

Rochester Institute of Technology

## RIT Digital Institutional Repository

---

Theses

---

4-1-1997

### Evolution of solutions to real-time problems

Greg Semeraro

Follow this and additional works at: <https://repository.rit.edu/theses>

---

#### Recommended Citation

Semeraro, Greg, "Evolution of solutions to real-time problems" (1997). Thesis. Rochester Institute of Technology. Accessed from

This Thesis is brought to you for free and open access by the RIT Libraries. For more information, please contact [repository@rit.edu](mailto:repository@rit.edu).

# Evolution of Solutions to Real-Time Problems

by

**Greg P. Semeraro**

A Thesis Submitted  
in  
Partial Fulfillment of the  
Requirements for the Degree of  
MASTER OF SCIENCE  
in  
Computer Engineering

Approved by:

---

Dr. Tony H. Chang, Professor

---

Dr. Roy Czernikowski, Professor and Department Head

Department of Computer Engineering  
College of Engineering  
Rochester Institute of Technology  
Rochester, New York  
April, 1997

# THESIS RELEASE PERMISSION FORM

ROCHESTER INSTITUTE OF TECHNOLOGY

COLLEGE OF ENGINEERING

Title: The Evolution of Solutions to Real-Time Problems

I, Greg P. Semeraro, hereby grant permission to the Wallace Memorial Library to reproduce my thesis in whole or part.

---

Greg P. Semeraro

*29 April 1997*

Date

# **Evolution of Solutions to Real-Time Problems**

by

**Greg P. Semeraro**

Copyright © 1997 by Greg P. Semeraro

All Rights Reserved

# Abstract

This thesis develops the theory and tools necessary for the determination of a near optimal Real-Time Operating System (RTOS) scheduling policy for an arbitrary multitasking problem specification. The solution is determined using a Genetic Algorithm (GA).

All real-time operating systems provide some means of 'tuning' the characteristics of the scheduling policy to accurately meet the application requirements. This thesis shows the applicability of using a GA to determine these parameters for an arbitrary application. In addition, the RTOS parameters considered are broad enough to allow the results to be used for specifying and/or choosing an RTOS for the actual implementation of a real-time system.

The domain of real-time applications which is of particular interest is that of embedded systems. In the embedded systems domain, real-time multitasking problems are specified by a series of timing constraints, time deadlines and practical available resources. These constraints guide the analysis of the results.

A PC-based RTOS/GA tool set is the end result of this thesis and can be used for the analysis of arbitrary real-time applications.

# Table of Contents

1. Introduction.....	1
2. Background.....	10
2.1 Real-Time Applications.....	10
2.1.1 Formal Methods for Real-Time Analysis.....	15
2.2 Genetic Algorithms.....	18
2.2.1 Genetic Operators .....	26
3. Evolution of Solutions to Real-Time Problems .....	31
3.1 Real-Time Operating System (TASKING) .....	34
3.1.1 Facilities and Capabilities.....	37
3.1.2 Configuration Parameters .....	49
3.1.3 RTOS Performance.....	53
3.2 Genetic Algorithm .....	59
3.2.1 Genetic Representation.....	61
3.2.2 Mutation.....	62
3.2.3 Crossover .....	63
3.3 Problem Specification.....	64
3.3.1 Specification Parameters .....	64
3.3.2 Specification File Format .....	66
4. Evolution Tool Set.....	68
4.1 Custom RTOS (TASKING).....	68

4.2 Real-Time Application (RTOS-APP).....	71
4.3 Genetic Algorithm (GRTOS-GA) .....	71
5. Test Suite Description.....	74
5.1 Verification of Capabilities .....	75
5.2 Verification of Performance .....	77
5.3 Test of Capabilities .....	79
5.4 Test of Performance.....	85
6. Conclusions.....	92
7. Bibliography .....	94
8. Appendix A RTOS Source Code.....	96
8.1 TASKING.INT .....	96
8.2 TASKING.PAS.....	99
9. Appendix B Genetic Algorithm Source Code.....	117
9.1 RTOS-APP.PAS .....	117
9.2 GRAPH-GA.PAS.....	123
9.3 RTOS_GA.PAS .....	132
9.4 MATH_GA.PAS.....	136
10. Appendix C Support Software Source Code.....	140
10.1 DINING.PAS .....	140
10.2 TSK-BNCH.PAS .....	142
10.3 TASKS-A.PAS .....	146
10.4 TASKS-B.PAS.....	147

10.5 GRAPHICS.PAS .....	147
10.6 BMP_UTIL.PAS.....	148
10.7 MS_MOUSE.PAS .....	150



# Table of Figures

FIGURE 1 - TASK EXECUTION PROFILE.....	6
FIGURE 2 - GENETIC ALGORITHM .....	8
FIGURE 3 - REACTION RATE CONTROLLER (RRC).....	12
FIGURE 4 - GENERAL GENETIC ALGORITHM .....	20
FIGURE 5 - SAMPLE OPTIMIZATION FUNCTION PLOT .....	24
FIGURE 6 - SAMPLE OPTIMIZATION FUNCTION PLOT (LOCALIZED).....	25
FIGURE 7 - OPERATING SYSTEM FACILITIES/STRUCTURE .....	39
FIGURE 8 - UNBOUNDED PRIORITY INVERSION .....	51
FIGURE 9 - PRIORITY INHERITANCE PROTOCOL.....	52
FIGURE 10 - TASKING BENCHMARK (RELATIVE PERFORMANCE).....	55
FIGURE 11 - TASKING BENCHMARK (CONTEXT SWITCHING OVERHEAD).....	56
FIGURE 12 - TASKING BENCHMARK (CONTEXT SWITCHING PERCENTAGE).....	57
FIGURE 13 - TASKING BENCHMARK (SYSTEM CLOCK IMPACT).....	59
FIGURE 14 - GENETIC ALGORITHM PROCESS FLOW .....	60
FIGURE 15 - TASK EXECUTION PROFILE.....	65
FIGURE 16 - SAMPLE TASKING.RPT FILE.....	70
FIGURE 17 - SAMPLE RTOS-APP.RPT FILE.....	71
FIGURE 18 - GRTOS-GA STARTUP SCREEN .....	72
FIGURE 19 - REAL-TIME PROBLEM #1 .....	75

FIGURE 20 - PROBLEM #1, 'RANDOM' ASSIGNMENT ANALYSIS - RANK #3 .....	77
FIGURE 21 - REAL-TIME PROBLEM #2 .....	78
FIGURE 22 - REAL-TIME PROBLEM #3 .....	80
FIGURE 23 - PROBLEM #3, 'RANDOM' ASSIGNMENT ANALYSIS - RANK #2 .....	82
FIGURE 24 - PROBLEM #3, 'RANDOM' ASSIGNMENT ANALYSIS - RANK #3 .....	83
FIGURE 25 - PROBLEM #3, 'RANDOM' ASSIGNMENT ANALYSIS - RANK #5 .....	84
FIGURE 26 - REAL-TIME PROBLEM #4 .....	88
FIGURE 27 - PROBLEM #4, 'RANDOM' ASSIGNMENT ANALYSIS - RANK #3 .....	90
FIGURE 28 - PROBLEM #4, 'RANDOM' ASSIGNMENT ANALYSIS - RANK #5 .....	91

# Table of Tables

TABLE 1 - RRC REQUIREMENTS .....	12
TABLE 2 - GA SOLUTIONS TO SAMPLE FUNCTION .....	26
TABLE 3 - RTOS GENOTYPE .....	62
TABLE 4 - REAL-TIME APPLICATION SPECIFICATION .....	67
TABLE 5 - GRTOS-GA.EXE RESULTS, PROBLEM #1 .....	76
TABLE 6 - GRTOS-GA.EXE RESULTS, PROBLEM #2 .....	79
TABLE 7 - GRTOS-GA.EXE RESULTS, PROBLEM #3 .....	81
TABLE 8 - GRTOS-GA.EXE RESULTS, PROBLEM #4 .....	89

# Table of Equations

EQUATION 1 - SAMPLE OPTIMIZATION FUNCTION.....	23
EQUATION 2 - SAMPLE OPTIMIZATION FUNCTION (LOCALIZED).....	25
EQUATION 3 - EXPECTED ALLELE COVERAGE (BINARY REPRESENTATIONS) .....	28
EQUATION 4 - EXPECTED ALLELE COVERAGE (NON-BINARY REPRESENTATIONS).28	
EQUATION 5 - TASK SCHEDULABILITY.....	74
EQUATION 6 - SCHEDULABILITY, PROBLEM #1 .....	75
EQUATION 7 - SCHEDULABILITY, PROBLEM #2 .....	78
EQUATION 8 - SCHEDULABILITY, PROBLEM #3 .....	80
EQUATION 9 - PRIORITY ASSIGNMENT CHARACTERISTICS .....	83
EQUATION 10 - SCHEDULABILITY, PROBLEM #4.....	88

# Glossary

Note: All terms appearing in this glossary will appear for the first time as underlined text in the main body of the thesis.

