 1;
          DELETE(K) = -1;
        END;
      END;

6700 2 1
6720 2 2
6730 2 2
6740 2 2
6750 2 2
6760 2 1

6770 2 0 PUT EDIT ('LR(1) STATES TO DELETE:') (A) SKIP(2);

6780 2 0 PUT EDIT (' ') (A) SKIP;
6790 2 0 DO J = 1 TO 15 WHILE (DELETE(J) /= -1);
6800 2 1 PUT EDIT (DELETE(J)) (PZZ9);
6810 2 1
6820 2 1
6830 2 0
6840 2 0 PUT EDIT ('LR(1) SKIP;') (A) SKIP(2);
6850 2 0 DO J = 0 TO NEWSTATE;
6860 2 1 PUT EDIT (J) (PZZ9);
      END;

/* FILL DELETED STATES WITH VALID STATES */
K = 1;
DO I = 0 TO STATE-1;
  IF TRANSFORM(I) /= I
    THEN /* DELETED STATE - NU SWEAT */

```

```

ZEF06440
ZEF06450
ZEF06460
ZEF06470
ZEF06480
ZEF06490

ZEF06500
ZEF06510
ZEF06520
ZEF06530
ZEF06540
ZEF06550
ZEF06560
ZEF06570
ZEF06580
ZEF06590
ZEF06600
ZEF06610
ZEF06620
ZEF06630
ZEF06640

ZEF06650
ZEF06660
ZEF06670
ZEF06680
ZEF06690
ZEF06700
ZEF06710
ZEF06720
ZEF06730
ZEF06740
ZEF06750
ZEF06760

ZEF06770

ZEF06780
ZEF06790
ZEF06800
ZEF06810
ZEF06820
ZEF06830
ZEF06840
ZEF06850
ZEF06860

ZEF06870
ZEF06880
ZEF06890
ZEF06900
ZEF06910

```

NUMBER LEV NT

```

6920 2 1      ELSE
        IF TRANSFORM(I) < LALR_ST
        THEN: /* STATE WITHIN NEW LIMITS */
        ELSE /* BEYOND UPPER LIMIT - USE VALID */
        DO: /* DELETED STATE.
            IF DELETE(K) >= LALR_ST
            THEN DO:
                PUT EDIT ('54*0',K) (A,P'Z9') SKIP:
                SIGNAL ERROR;
            END;
            L = TRANSFORM(I);
            DO J = 0 TO STATE:
                IF TRANSFORM(J) = L
                THEN TRANSFORM(J) = DELETE(K);
            END;
            GOTOS(*,DELETE(K)) = GOTOS(*,L);
            ACTION(*,DELETE(K)) = ACTION(*,L);
            K = K+1;
        END;
        END;
        PUT EDIT ('LALR') (A) SKIP;
        DO J = 0 TO NEWSTATE:
            DO PUT EDIT (TRANSFORM(J)) (P'Z9');
        END;
        DO I = 1 TO NO_SYM UNTIL (SYMBOLS(I).SYM = ' '); END;
        LALR_SZ.M = I;
        LALR_SZ.N = LALR_ST - 1;
        ALLOCATE LALR_GOTO;
        ALLOCATE LALR_ACT;

        /* FILL IN ENTRIES IN LALR ARRAYS */
        DO I = 1 TO LALR_SZ.M;
        DO J = 0 TO LALR_SZ.N;
        IF GOTOS(I,J) >= 0
        THEN LALR_GOTO (I,J) = TRANSFORM (GOTOS(I,J));
        ELSE LALR_GOTO (I,J) = GOTOS (I,J);
        IF ACTION (I,J) >= 0
        THEN LALR_ACT (I,J) = TRANSFORM (ACTION(I,J));
        ELSE LALR_ACT (I,J) = ACTION (I,J);
        END;
        END;
        END; /* LALR */
7320 2 0      END;

```

ZEF06920

ZEF06930

ZEF06940

ZEF06950

ZEF06960

ZEF06970

ZEF06980

ZEF06990

ZEF07000

ZEF07010

ZEF07020

ZEF07030

ZEF07040

ZEF07050

ZEF07060

ZEF07070

ZEF07080

ZEF07090

ZEF07100

ZEF07110

ZEF07120

ZEF07130

ZEF07140

ZEF07150

ZEF07160

ZEF07170

ZEF07180

ZEF07190

ZEF07200

ZEF07210

ZEF07220

ZEF07230

ZEF07240

ZEF07250

ZEF07260

ZEF07270

ZEF07280

ZEF07290

ZEF07300

ZEF07310

ZEF07320

NUMBER LEV NT