<u>Term</u>	<u>Page</u>	<u>Description</u>
Allele	21	Any of a group of possible mutational forms of a gene.
Epistatic Interaction	7	A genetic interaction which is nonrecipricating and is between nonalternative forms of genes in which one gene suppresses the expression of another affecting the same part of an organism.
Expected Allele Coverage	27	Given the following: A binary string encoding of length $L$ , (with are $2L$ possible alleles) and a random population of $N$ encodings. The expected allele coverage is the expected proportion of all possible alleles that occur in the population and is $1-(1/2)^N$ [Ref. 19, pg. 32].
GA	1	A Genetic Algorithm (GA) is the general term applied to computer algorithms which search for a solution to a given problem by

<u>Term</u>	<u>Page</u>	<u>Description</u>
		imitating the 'natural selection' which steers the evolution of living organisms.
Genotype	21	The genetic make-up of an organism.
NP-Complete	1	A problem is said to be NP-Complete if computing an exact solution to the problem requires <u>n</u> ondeterministic polynomial time and it has been proven to be Complete.
Nyquist Theorem	11	Nyquist Theorem [Ref. 16, pg. 519]: Let $x(t)$ be a bandlimited signal with $X(\omega) = 0$ for $ \omega  > \omega_M$ . Then $x(t)$ is uniquely determined by its samples $x(nT)$ , $n = 0, \pm 1, \pm 2, \dots$ if $\omega_S > 2\omega_M$ where $\omega_S = 2\pi/T$ .
Phenotype	21	The environmentally and genetically determined observable appearance of an organism, especially as considered with respect to all possible genetically influenced expressions of one specific character.
RTOS	1	A Real-Time Operating System (RTOS) is a computer operating

<u>Term</u>	<u>Page</u>	<u>Description</u>
TSR	59	<p>systems which provide the necessary constructs to allow for the creation of multitasking applications capable of communicating with each other and meeting critical timing constraints.</p> <p>A Terminate and Stay Resident (TSR) program is an MS-DOS program that remains in memory after the program has finished executing. TSRs typically hook into interrupt vectors and appear to execute in the background under MS-DOS. Since their execution can be sporadic, their impact on system performance is not defined.</p>

# 1. Introduction

The objective of this thesis is to develop a tool to assist the user in determining a near optimal Real-Time Operating System (RTOS) scheduling policy for an arbitrary multitasking problem specification. The solution is determined using a Genetic Algorithm (GA). The basic idea is that determining the 'best' RTOS scheduling policy is a very complicated task. This determination is usually done by an individual, or small group of individuals, with significant knowledge about the problem domain of the application. With multitasking software becoming increasingly complex, knowing that the RTOS scheduling policy that is being used is the most effective for the specific application is no longer something that can be determined merely by understanding the problem domain. What is needed is detailed knowledge about the interaction of the tasks that make up the multitasking application. Even with this knowledge, the problem of determining an optimal RTOS configuration is NP-Complete.

"The problem of determining that a given schedule is optimal has been proven to be an NP-complete problem in general. ... The most notable exception from NP-completeness is determining a schedule for a monoprocessor system, in which all tasks are initially released simultaneously. If a problem is NP-complete, it is not likely that an algorithm exists that performs better than an algorithm based on enumerative methods."

[Ref. 2, pg. 33]



For many optimization problem classes, it has been shown that GAs are uniquely qualified to search complex problem spaces for a global optimum, especially when that search space contains numerous local optima.

“... The developed evolution program displayed some qualities not always present in the other (gradient-based) systems:

- The optimization function for the evolution program need not be continuous. At the same time some optimization packages will not accept such functions at all.
- Some optimization packages are all-or-nothing propositions: the user has to wait until the program completes. Sometimes it is not possible to get partial (or approximate) results at some early stages. Evolution programs give the users additional flexibility, since the user can monitor the ‘state of the search’ during the run time and make appropriate decisions. In particular, the user can specify the computation time (s)he is willing to pay for (longer time provides better precision in the answer).
- The computational complexity of evolution programs grows at a linear rate; most of other search methods are very sensitive to the length of the optimization horizon. As usual, we can easily improve the performance of the system using parallel implementations; often this is difficult for other optimization methods”

[Ref. 12, pg. 117].

It is for this reason that a GA will be used to determine the 'best' RTOS scheduling policy for the given problem specification.

All real-time operating systems provide some means of 'tuning' the characteristics of the scheduling policy to accurately meet the application requirements. The main objective of this thesis is to show the applicability of using a GA to determine these parameters for an arbitrary application. In so doing, a tool set will be developed which completely automates this tuning process. In addition, the RTOS parameters that are considered are broad enough to allow the results to be used for specifying and/or choosing an RTOS for the actual implementation of the real-time system.

The domain of real-time applications which is of particular interest is that of embedded systems. In the embedded systems domain, real-time multitasking problems are specified by the usual real-time properties, such as a series of timing constraints and time deadlines. In addition, embedded systems are further confined to using processing resources that are both practical and available in the particular application environment. These additional constraints are probably the most significant attributes of embedded real-time systems which must be considered during the design of the system. These constraints arise from the very nature of embedded systems and take the form of size, weight, power and cost restrictions.

Therefore, in addition to evaluating the ability of various RTOS scheduling policies to meet the real-time timing constraints and time deadlines (both hard and soft), the effectiveness of the scheduling policy is evaluated with regard to

CPU utilization. One of the most common means of reducing the size, weight, power and cost of an embedded system is to reduce the required CPU processing power. As a result of reducing the CPU processing power, the clock frequency and memory speed can often be reduced as well. This is the area of the system design where size, weight, power and cost savings are achieved. Therefore, the 'best' RTOS scheduling policy (which is application-specific) is the one which maximizes the following criteria:

- Percentage of timing constraints met.
- Weighted percentage of deadlines met (recognizing that for applications with hard deadlines, anything less than 100% is unacceptable).
- Percentage of CPU processing bandwidth available (i.e., surplus processing power).

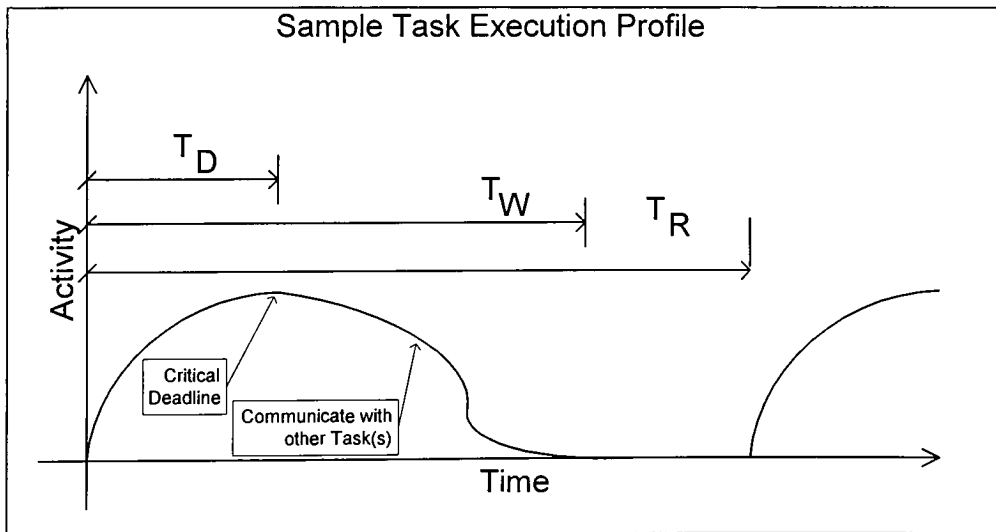
Real-time operating systems typically provide relatively few parameters which can be 'tuned' to a particular application. This thesis evaluates the effectiveness of the following RTOS parameters (this list is intended to include all configuration parameters likely to be found in most available RTOSs):

- Tasking Model - Both cooperative and preemptive multitasking models are used.
- Timeslice - The processing timeslice ranges from 50  $\mu$ Sec to 65 mSec (only applicable for preemptive multitasking).

- Priority Inheritance - Scheduling policies with and without priority inheritance are used (enabling priority inheritance guarantees that unbounded priority inversion cannot occur).
- Priority Allocation - During run-time, the task priorities are either fixed at the initial assigned value (static priority allocation) or changed through the use of a rotation scheme (dynamic priority allocation).
- Initial Priority Assignment - The priority of each task is initially assigned using one of the following algorithms:
  - ◊ Uniform (constant) Assignment - All tasks are assigned the same priority regardless of execution profile.
  - ◊ Ad-hoc (random) Assignment - Tasks are assigned random priorities, uniformly distributed within the allowable range, without regard to execution profile.
  - ◊ Rate Monotonic Assignment - Tasks are assigned priorities based on execution period. Tasks with high execution rate (low period) are assigned high priorities; lower execution rates (higher periods) are assigned proportionally lower priorities. Execution period is defined as the amount of time which elapses between consecutive executions of a task (periodic execution is assumed). See *Figure 1 - Task Execution Profile*,  $T_R$ .
  - ◊ Deadline Monotonic Assignment - Tasks are assigned priorities based on execution deadlines. Tasks with short deadlines are

assigned high priorities; longer deadlines are assigned lower priorities. Execution deadlines are defined as the amount of time which can elapse between when a task begins executing and when it must reach a critical point in its execution. See *Figure 1 - Task Execution Profile,  $T_D$* .

- ◇ Workload Monotonic (greedy) Assignment - Tasks are assigned priorities based on the amount of work that the task must perform. Tasks with high workloads are assigned high priorities; lower workloads are assigned proportionally lower priorities. Task workload is defined as the amount of time which elapses between when a task begins executing and when it completes executing. See *Figure 1 - Task Execution Profile,  $T_W$* .



*Figure 1 - Task Execution Profile*

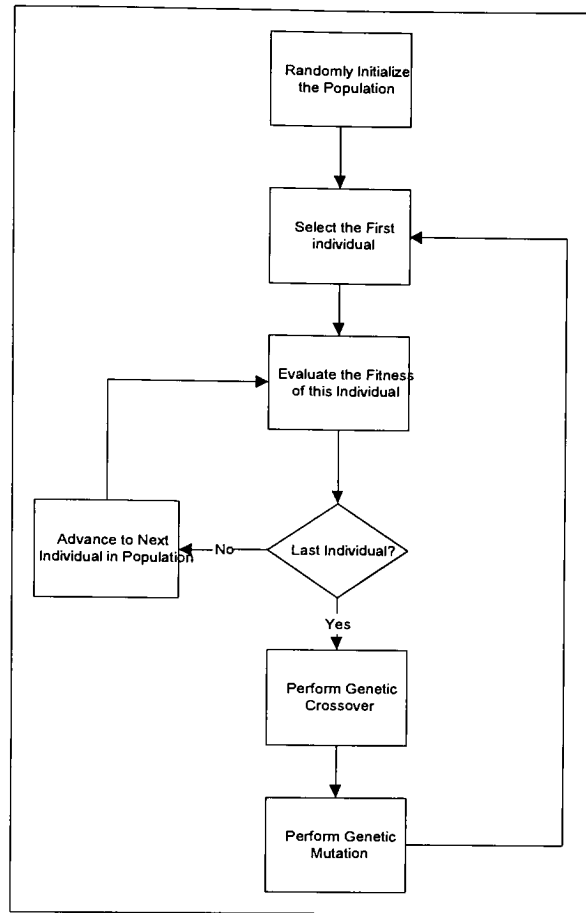
The Genetic Algorithm used to perform this RTOS evaluation is a generalized algorithm with application-specific genetic operators (population

crossover and mutation), see *Figure 2 - Genetic Algorithm*. It has been shown that genetic operators are most capable of performing the required search when they are designed with prior knowledge of the problem domain.

"... although the gains are not dramatic, RAR<sub>2</sub> out-performs uniform crossover with a simple penalty function on a family of problems in which the solution is a fixed-size set, using fitness functions which vary from highly epistatic to non-epistatic. This therefore provides some evidence that forma [sic] analysis is capable of leading to effective genetic search in problem domains not well suited to conventional binary representations."

[Ref. 17, pg. 27]

Therefore, the approach taken is to fully define the problem domain (i.e., RTOS configuration parameters) and then design the genetic operators so as to properly evolve a solution to this specific problem domain.



*Figure 2 - Genetic Algorithm*

It is obvious from *Figure 2* above that the GA itself is extremely simple. The power of the algorithm comes from the fact that as it progresses, it does two distinct things: 1) it continuously improves and, 2) it explores solutions which may provide additional improvements. Both of these operations are encompassed in the genetic operators of population crossover and mutation. These operators manipulate the genes of the individuals to produce the continuously improving and experimentation properties of the GA. Since the genes of the individuals are used to determine how it behaves (i.e., how well it solves the problem), and the genetic operators manipulate those genes, the genetic operators must be intricately tied to the representation of the genes. This

results in genetic operators that are specific to the problem domain. Significant research has been done, attempting to determine 'universal genetic operators' based on 'universal gene representations'. Unfortunately, as the following quote summarizes, these attempts have not been very successful:

"The EC (Evolutionary Computation) community differs widely on opinions and strategies for selecting appropriate representations, ranging from universal binary encodings to problem-specific encodings for TSP (Traveling Salesman Problem) problems and real-valued parameter optimization problems. The tradeoffs are fairly obvious in that universal encodings have a much broader range of applicability, but are frequently outperformed by problem-specific representations which require extra effort to implement but exploit additional knowledge about a particular problem class."

[Ref. 7, pg. 619]

It is for these reasons that the GA developed for this thesis will use application specific genetic operators.



## 2. Background

In order to provide a firm foundation from which to establish this thesis, the concepts of real-time applications and genetic algorithms must be fully defined. This section will provide a generalized, detailed description of these very important concepts by using two simple, but concrete, examples.

### *2.1 Real-Time Applications*

It is important to first discuss the characteristics of real-time problems in general, after which, the problems associated with determining the appropriate configuration parameters of an RTOS will become apparent. The following example will be used to illustrate the general characteristics of real-time applications.

This example will consider the problem of a Reaction Rate Controller (RRC) of a chemical processing plant [Ref. 1]. The RRC is responsible for managing the reaction rate for a single chemical process. For this simplified example, the reaction rate is a function of the relative proportions of only two reagents and the temperature and pressure within the reaction chamber.

The user interface to the RRC is an operator console; the RRC must accept the desired reaction rate from the operator and provide reaction status to the operator. This operator interaction is relatively slow due to the fact that human operators respond slowly and require some time to assimilate information. It is assumed that the operator cannot request reaction rate changes more frequently than once every two seconds and that the operator cannot recognize status

changes more frequently than four times per second. In order to ensure that the operator is always informed about the status of the reaction, the RRC must update the operator display no less than once per second.

The RRC must also accommodate the fact that the reaction rate within the reaction chamber can change within 100 milliseconds. In order to maintain a safe reaction environment, the RRC must respond to that change by re-positioning the reagent valves within an additional 100 milliseconds. This means that in the worst case, the reagent valves may be re-positioned as frequently as every 200 milliseconds. The RRC must also ensure that the reaction pressure and temperature are sampled at a sufficiently high rate to ensure that all changes are detected. Again, since the reaction rate within the reaction chamber can change within 100 milliseconds, the Nyquist Theorem states that the sampling period for the pressure and temperature sensors must be no more than 50 milliseconds.

The system software architecture for the RRC is shown in *Figure 3 - Reaction Rate Controller (RRC)*, which shows data and control flow within the system, timing parameters and software tasks.

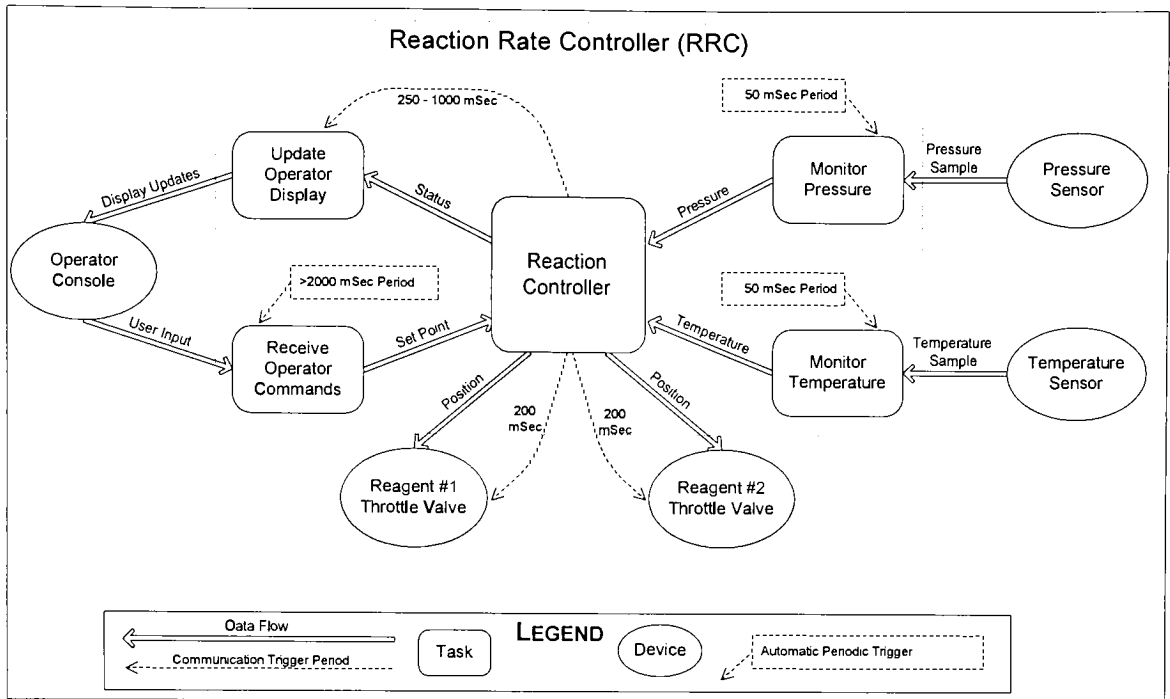


Figure 3 - Reaction Rate Controller (RRC)

The physical input/output devices connected to the RRC, the information flow to/from the RRC and a description of the timing requirements imposed on the RRC are summarized in *Table 1* below.

ID #	I/O Device	Data Flow	Direction	Timing Requirements
1	Operator Console	Desired Reaction Rate	To RRC	Non-Periodic: >2000 mSec
2	Operator Console	Actual Reaction Rate & Reaction Status	From RRC	Periodic: 250 mSec (Design Goal), <1000 mSec (Requirement)
3	Reagent #1 Valve	Valve Position	From RRC	<100 mSec after Reaction Change
4	Reagent #2 Valve	Valve Position	From RRC	<100 mSec after Reaction Change
5	Pressure Sensor	Pressure Sample	To RRC	Periodic: 50 mSec
6	Temperature Sensor	Temperature Sample	To RRC	Periodic: 50 mSec

Table 1 - RRC Requirements

From this example it is readily apparent that real-time applications are most notably characterized by deadlines. That is, real-time applications are specified by a series of deadlines that the software (in this example, the RRC) must meet. These deadlines typically arise from the very nature of the

application and the processing that the application must perform. Careful analysis of the problem can expose all of the deadlines that a system must meet. In addition to identifying the deadlines, it is also important to characterize the nature of each deadline. For example, some deadlines are more 'important' to the proper operation of the system than others. In the RRC example above, it is obvious from *Table 1 - RRC Requirements* that requirement #2 (Display Update) is 'least' important because a range of allowable performance values is specified. Intuitively that makes sense as well because failure to update the operator display as frequently as four times per second will result in a graceful performance degradation. On the other hand, failure to properly control the reagent valves may result in an unsafe reaction; therefore the requirements associated with actually controlling the reaction are 'more' important to the correctness of the system.