```

7330 1 0 IDENT: PROC (FIRST,SECOND) RETURNS (BIT);
      /* RETURN '1' IF FIRST STATE IDENTICALLY
      /* EQUALS SECOND STATE.
      /* CHECK FOR IDENTITY IF ITEM HAS SAME PRODUCTION
      /* AND SAME NO. OF RECOGNIZED POSITIONS. */
      ZEF07330
      ZEF07340
      ZEF07350
      ZEF07360
      ZEF07370

7380 2 0 UCL (FIRST,SECOND,I,J,K,L,M) FIXED BIN(15);
      DO I = 1 TO NEXT_I-1 WHILE (ITEMS(I).STATE ^= FIRST); END;
      DO J = 1 TO NEXT_J-1 WHILE (ITEMS(J).STATE ^= SECOND); END;
      L=J;
      DO M = 1 TO 2;
      DO UNTIL (I = -1);
      IF M = 1
      THEN J = L;
      ELSE J = K;
      DO UNTIL (J = -1);
      IF (ITEMS(I).PROD = ITEMS(J).PROD) &
      (ITEMS(I).POSIT = ITEMS(J).POSIT)
      THEN GOTO IDENT1;
      IF J = -1
      THEN
      ELSE J = ITEMS(J).NEXT;
      END;
      RETURN ('0'B); /* STATES NOT EQUAL */
      ZEF07380
      ZEF07390
      ZEF07400
      ZEF07410
      ZEF07420
      ZEF07430
      ZEF07440
      ZEF07450
      ZEF07460
      ZEF07470
      ZEF07480
      ZEF07490
      ZEF07500
      ZEF07510
      ZEF07520
      ZEF07530
      ZEF07540
      ZEF07550
      ZEF07560

7570 2 2 IDENT1:
      IF I = -1
      THEN
      ELSE I = ITEMS(I).NEXT;
      END;
      I = L; /* NOW ENSURE ALL ITEMS IN SECOND STATE */
      J = K; /* ARE CONTAINED IN FIRST STATE. */
      END;
      RETURN ('1'B);
      ZEF07570
      ZEF07580
      ZEF07590
      ZEF07600
      ZEF07610
      ZEF07620
      ZEF07630
      ZEF07640
      ZEF07650

7660 2 0 END; /* IDENT */
      ZEF07660

7670 1 0 END; /* APARSE */
      ZEF07670

```

OPTIONS SPECIFIED

MAR(2 72 1) NUM GN MACRO NIS TERM;

OPTIONS USED

AGGREGATE	NOCOUNT	ATTRIBUTES(SHORT)
COMPILE	NODECK	CHARSET(60,EBCDIC)
CONUMBER	NODESU	FLAG(1)
MESSAGE	NOFLOW	LINECOUNT(90)
MACRO	NOJUSTMT	MARGINS(2,72,1)
MAP	NOIMPRECISE	OPTIMIZE(TIME)
TEST	NOINCLUDE	SEQUENCE(75,80)
NUMBER	NOINTERUPT	SIZE(32740)
OBJECT	NOINSOURCE	NUSYNTEX(5)
OPTIONS	NOLIST	XREF(SHORT)
SOURCE	NOMARGIN1	TERMINAL(NOAGGREGATE,
STORAGE	NODECK	NOATTRIBUTES,
	NUSTMT	NODESU,
		NOINSOURCE,
		NOLIST,
		NOMAP,
		NOMFSET,
		NNOPTIONS,
		NUSOURCE,
		NOSTORAGE,
		NOMXREF)

NO PREPROCESSOR DIAGNOSTIC MESSAGES PRODUCED

SOURCE LISTING

NUMBER LEV NT

R