The different relative importance of real-time deadlines is expressed as to whether a given deadline is 'hard' or 'soft'. Hard deadlines are those which must always be met in order for the system to be said to operate correctly. Soft deadlines are those which can be missed (occasionally) without severely impacting the correctness of the system. In other words, if a system fails to meet a hard deadline, a catastrophic failure usually results. Failure to meet a soft deadline usually results in graceful degradation of some aspect of system performance. In the above example, failure to meet a hard deadline associated with controlling the reaction rate may result in an unsafe reaction which could result in loss of life or property. Failure to meet the soft deadline associated with

updating the operator display will result in the inability of the operator to adjust the reaction rate to the desired level (if the display is not updated, the operator does not know that the reaction rate is something other than the desired rate). Therefore, in the above example, all deadlines are hard except for update of the operator display. In essence, the system can be said to be operating properly even if the operator display is updated less frequently than four timer per second.

In large, complex systems, the concept of strictly hard or soft deadlines is usually too restrictive to adequately describe all of the deadlines present within the system. Obviously, some deadlines will still be strictly hard or soft, but some soft deadlines may be more important than other soft deadlines. It is typically necessary to extend hard and soft deadlines to include a 'degree of hardness' so that all of the subtle aspects of required system performance can be captured. This allows the system to be evaluated based on how well the hard deadlines are met ('Is the system operating correctly/safely?'), as well as how well the varying degrees of hardness deadlines are met ('Is the performance of the system adequate?'). Another method of specifying the relative importance of a soft deadline is to use the benefit accrual model of timeliness [Ref. 2, pg. 30]. In this model, all deadlines are modeled as a function of time,  $f(t)$ , which corresponds to the benefit achieved to the system by completing the deadline at time  $t$ . In this way the complete spectrum of hard and soft deadlines can be modeled consistently and thoroughly.

The problem associated with determining the appropriate configuration parameters of an RTOS for a given real-time application is extremely difficult.

One reason for this difficulty is that there is often very little to guide a designer in choosing the 'best' RTOS parameters. Even though it may be possible to fully specify the real-time application in terms of its deadlines, that alone in no way guarantees that the deadlines will actually be met. The RTOS configuration will affect the system performance which will in turn determine to what extent the system performs correctly (hard deadlines) or to what degree the system performance is degraded from the design goals (soft deadlines).

### **2.1.1 Formal Methods for Real-Time Analysis**

Recently, there has been significant research into defining formal methods for analyzing the correctness of real-time systems [Ref. 9]. The appeal of this approach is that by using some of these formal methods, the correctness of a real-time system design can often be proven analytically. Unfortunately, for large complex systems, the applicability of these methods can be limited.

The techniques that have been developed for analyzing real-time systems generally fall into two distinct categories: model-based reasoning and proof-based reasoning [Ref. 9, pg. 14]. Each of these has associated advantages and disadvantages, and each has been used successfully to model real-time systems. The remainder of this section outlines the major accomplishments in defining formal methods for the analysis of real-time systems.

#### ***2.1.1.1 Model-Based Analysis***

For model-based analysis, the real-time system is typically specified using a graphical model tool set. The main advantage to this approach is that the model can be created and maintained by individuals with no formal mathematical

background, and the analysis is automatically performed by the tool set. In addition, the graphical model provides a description of the system that can be used to document the real-time design. One disadvantage of this approach is that it does not scale well to large complex systems. The problem of creating and maintaining a single large model that fully describes the entire system can easily result in errors and areas of model inadequacy. Another problem is that as the model grows, the analysis process becomes increasingly more time consuming. This is due to the fact that the complexity of the analysis grows as  $O(2^N)$  where  $N$  is the number of states in the real-time system.

One promising method of model-based analysis that has been introduced is based on the Modechart specification language [Ref. 14]. This specific method provides a set of tools which is used to create the system specification, simulate the model to exhibit individual behaviors of the system, and verify global properties of the system. This method is promising because it provide significant automation of the analysis using the tool set provided. Unfortunately this, like all state-based methods suffers from the state explosion problem. That is, as the real-time system being modeled more closely parallels a real-world problem, the number of states required to accurately represent the system becomes unmanageably large. Ultimately what often happens with such a system is that, over time, the model loses synchronization with the actual system being designed and becomes obsolete. Another potential problem with state-based models is that the computations that must be performed to verify the design grow of order  $2^N$ . Again, the number of states creates an often unacceptable computation time

which is not proportional to the increase in design complexity. Even a marginal increase in complexity of the system can result in a tremendous increase in the complexity of the computations necessary to ‘prove’ the validity of the design. For small-scale real-time systems this method may be feasible but it is unlikely to be applicable for large, complex systems.

### ***2.1.1.2 Proof-Based Analysis***

For proof-based analysis, standard logical deduction methods are used to formally prove that one specification is equivalent to another specification. With the system design and system requirements being used as the starting and ending points of the proof, this method can prove that a specific real-time design specification is equivalent to the corresponding system requirements specification. The main advantage to this approach is that a formal proof of the design meeting the requirements is achieved. The disadvantages are that producing the proof requires formal mathematical training, and since the proof must be done by hand, changes in the design and/or requirements can necessitate the complete recalculation of the proof.

One promising method of proof-based analysis that has been introduced is based on a timed process algebra which supports time consuming actions, instantaneous events, and the concept of prioritized interactions [Ref. 11]. In order to utilize this method, the real-time system and the requirements must be specified in terms of the ACSR specification language [Ref. 11, pg. 169-179]. This may not seem like a severe limitation but it often is. First, expertise in the ACSR specification language is necessary, although this is not significant it is



knowledge that is not often present within the real-time development community. Secondly, and more importantly, using this analysis method requires that not only the system being designed but also the system requirements be specified in the ACSR specification language. System requirements are very infrequently defined to a degree that would allow a formal specification to be developed. More often than not the system requirements are being developed at the same time that the system design is being developed. Again, this is not to say that this method is not without merit, it is just that for real-world problems where the requirements of the system are dynamic this method may not be feasible.

## **2.2 Genetic Algorithms**

It is important to first discuss the characteristics of genetic algorithms (GAs) in general, after which, the particular application of a GA to determining appropriate configuration parameters of an RTOS can be investigated. One way to understand GAs is to view them as algorithms which search the space of possible solutions to a problem for the optimal solution. In this regard, GAs are very similar to other parametric search algorithms (hill climbing, simulated annealing, etc.). The most important difference between GAs and other search methods is that GAs attempt to model the natural world laws of ‘random perturbations’ and ‘survival of the fittest’. This property is unique to GAs.

The basic architecture of a GA is one of a collection of individuals, where each individual is a candidate solution to the problem being solved. In turn, the basic architecture of each individual is one of a collection of ‘genes’, where each gene represents one component of the parameterized solution to the problem.

The GA uses this ‘population’ of individuals and ‘gene pool’ of genes together with the natural laws of ‘random perturbations’ and ‘survival of the fittest’ to allow the ‘best’ solution to evolve over many generations, see *Figure 4 - General Genetic Algorithm*. From this discussion it is apparent that the GA is parameterized as follows:

- Population Size,  $N$  - The number of individuals that will be used to comprise the population. This value is chosen such that statistical significance within the population is maintained and such that the initial (and potential) gene pool is sufficiently diverse.
- Mutation Rate,  $P_m$  - The rate at which ‘random perturbations’ will be introduced into the gene pool. This value is chosen in such a way as to ensure that any useful genetic material which is lost due to crossover, is re-introduced into the gene pool and to allow the GA to explore areas of the search space that have not yet been explored.
- Crossover Rate,  $P_c$  - The rate at which individuals participate in genetic recombination, i.e., the rate at which genes are recombined from individuals of the present generation to form individuals of the next generation. This is the means that ‘fit’ genetic material ‘survives’ from generation to generation. The value is chosen in a manner that allows for a reasonable rate of convergence, a trade-off between execution time and probability of false convergence must be made.

It is obvious that this approach does not constitute an exact or formal method. Instead, as the example that follows will illustrate, the GA seems to

meander about in its search for the ‘best’ solution, never really knowing where the search will lead but amazingly enough almost always able to find a reasonably appropriate solution.