```

20      0  ZLPSDL:  PROC OPTIUNS(MAIN) REORDER;
/*-----
/* FILE DEFINITIONS
60      1  0  DCL SYSPRINT STREAM OUTPUT PRINT FILE
      ENV ( V8 RECSIZE(121) BLKSIZE(125));
80      1  0  DCL STRINGS FILE RECORD INPUT; /* INPUT STRINGS FROM GRAMMAR */
90      1  0  DCL SYMTAB FILE RECORD INPUT; /* SYMBOL TABLE */
100     1  0  DCL PROUTAB FILE RECORD INPUT; /* PRODUCTIONS TABLE */
110     1  0  DCL LALKSZ FILE RECORD INPUT; /* LALK TABLE SIZES */
120     1  0  DCL LALRTAB FILE RECORD INPUT; /* LALK TABLES */

150     1  0  DCL NO_SYM FIXED BIN(15) STATIC INIT( 40);

160     1  0  UCL 1 SYMBOLS ( 40) STATIC;
      2 TYPE CHAR(1) INIT('A','S','C','B','U',' ');
      2 L_FULL FIXED BIN(15);
      2 L_FULL CHAR(30);

230     1  0  DCL NO_PROD FIXED BIN(15) STATIC INIT( 20);

240     1  0  UCL 1 PROD ( 20) STATIC;
      2 NO FIXED BIN(15);
      2 LP CHAR(1) INIT('R','S','C','C');
      2 NO_KP FIXED BIN(15) INIT(1,2,2,1);
      2 RP CHAR(6) INIT('S','CC','BC','D');

290     1  0  DCL 1 LALK_SZ, /* LALK GOTO & ACTION ARRAYS */
      2 M FIXED BIN(15); /* NO. OF GRAMMAR SYMBOLS */
      2 N FIXED BIN(15); /* NO. OF STATES */

320     1  0  DCL GOTOS (LALK_SZ,M,0:LALK_SZ,M) FIXED BIN (15) CTL;
330     1  0  UCL ACTION (LALK_SZ,M,0:LALK_SZ,M) FIXED BIN (15) CTL;

340     1  0  DCL STACK(20) FIXED BIN(15) STATIC INIT(1); /* PROCESSING STACK */
350     1  0  DCL SPTR FIXED BIN(15) STATIC INIT(1); /* STACK POINTER */

360     1  0  DCL INPUT CHAR(60) STATIC INIT(' ');
370     1  0  UCL A_INPT(80) CHAR(1) BASED(P_INP);
380     1  0  DCL P_INP PTR INIT(ADDR(INPUT));

390     1  0  UCL SYM CHAR(1);
400     1  0  UCL N_SYM FIXED BIN(15);

```

ZEF00010
ZEF00020

ZEF00030
ZEF00040
ZEF00050
ZEF00060
ZEF00070
ZEF00080
ZEF00090
ZEF00100
ZEF00110
ZEF00120

ZEF00150

ZEF00160
ZEF00170
ZEF00180
ZEF00190
ZEF00200

ZEF00230

ZEF00240
ZEF00250
ZEF00260
ZEF00270
ZEF00280

ZEF00290
ZEF00300
ZEF00310

ZEF00320
ZEF00330

ZEF00340
ZEF00350

ZEF00360
ZEF00370
ZEF00380

ZEF00390
ZEF00400

NUMBER LEV NT

/* MISC DECLARATIONS */

```

420 1 0 DCL (I,J,K,L) FIXED BIN(15);
430 1 0 DCL (M,N,P) FIXED DEC(3,0);
440 1 0 DCL ADDR BUILTIN;
450 1 0 DCL (MORE_IN,NEW_STK) BIT(1) STATIC INIT('1'b);

```

R

```

ZEF00410
ZEF00420
ZEF00430
ZEF00440
ZEF00450

```

NUMBER LEV NT

R

```

470 1 0 /* ON UNITS AND INITIALIZATION */
480 1 0 UN SUBRG SIGNAL ERRK;
      ON STRG SIGNAL ERRK;

490 1 0 UN ERKGA SNAP
510 2 0 BEGIN;
      END;