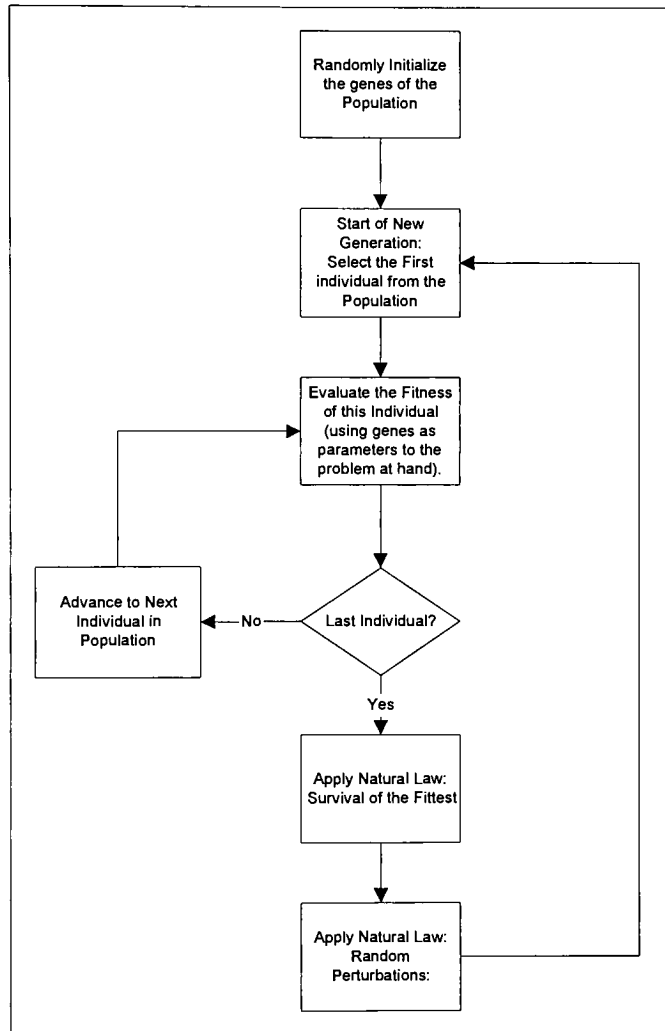


Figure 4 - General Genetic Algorithm

Note that the answer found by the GA is not necessarily optimal. This arises from the fact that the GA is never really finished, in theory the solutions would evolve over an infinite number of generations (notice that there is no exit condition specified in the *Figure 4*). This means that in practice, one of the more difficult problems associated with using a GA is determining when to stop the

search. This decision is often based on statistical measurements of the population and/or gene pool, as stated below.

“There are two basic categories of termination conditions, which use the characteristic of the search for making termination decisions. One category is based on the chromosome structure (genotype); the other - on the meaning of a particular chromosome (phenotype). Termination conditions from the first category measure the convergence of the population by checking the number of converged alleles, where an allele is considered converged if some predetermined percentage of the population have the same (or similar - for non-binary representations) value in this allele. If the number of converged alleles exceeds some percentage of total alleles, the search is terminated. Termination conditions from the second category measure the progress made by the algorithm in a predefined number of generations; if such progress is smaller than some epsilon (which is given as a parameter of the method), the search is terminated.”

[Ref. 12, pg. 65].

The natural laws of ‘random perturbations’ and ‘survival of the fittest’ were first articulated by Charles Darwin in The Origin Of Species. Even though these concepts were focused entirely on biological organisms and their interactions, it is in this rich background that the field of genetic algorithms is based.

As shown below, Darwin uses the term ‘individual differences’ to express the concept of random perturbations as it relates to biological organisms; the

exact same principles are used when applying genetic algorithms to computational problems, see section 3.2 *Genetic Algorithm*.

“The many slight differences which appear in the offspring from the same parents, or which it may be presumed have thus arisen, from being observed in the individuals of the same species inhabiting the same confined locality, may be called individual differences. No one supposes that all the individuals of the same species are cast in the same mould. These individual differences are of the highest importance for us, for they are often inherited, as must be familiar to every one; and they thus afford materials for natural selection to act on and accumulate, in the same manner as man accumulates in any given direction individual differences in his domesticated productions.”

[Ref. 5, pg. 59]

As Darwin indicates below, individual biological organisms within a population compete with each other for the right to survive. Together with random perturbations, natural selection creates a continuously improving population. This is true even though, for any given generation, there will exist individuals that are less fit for survival.

“Can it, then, be though improbable, seeing that variations useful to man have undoubtedly occurred, that other variations useful in some way to each being in the great and complex battle of life, should occur in the course of many successive generations. If such do occur, can we doubt (remembering that many more individuals are born than can possibly survive) that individuals having any

advantage, however slight, over others, would have the best chance of surviving and of procreating their kind? On the other hand, we may feel sure that any variation in the least degree injurious would be rigidly destroyed. This preservation of favourable individual differences and variations, and the destruction of those which are injurious, I have called Natural Selection, or the Survival of the Fittest. Variations neither useful nor injurious would not be affected by natural selection, and would be left either a fluctuating element, as perhaps we see in certain polymorphic species, or would ultimately become fixed, owing to the nature of the organism and the nature of the conditions.”

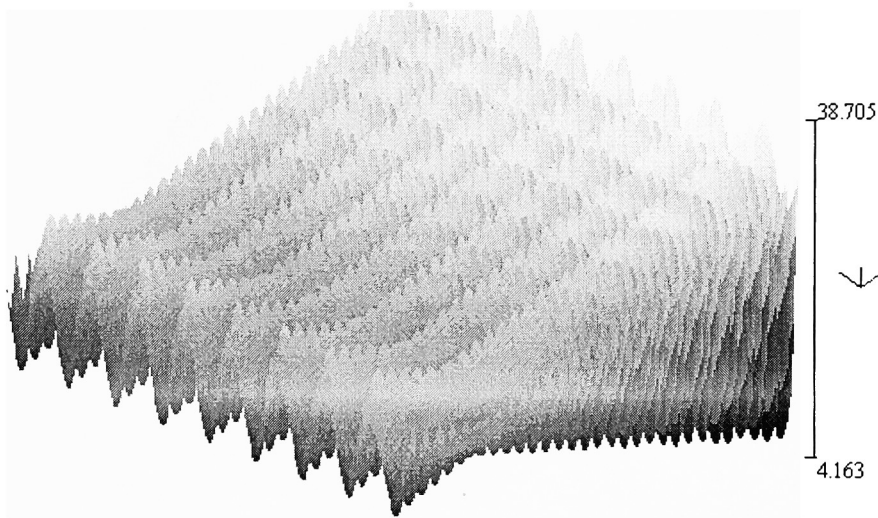
[Ref. 5, pg. 87-88]

The following simple example will be used to illustrate the general characteristics of genetic algorithms by determining the maximal value of a complex trigonometric function. Consider the following function of two real-valued variables:

$$\begin{aligned} f(x, y) &= 21.5 + x \cdot \sin(4 \cdot \pi \cdot x) + y \cdot \sin(10 \cdot \pi \cdot y) \\ -3.0 &\leq x \leq 12.1 \\ 4.1 &\leq y \leq 5.8 \end{aligned}$$

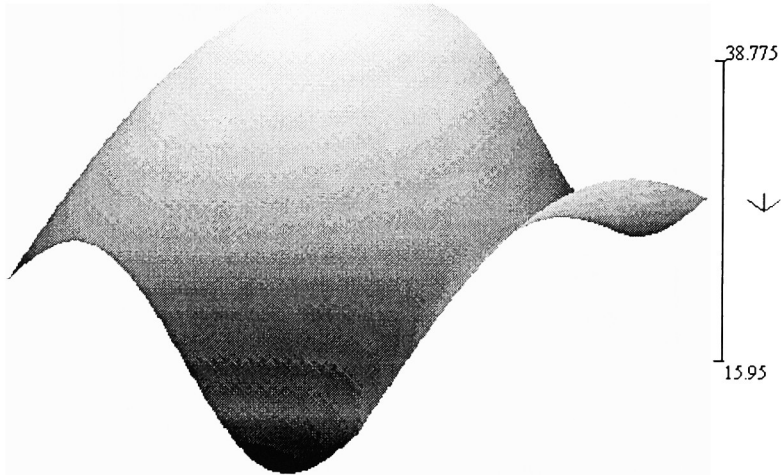
*Equation 1 - Sample Optimization Function*

It is not at all obvious what values of  $x$  and  $y$  result in the maximal value of this function. The plot below, *Figure 5 - Sample Optimization Function Plot*, shows the function for the given range.



*Figure 5 - Sample Optimization Function Plot*

Because the function has so many local maxima, most optimization algorithms would ‘find’ a local maximum instead of the global maximum unless an extremely close initial value was used, GAs do not possess such a limitation. After localizing the exhaustive search to the region shown in *Equation 2 - Sample Optimization Function (Localized)*, the maximum value is determined. The graph is shown in *Figure 6 - Sample Optimization Function Plot (Localized)* shows the maximum value of the function is  $f(11.626, 5.65) = 38.775$ .



*Figure 6 - Sample Optimization Function Plot (Localized)*

$$\begin{aligned}
 f(x, y) &= 21.5 + x \cdot \sin(4 \cdot \pi \cdot x) + y \cdot \sin(10 \cdot \pi \cdot y) \\
 11.5 &\leq x \leq 11.7 \\
 5.5 &\leq y \leq 5.7
 \end{aligned}$$

*Equation 2 - Sample Optimization Function (Localized)*

The following table shows the results of a single run of `GMATH-GA.EXE`, the results show the percent error after 903 generations. Notice that even without any knowledge of the optimal solution, and searching the entire range (see *Equation 1*), the GA performed remarkably well at determining the maximum value of the function. This function optimization example will be revisited in the remainder of this section as the various aspects of genetic algorithms are detailed.



$f(x,y)$	$x$	$y$
38.668 (0.276%)	11.615 (0.095%)	5.650 (0.000%)
38.664 (0.286%)	11.615 (0.095%)	5.651 (0.018%)
38.662 (0.291%)	11.615 (0.095%)	5.652 (0.354%)
38.657 (0.304%)	11.615 (0.095%)	5.648 (0.354%)
38.603 (0.443%)	11.615 (0.095%)	5.645 (0.088%)

*Table 2 - GA Solutions to Sample Function*

### 2.2.1 Genetic Operators

The most important aspect of the genetic algorithm is the internal representation of the genes. There is nothing which is more intrinsic to a specific GA for a problem than the gene representation and the genetic operators which manipulate those genes. There are two genetic operators which essentially comprise the genetic algorithm: mutation(...) and crossover(...). These operators are responsible for the behavior of the GA, its ability to converge to a solution, and the rate at which the convergence occurs.

It is important to note that both of these operators must be designed with complete knowledge of the precise genetic representation which is being used. There are two different representation methods which are typically used to create a GA: pure binary genes and representational genes. There is significant debate within the GA research community as to whether one method is superior to the other, most indications are that pure binary genetic representations may be outperformed by representations which more closely match the problem domain [Ref. 17, pg. 27].

The following sections will describe the mutation(...) and crossover(...) operators for both type of genetic representation.

### **2.2.1.1 Mutation**

The mutation(...) operator attempts to model the natural law of ‘random perturbations’. Quite simply, mutation is the random altering of genetic material of some members of the population. One of the parameters of a GA is the mutation rate (or probability of mutation). This parameter controls how frequently (or with what probability) genes will be mutated.

Determining the ‘correct’ mutation rate for a given genetic algorithm optimization problem is itself an optimization problem, this is also true for determining other parameter values for a GA. Since it is not feasible to construct a GA which determines the parameters for another GA, although this has, been done [Ref. 6, pg. 92], typically, general principles are used to select approximate GA parameter values.

One of the most important measurements of the GA parameters is the expected allele coverage of the population. The mutation rate and the population size directly effect the expected allele coverage, 99% expected allele coverage is typically used [Ref. 19]. The way that the expected allele coverage is calculated is different for pure binary genes than for more representational genes. *Equation 3* and *Equation 4* provide guidelines for determining the mutation rate or population size for binary and representational genes respectively [Ref. 19].

For pure binary genes the genetic material is uniform, therefore expected allele coverage is simply a function of the population size and the number of

choices for each bit (i.e., 2), see *Equation 3 - Expected Allele Coverage (Binary Representations)* below [Ref. 19, pg. 33].

$$E[ac] = 1 - (1/2)^N$$

Where:  $N$  is the population size

*Equation 3 - Expected Allele Coverage (Binary Representations)*

For more representational genes the genetic material is not uniform, therefore a summation is required to accumulate the weighted expected allele coverage, see *Equation 4 - Expected Allele Coverage (Non-Binary Representations)* below [Ref. 19, pg. 33].

The probability that exactly  $m$  distinct symbols are used for any given gene within the entire population is given by:

$$p_m(N, K) = \frac{1}{K^N} \cdot \binom{K}{m} \cdot \sum_{v=0}^{K-m-1} (-1)^v \cdot \binom{K-m}{v} \cdot (K-m-v)^N$$

Where:  $K$  is the number of distinct symbols used in the encoding and  $N$  is the population size.

Since all gene positions are disjoint with respect to alleles, the expected allele coverage is given by:

$$E[ac] = \frac{1}{K} \cdot \sum_{m=1}^K m \cdot p_m(N, K)$$

*Equation 4 - Expected Allele Coverage (Non-Binary Representations)*

For pure binary representations, carrying out the mutation operation is very straightforward. Since each individual's genotype is represented as a single bit string of length  $m$ , where individual genes that comprise the genotype are assigned sub-regions of the bit string, and the population is composed of  $n$  individuals, the entire gene pool is composed of  $n \cdot m$  bits. The mutation

operation is performed once every generation where each bit is mutated (i.e., inverted) with probability  $P_m$ .

For representational genes, carrying out the mutation operation is slightly more complex. Since it is possible that simply mutating single bits will result in illegal (i.e. meaningless) genes, the mutation operation must always produce a gene which is valid within the problem domain. For this reason the mutation operator must be closely tied to the actual problem being solved. Again, the mutation operation is performed once every generation but in this case each gene is mutated (the exact meaning of which is application specific) with probability  $P_m$ .

In the example above, the genes are two real-valued numbers in the ranges specified in *Equation 1 - Sample Optimization Function*. Since the genes are not binary encodings, mutation for this example is defined as replacing the gene to be mutated with a randomly selected value from the corresponding range. This is typical of a representational genotype. The mutation rate,  $P_m$ , is determined by choosing a population size,  $N$ , of 175 individuals, setting the desired expected allele coverage to 97% and evaluating *Equation 4*, above. The resulting mutation rate, for this example, is  $P_m = 3\%$ , resulting in accuracy of  $1/50^{\text{th}}$  of the respective ranges of each variable.

### **2.2.1.2 Crossover**

The crossover(...) operator attempts to model the natural law of ‘survival of the fittest’. Genetic crossover allows individuals to pass their genetic material on to individuals of the next generation. To do this, those individuals that are

most fit are chosen to participate in genetic crossover. When crossover(...) is applied to the population, two relatively fit individuals are chosen at random and their genetic material is combined in some way to produce an offspring. The most common type of crossover is ‘single-point crossover’. Single-point crossover occurs when a randomly selected gene is used to split the genotypes of the two parents, the offspring inherits one half of the genotype from each parent. There are important consequences of this operation, as described below.

“One important feature of one-point crossover you should be aware of is that it can produce children that are radically different from their parents. ... Another important feature is that one-point crossover will not introduce differences for a bit in a position where both parents have the same value. ... An extreme instance occurs when both parents are identical. In such cases, crossover can introduce no diversity in the children.”

[Ref. 6, pg. 17]

There is no theoretical guideline for determining the proper crossover probability,  $P_c$ , for a GA, typical values range from 25% to 75%. Values used for GAs in this thesis will be determined empirically but will be from the typical range.

In the example above, single-point crossover with a crossover probability of  $P_c=25\%$  was used.

### 3. Evolution of Solutions to Real-Time Problems

The most promising applications of genetic algorithms arise in the area of engineering design. This is true because many design problems are very complex (often NP-Complete or at least NP-Hard) and once the engineering and business trade-offs have been made, these problems can be expressed as an optimization problem. GAs are uniquely qualified for these exact type of optimization problems where there are far too many system constraints for an engineer (or small group of engineers) to keep in mind during the design phase. If sufficiently large computers are available there is really no limit to the number of constraints that can be managed by a GA. In addition, GAs are capable of solving problems without actually requiring that the specific design problem be solved (or solvable) at all, as the following summarizes:

“Assume, then, that (for some reason) we have to (or like to) build a new system to solve a nontrivial optimization problem. ... Then we have to make a choice: either we can try to construct an evolution program or we can approach the problem using some traditional (heuristic) methods. It is interesting to note that in a traditional approach it usually takes three steps to solve an optimization problem:

1. Understand the problem
2. Solve the problem
3. Implement the algorithm found in the previous step

In the traditional approach, a programmer should solve the problem - only then may a correct program be produced. However, very often an algorithmic solution of a problem is not possible, or at least is very hard. On top of that, for some applications, it is not important to find the optimal solution - any solution with a reasonable margin of error (relative distance from the optimum value) will do. ... An evolution programming approach usually eliminates the second, most difficult step. Just after we understand the problem, we can move to the implementational issues. The major task of a programmer in constructing an evolution program is a selection of appropriate data structures as well as genetic operators to operate on them (the rest is left for the evolution process).”

[Ref. 12, pg. 302]

In addition to the all too often cited example of VLSI circuit design, there are many other examples of systems that have been designed using, at least in part, a genetic algorithm. One such example is the parametric design of aircraft:

“The problem addressed here is parametric design using a design concept typical of modern fighter aircraft. The problem is representative of real-world aircraft conceptual design problems. It is similar to practical problems that are currently being solved using methods described here. The aircraft takes off, accelerates subsonically and then perhaps supersonically, transits to an engagement area, conducts combat operations, returns to base, loiters, and lands. The parameters describing the mission include altitude, duration and velocity of those activities, range, and acceleration and maneuverability requirements for combat.

Representative parameters of the aircraft's geometry and configuration include fuselage width, height and length; wing planform and thickness characteristics; vertical and horizontal tail surface dimensions; and engine size (that is, aircraft thrust-to-weight ratio). Parameters of the design concept include unscaled engine weight and dimensions, and weights of aircraft subsystems and the weapons complement. The total number of parameters for this problem is fifty.”

[Ref. 3, pg. 113]

The engineering design problem that is being considered in this thesis is the system software architecture design for an embedded, real-time application. This design problem can be viewed as an optimization problem. The ‘function’ to be optimized is the applicability of a real-time operating system (RTOS) configuration to a given problem specification. The parameters of the ‘function’ are the configuration items for the RTOS. It is reasonable to hypothesize that there is no optimal RTOS configuration for all real-time applications, instead the best RTOS configuration for any given application will, in some unknown manner, be a function of the application itself. Therefore, the tools developed must be capable of accepting a user specified real-time application. The tools will then generate a list of the ‘best’ RTOS configurations for that application. A genetic algorithm will be used to determine these RTOS configurations. In general, a genetic algorithm can be used to find the optimal solution to a given problem but only if given infinite time. Since infinite time is not feasible, the solutions provided by the tools developed in this thesis will produce sub-optimal



solutions. The following sections describe the RTOS, genetic algorithm and real-time application specification aspects of the tool set.

### **3.1 Real-Time Operating System (*TASKING*)**

*TASKING* provides all of the RTOS services necessary for real-time application development. It is important to note that the services provided merely allow for the development of real-time applications -- they do not ensure that a system design will actually perform in real-time. The services provided are easily segregated into the following categories: task management, software events, binary semaphores, counting ("Dijkstra") semaphores, message passing, condition variables and interactive I/O handling.