520 1 0 UN ENDFILE(STRINGS)
      BEGIN;
          A_INPUT(1) = 'S';
          MORE_IN = 'O'B;
      END;

540 2 0
550 2 0
560 2 0

570 1 0 UPEN FILE(SYSPRINT),
      FILE(STKINGS),
      FILE(SYMTAB),
      FILE(PRODTAB),
      FILE(LALRTAB),
      FILE(LALRSZ);

630 1 0 READ FILE(SYMTAB) INTO(SYMBOLS);
640 1 0 DO I = 1 TO NO_SYM UNTIL(SYMBOLS(1).SYM = ' '); END;
650 1 0 SYMBOLS(1).SYM = 'S';
660 1 0 SYMBOLS(I+1).SYM = 'S';
670 1 0 SYMBOLS(1).TYPE = 'O'B;
680 1 0 READ FILE(PRODTAB) INTO(PROD);
690 1 0 READ FILE(LALRSZ) INTO(LALKSZ);
700 1 0 ALLOCATE GOTDS;
710 1 0 ALLOCATE ACTION;

      DCL TBL(2) FIXED BIN(15);
      DCL C_TBL CHAR(4) BASED (P_TBL);
      DCL P_TBL PTR;
      P_TBL = ADDR (TBL(1));

      DO I = 1 TO LALR_SZ:M;
          DO J = 0 TO LALKSZ:N;
              READ FILE (LALRTAB) INTO (C_TBL);
              GOTUS (1,J) = TBL(1);
              ACTION (I,J) = TBL(2);
          END;
      END;

      CLOSE FILE(SYMTAB);
      FILE(PRODTAB);
      FILE(LALRTAB);
      FILE(LALRSZ);

```

ZEF00460
ZEF00470
ZEF00480
ZEF00490
ZEF00500
ZEF00510
ZEF00520
ZEF00530
ZEF00540
ZEF00550
ZEF00560
ZEF00570
ZEF00580
ZEF00590
ZEF00600
ZEF00610
ZEF00620
ZEF00630
ZEF00640
ZEF00650
ZEF00660
ZEF00670
ZEF00680
ZEF00690
ZEF00700
ZEF00710
ZEF00720
ZEF00730
ZEF00740
ZEF00750
ZEF00760
ZEF00770
ZEF00780
ZEF00790
ZEF00800
ZEF00810
ZEF00820
ZEF00830
ZEF00840
ZEF00850
ZEF00860

NUMBER' LEV NT

R

```

/*-----
MAIN PROCESSING
-----*/

900 1 0 CALL GETSYM;
910 1 0 DO WHILE(MORE IN);
920 1 1 SELECT (ACTION(N_SYM,STACK(SPTR)));
930 1 2 WHEN (-1)
940 1 3 DO; PUT EDIT (, INVALID STRING FOR THIS GRAMMAR*,
950 1 4 (A,SKIP,A) SKIP;
960 1 3 DO WHILE(SYM ^= '$$');
970 1 4 CALL GETSYM;
980 1 3 ENU;
990 1 4 CALL GETSYM;
1000 1 3 ENU;
1010 1 4 CALL GETSYM;
1020 1 3 ENU;
1030 1 2 WHEN (0)
1040 1 3 DO; PUT EDIT (, REDUCE USING PRODUCTION NO. *,1)
1050 1 4 (A,P,ZZ9') SKIP;
1060 1 3 PUT EDIT (, VALID INPUT STRING*) (A) SKIP;
1070 1 4 CALL GETSYM;
1080 1 3 ENU;
1090 1 4 CALL GETSYM;
1100 1 2 WHEN (-2) /* SHIFT */
1110 1 3 DO; CALL PUSH(N_SYM);
1120 1 4 CALL PUSH(GOTOS(N_SYM,STACK(SPTR-1)));
1130 1 3 M = STACK(SPTR);
1140 1 4 PUT EDIT (, STACK INPUT CHARACTER AND STATE *,M)
1150 1 3 PUT EDIT (A,P,ZZ9') SKIP;
1160 1 4 CALL GETSYM;
1170 1 3 ENU;
1180 1 4 CALL GETSYM;
1190 1 2 OTHERWISE /* REDUCE */
1200 1 3 DO; 1 = ACTION(N_SYM,STACK(SPTR)); /* REDUCE PROD NO. */
1210 1 4 J = PROD(1).NU,RP * 2;
1220 1 3 SPTR = SPTR - 1; /* POP OFF STACK J ELEMENTS */
1230 1 4 K = PTRSYM(PROD(1).LP);
1240 1 3 CALL PUSH(K);
1250 1 4 CALL PUSH(GOTOS(K,STACK(SPTR-1)));
1260 1 3 IF STACK(SPTR) < 1
1270 1 4 THEN DO; PUT EDIT(, INVALID STRING FOR THIS GRAMMAR*,
1280 1 5 (A,SKIP,A);
1290 1 4 DO WHILE(SYM ^= '$$');
1300 1 5 CALL GETSYM;
1310 1 4 ENU;
1320 1 5 CALL GETSYM;
1330 1 4 ENU;
1340 1 5 CALL GETSYM;
1350 1 4 ENU;

```

ZEF00870
ZEF00880
ZEF00890

ZEF00900
ZEF00910
ZEF00920

ZEF00930
ZEF00940
ZEF00950
ZEF00960
ZEF00970
ZEF00980
ZEF00990
ZEF01000
ZEF01010
ZEF01020

ZEF01030
ZEF01040
ZEF01050
ZEF01060
ZEF01070
ZEF01080
ZEF01090

ZEF01100
ZEF01110
ZEF01120
ZEF01130
ZEF01140
ZEF01150
ZEF01160
ZEF01170
ZEF01180

ZEF01190
ZEF01200
ZEF01210
ZEF01220
ZEF01230
ZEF01240
ZEF01250
ZEF01260
ZEF01270
ZEF01280
ZEF01290
ZEF01300
ZEF01310
ZEF01320
ZEF01330
ZEF01340
ZEF01350

NUMBER LEV NT

1360 1 4
1370 1 3
1380 1 3
1400 1 3
1410 1 2
1420 1 1

END: M = I;
PUT EDIT (' REDUCE USING PRODUCTION NO. ',M)
(A,P,ZZ9;) SKIP;
END;
END: /* SELECT */
END;

R

ZEF01360
ZEF01370
ZEF01380
ZEF01390
ZEF01400
ZEF01410
ZEF01420

NUMBER LEV NT

/*-----
 -----*/
 PUSH PRDC

1460	1	0	PUSH: PROC(VAL);
1470	2	0	DCL VAL FIXED BIN(15);
1480	2	0	SPR = SPTR + 1;
1490	2	0	STACK(SPTR) = VAL;
1500	2	0	END; /* PUSH */

/*-----
 -----*/
 PTR_SYM PROC

1540	1	0	PTR_SYM: PROC(SYM);
1550	2	0	DCL SYM CHAR(1);
1560	2	0	DCL I FIXED BIN(15);
1570	2	0	DO I = 1 TO NO_SYM UNTIL(SYMBOLS(I).SYM = SYM);
1580	2	0	RETURN(I);
1590	2	0	END; /* PTR_SYM */

R

ZEF01430
 ZEF01440
 ZEF01450

ZEF01460
 ZEF01470

ZEF01480
 ZEF01490

ZEF01500

ZEF01510
 ZEF01520
 ZEF01530

ZEF01540

ZEF01550
 ZEF01560

ZEF01570
 ZEF01580

ZEF01590

PL/I OPTIMIZING COMPILER

/* PARSER */

NUMBER LEV NT

R