The task management services primarily allow tasks to be dynamically created and destroyed. In addition, services are provided so that tasks can be suspended, resumed or change their priority. These services are sufficient to allow for the design of a multitasking application; however, additional services are necessary to allow the tasks to communicate and interact. The remaining categories of services provide these necessary capabilities.

Inter-task communication and interaction are achieved by the use of one (or more) of the following constructs: software events, binary semaphores, counting ("Dijkstra") semaphores, message passing or condition variables. These constructs are actually a superset of those that are required for inter-task communication. It has been shown that at the most basic level, all of these constructs can be considered to be a special case of a simple semaphore.

“In the previous sections we have studied four different interprocess communication primitives. Over the years, each one has accumulated supporters who maintain that their favorite way is the best way. The truth of the matter is that all these methods are essentially semantically equivalent (at least as far as single CPU systems are concerned). Using any of them, you can build the other ones.”

[Ref. 18, pg. 52]

Nonetheless, in order to ease application development, they are all provided by `TASKING`. Note that when these constructs are created they are just like any other variable in that their value is undefined. Before the construct is used it should be initialized to a value appropriate for the given construct. In addition, task interaction with the application user is performed by a set of blocked I/O handling routines.

Software events are one of the more simple constructs; events can take on only one of two possible values: signaled or unsignaled. In addition to this restriction, events can retain only a single signal, i.e., once an event has been signaled, all subsequent signals to that event are superfluous (the event remains signaled) but are not considered errors. Tasks may communicate simple information with events but most often events are used for task synchronization. Quite simply one task would wait for an event to be signaled by another task which would indicate that some temporal situation was satisfied (e.g., a data buffer had just become non-empty or non-full, etc.). At that point the tasks would be synchronized.

Whereas software events are typically used to task synchronization, binary semaphores are typically used to ensure mutually exclusive access to a resource which is shared between two (or more) tasks. Binary semaphores, like software events, can only take on two possible values: 0 or 1. In addition, it is logically an error to signal an already signaled binary semaphore; `TASKING` will enforce this restriction by detecting and reporting this as an error.

"Dijkstra", or counting, semaphores are similar to binary semaphores except that they can retain multiple signals because `TASKING` maintains information so that it knows how many signals are stored in the semaphore at all times. Whenever a given semaphore is signaled the count is incremented. Tasks that wait on the semaphore are only suspended when the semaphore contains no signals (i.e., the semaphore value is zero). This allows "Dijkstra" semaphores to provide control over access to a resource that has more than one instance. That is  $n$  tasks can access  $m$  similar resources, where  $n \geq m$ .

`TASKING` also supports passing arbitrary, application-specific messages between tasks; this is accomplished when the source task allocates memory for the message and sends a pointer to the message to the destination task. The receiving task is responsible for de-allocating the memory for the message after it has been consumed. There are no restrictions on the size, structure or number of messages that can be used within the application. `TASKING` manages mailboxes for the destination of the messages and all messages are delivered in a non-prioritized, first-in-first-out manner. Extreme care must be taken to ensure that

these message queues do not cause unbounded priority inversion; note that this form of priority inversion cannot be detected or prevented by `TASKING`.

The last construct is condition variables. Condition variables do not have any memory associated with them; if a task is already waiting on a condition variable when it is signaled it will become unblocked, but if a condition variable is signaled and no task is waiting on it, then the signal is lost. Condition variables are typically used in the implementation of Hoare Monitors or for simple forms of task synchronization.

The other service that `TASKING` can provide is blocked I/O. Both keyboard and MS-Mouse I/O can be performed in a completely blocking manner where the task is blocked (suspended) until input is available. The only restriction is that only one task each can perform blocked keyboard input and blocked MS-Mouse input.

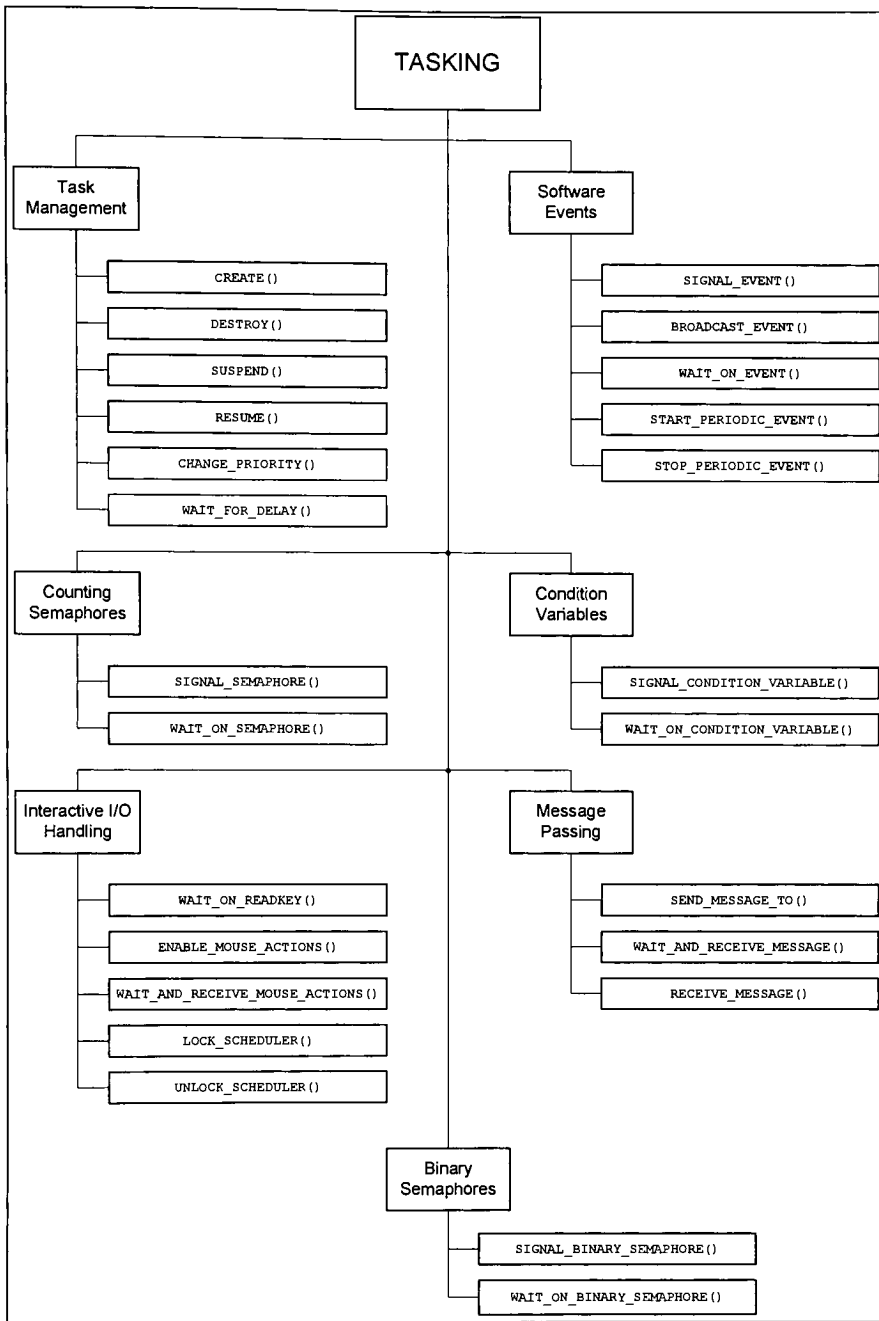
The detailed design of each service and the general operating system philosophy are presented in the remainder of this section.

### **3.1.1 Facilities and Capabilities**

`TASKING` allows for the creation of up to 2047 simultaneous, parallel executing tasks, or threads of execution, within the application. The tasks execute in an environment of 1000 distinct priority levels. The error handling options provided by `TASKING` are quite flexible; each task, or set of tasks, can have its own specific error handling routines; all tasks can use global application error handlers; or `TASKING` can handle any/all of the application errors. When default error handlers are used, `TASKING` will attempt to continue executing in

the presence of errors, but it should be noted that under error conditions, task deadlock is likely to occur (under task deadlock conditions `TASKING` will terminate the application). As regards inter-task constructs, the number of these constructs used by the application is limited only by the amount of available memory in the system. `TASKING` imposes no limits on these constructs. In addition to these features, `TASKING` provides performance and operational statistic reports which can be used to determine how well `TASKING` has been configured to match the needs of the application (reports may also be useful for determining the needs of the application tasks).

The structure of `TASKING` is given in *Figure 7 - Operating System Facilities/Structure*. Detailed component descriptions are provided in the following section and complete `TASKING` source code can be found in *Appendix A RTOS Source Code [TASKING.INT and TASKING.PAS]*.



*Figure 7 - Operating System Facilities/Structure*

The basic architecture of a TASKING application consists of a program ‘main’ body, used only to initialize global data and create initial tasks, and the tasks themselves. No tasks that are created by the ‘main’ body will begin executing until the ‘main’ body actually terminates. Essentially, the main program is used only to initialize the multitasking application. A task

automatically ceases to exist when the body of the task terminates. The program as a whole terminates when all tasks have terminated or when the *halt(...)* system call is made. *Appendix C Support Software Source Code [DINING.PAS]* includes a sample `TASKING` program, it is a solution to the classic multitasking problem proposed and solved by Dijkstra in 1965 called *The Dining Philosophers Problem* [Ref. 18, pg. 56].