```

1630 1 0 GETSYM: PROC;
1640 2 0 UCL 1 FIXED BIN(15);
1650 2 0 IF NEW_STK
      THEN DO;
1670 2 1 PUT EDIT ('INPUT STRING: ') (A) PAGE;
1680 2 1 NEW_STR = '0'B;
1690 2 1 STACK(1) = 0;
1700 2 1 SSTR = 1;
1710 2 1 END;
1720 2 0 GETC: DO 1 = 1 TO 80 UNTIL(A_INPUT(1) ^= ' '); END;
1730 2 1 IF 1 >= 80
1740 2 0 THEN
      DO;
1770 2 1 READ FILE(STRINGS) INTO(INPUT);
1780 2 1 PUT EDIT (INPUT) (A) SKIP;
1790 2 1 PUT EDIT (' ') (A) SKIP;
1800 2 1 M=1;
1810 2 1 N=1;
1820 2 1 AGAIN:
1830 2 1 DO J = M TO 80 UNTIL(A_INPUT(J) ^= ' '); END;
1840 2 1 IF J >= 80
      THEN DO;
1860 2 2 A_INPUT(N) = ' ';
1870 2 2 N = N+1;
1880 2 2 SUBSTR(INPUT,N,80-N) = ' ';
1890 2 2 GOTO GETC;
1900 2 2 END;
1910 2 2 DO K = J BY 1 UNTIL(A_INPUT(K) = ' '); END;
1920 2 2 A_INPUT(N) = INT(SUBSTR(INPUT,J,K-J));
1930 2 2 N = N+1;
1940 2 2 M = K;
1950 2 2 GOTO AGAIN;
1960 2 2 END;
1970 2 0 SYM = A_INPUT(1);
1980 2 0 A_INPUT(1) = ' ';
1990 2 0 N_SYM = PTR_SYM(SYM);
2000 2 0 PUT EDIT
      (SUBSTR(SYMBOLS(N_SYM),FULL,1,SYMBOLS(N_SYM).L_FULL),) => ' ,SYM)
      IF SYM = '3'A SKIP;
2030 2 0 THEN NEW_STR = '1'B;
2050 2 0 END; /* GETSYM */
      GETSYM PROC
    
```

ZEF01600
 ZEF01610
 ZEF01620
 ZEF01630
 ZEF01640
 ZEF01650
 ZEF01660
 ZEF01670
 ZEF01680
 ZEF01690
 ZEF01700
 ZEF01710
 ZEF01720
 ZEF01730
 ZEF01740
 ZEF01750
 ZEF01760
 ZEF01770
 ZEF01780
 ZEF01790
 ZEF01800
 ZEF01810
 ZEF01820
 ZEF01830
 ZEF01840
 ZEF01850
 ZEF01860
 ZEF01870
 ZEF01880
 ZEF01890
 ZEF01900
 ZEF01910
 ZEF01920
 ZEF01930
 ZEF01940
 ZEF01950
 ZEF01960
 ZEF01970
 ZEF01980
 ZEF01990
 ZEF02000
 ZEF02010
 ZEF02020
 ZEF02030
 ZEF02040
 ZEF02050

NUMBER LEV NT

R

2060 1 0 INT: PRGC(GSYM) RETURNS(CHAR(1)):

ZEF02060

```

/*-----
  GIVEN FULL GRAMMAR SYMBOL THIS PROCEDURE RETURNS INTERNAL
  GRAMMAR SYMBOL.
  -----*/

```

ZEF02070
ZEF02080
ZEF02090
ZEF02100

```

2110 2 0 DCL GSYM CHAR(30);
2120 2 0 DCL SYM CHAR(1);
2130 2 0 DCL (1,J) FIXED BIN(15);

```

ZEF02110
ZEF02120
ZEF02130

```

2140 2 0 DO I = 1 TO NO-SYM;
2150 2 1 J = SYMBOLS(I).L-FULL;
2160 2 1 IF SUBSTR(SYMBOLS(I).FULL,1,J) = SUBSTR(GSYM,1,J)
    THEN DO;
    SYM = SYMBOLS(I).SYM;
    RETURN (SYM);
    END;

```

ZEF02140
ZEF02150
ZEF02160
ZEF02170
ZEF02180
ZEF02190
ZEF02200
ZEF02210
ZEF02220
ZEF02230
ZEF02240
ZEF02250

```

2180 2 2 PUT EDIT ('UNDEFINED GRAMMAR SYMBOL: ',GSYM)
2190 2 2 SYM = 'X';
2200 2 2 RETURN (SYM);
2210 2 2
2220 2 0 END; /* INT */

```

ZEF02260

2270 1 0 END: /* ZEFSDL */

ZEF02270