#### **3.1.1.1 *CREATE(...)***

This routine creates all system data structures for a parallel executing task. The task will be made ready to run and will be available for execution. Actual execution occurs when the scheduler activates the task (i.e., precise execution time is unknown and depends on the task priority and the relative priority of other 'ready' tasks). In addition, no tasks are allowed to execute until the application 'main' has terminated. In this manner all tasks created in the 'main' body begin executing simultaneously. In addition to creating tasks in the 'main' body, tasks can also create other tasks. There is no concept of task parent-child relationships within `TASKING`; when a task is created it continues to exist until it is destroyed or terminates itself, i.e., the task existence is independent of the existence of the task that created it.

#### **3.1.1.2 *DESTROY(...)***

This routine forces the task to terminate and releases all of the task's resources so that they can be used by other tasks and will then remove all evidence of the task ever having existed. Tasks may only destroy other tasks, not themselves. Tasks destroy themselves by merely exiting the task procedure.

This operation should only be performed when the full consequences of task destruction are known. This is particularly important in cases when the task that is being destroyed communicates with other tasks. If the destroyed task has a resource locked (with a semaphore, for example), there may be no way to reclaim that resource. Another potential risk is that the destroyed task may be the recipient of a message. Since the mailbox structure for that task will cease to exist after the task is destroyed, any subsequent message sent to that task will cause an 'Illegal Operation' error. It is important to note that except for the most simple multitasking designs, deadlock is likely to result from the task destruction. In most cases `TASKING` will detect system deadlock and terminate the application.

### ***3.1.1.3 SUSPEND(...)***

This routine causes a previously ready task to become blocked unconditionally. Once this has occurred, the only way that the suspended task can execute again is if another task resumes it (see section 3.1.1.4 *RESUME(...)*). Note that suspending a task is distinctly different from destroying a task. A suspended task is essentially waiting for the conditions to exist which will cause another task to awaken it. A destroyed task can never be awakened (although it may be re-created).

### ***3.1.1.4 RESUME(...)***

This routine causes a previously suspended task to become ready to execute. After the task has been resumed, it is not necessarily the next task to execute, as is always the case when a task becomes unblocked. The task



priorities of all ready tasks and the scheduling policy determine which task will be next to execute. Note that although typically this action is performed on a task that has been suspended, `TASKING` will allow any task to be resumed, regardless of the reason that the task is blocked. It is important to realize that resuming a task prematurely is likely to cause logical errors in the way that the tasks interact; deadlock may occur.

#### ***3.1.1.5 CHANGE\_PRIORITY(...)***

This routine allows the currently running task to change its priority. It is not possible for a task to change the priority of another task. After the task's priority has been changed, the normal scheduling arbitration takes place and a new task is chosen to execute. It is important to realize that designing a real-time system where a task changes its priority is extremely difficult. There may be isolated instances where it is useful, but in general, it should be avoided.

#### ***3.1.1.6 WAIT\_FOR\_DELAY(...)***

This routine causes the currently running task to become blocked for a specified length of time. The time specified can be as long as thirty-two days with accuracy to one millisecond. When the time delay expires, the task becomes ready to execute. As is always the case when a task becomes unblocked, the task priorities of all ready tasks and the scheduling policy determine which task will be next to execute. Although it is possible to use this mechanism to create a task which performs some operation periodically, `TASKING` provides a more accurate means to do this (see sections *3.1.1.10 START\_PERIODIC\_EVENT(...)* and *3.1.1.11 STOP\_PERIODIC\_EVENT(...)*).

### **3.1.1.7 *SIGNAL\_EVENT(...)***

This routine signals the event specified; if there are tasks waiting on the event, then one task is made ready (based on a non-prioritized, first-in-first-out algorithm). If there are no tasks waiting on the event, then the signal is saved, the next task to wait on the event will consume the signal. Only one signal is maintained no matter how many times the event is signaled. Note that sending multiple signals to the same event is not considered to be an error, but can lead to logical problems on the receiving task's subsequent reactivation.

### **3.1.1.8 *BROADCAST\_EVENT(...)***

This routine is functionally identical to section 3.1.1.7 *SIGNAL\_EVENT(...)* with minor differences. If there are tasks waiting on the event when it is signaled, then they are all made ready to execute. If there are no tasks waiting on the event, then the signal is saved and the next single task to wait on the event will consume it.

### **3.1.1.9 *WAIT\_ON\_EVENT(...)***

This routine either causes the calling task to be suspended waiting for the event to be signaled or else consumes the signal already stored in the event. If the event has already occurred, then the calling task will not be suspended but instead will immediately become ready to execute. Although technically the task has not become unblocked, `TASKING` treats these two conditions identically. Therefore, the task priorities of all ready tasks and the scheduling policy determine which task will be next to execute.

### ***3.1.1.10 START\_PERIODIC\_EVENT(...)***

This routine causes `TASKING` to begin the periodic signaling of the event specified at the specified interval. This mechanism is most useful for creating periodic tasks. Since `TASKING` is responsible for signaling the event, the task is guaranteed that the signal will occur at the specified interval. It is important to note that the task priorities of all ready tasks and the scheduling policy will actually determine the exact execution period of the task. Requesting a periodic event is not the same as a real-time design which guarantees that the periodic execution will result.

### ***3.1.1.11 STOP\_PERIODIC\_EVENT(...)***

This routine causes `TASKING` to stop the periodic signaling of the event specified (see section 3.1.1.10 *START\_PERIODIC\_EVENT(...)*). Note that the event is not destroyed by this operation; it may still be used by restarting the periodic event.

### ***3.1.1.12 SIGNAL\_BINARY\_SEMAPHORE(...)***

This routine performs an 'Up' operation on the specified binary semaphore. Since the 'Up' operation is not defined for a binary semaphore that is already set, `TASKING` will detect this condition as an error. If there are tasks waiting on the semaphore when it is signaled, then one is made ready to execute. If there are none waiting then the semaphore is set.

### ***3.1.1.13 WAIT\_ON\_BINARY\_SEMAPHORE(...)***

This routine performs a 'Down' operation on the specified binary semaphore. If the binary semaphore is already zero, then the running task is

suspended until the semaphore is signaled. If it is non-zero, then the semaphore is cleared and execution continues, based on the priorities of all ready tasks and the scheduling policy that is in effect.

#### ***3.1.1.14 SIGNAL\_SEMAPHORE(...)***

This routine performs an 'Up' operation on the specified "Dijkstra" semaphore. If there are tasks waiting on the semaphore when it is signaled, then one is made ready to execute. If there are no tasks waiting, then the semaphore is incremented. In either case, execution continues, based on the priorities of all ready tasks and the scheduling policy that is in effect.

#### ***3.1.1.15 WAIT\_ON\_SEMAPHORE(...)***

This routine performs a 'Down' operation on the specified "Dijkstra" semaphore. If the "Dijkstra" semaphore is already zero, then the running task is suspended until the semaphore is signaled. If it is non-zero, then the semaphore is decremented and execution continues, based on the priorities of all ready tasks and the scheduling policy that is in effect.

#### ***3.1.1.16 SEND\_MESSAGE\_TO(...)***

This routine sends the message (actually, the pointer to the message) to the specified task by placing it in the task's mailbox (the mailbox will automatically be created if necessary). Although the calling task cannot block when sending a message, `TASKING` re-evaluates which task should be executing, and execution continues based on the priorities of all ready tasks and the scheduling policy that is in effect. Typically, messages are dynamically allocated (using `getmem(...)`)

before being sent, the receiver is responsible for deallocation (using *freemem(...)*) after the message has been used.

#### ***3.1.1.17 WAIT\_AND\_RECEIVE\_MESSAGE(...)***

This routine retrieves a message (actually, a pointer to a message) from the caller's mailbox (which is created if necessary). If there are no messages in the mailbox then the task is blocked waiting for a message to arrive. Because the caller is blocked until a message is available, the return pointer is guaranteed not to be nil. Typically, the receiving task will dispose of the message after it has been used. Messages are queued in first-in-first-out order into the task's mailbox (extreme caution should be used in the design of the message passing tasks to avoid unbounded priority inversion).

#### ***3.1.1.18 RECEIVE\_MESSAGE(...)***

This routine is functionally identical to section *3.1.1.17 WAIT\_AND\_RECEIVE\_MESSAGE(...)* with minor differences; if there are no messages present in the caller's mailbox at the time of the call, then the return value is nil. Therefore, the task cannot become blocked by calling this routine.

#### ***3.1.1.19 SIGNAL\_CONDITION\_VARIABLE(...)***

This routine signals the condition variable specified; if there is a task waiting on the condition variable then it is made ready. Multiple waiting tasks are handled in first-in-first-out order. If there are no tasks waiting on the condition variable when this routine is called, the signal is lost, (recall that condition variables have no storage capabilities).

### ***3.1.1.20 WAIT\_ON\_CONDITION\_VARIABLE(...)***

This routine is guaranteed to cause the calling task to become blocked. Because condition variables do not store signals, the task is guaranteed to be waiting on the condition variable as a result of calling this routine. As soon as another task signals the condition variable the waiting task will be made ready.

### ***3.1.1.21 WAIT\_ON\_READKEY(...)***

This routine provides the means for an application task to perform I/O in a 'blocking' manner. This means that while there is no input from the keyboard, the calling task is suspended. As soon as input arrives from the keyboard, the calling task is made ready to execute. In addition, the ASCII code of the key that was pressed is returned to the caller. This is an example of interrupt driven multitasking. The task that calls this function (of which there can be only one) provides the interface between the rest of the application tasks and the outside world (in this case, the user). Typically, the task responsible for user input would send messages to other tasks or signal some type of inter-task communication construct as a result of receiving the user input. Note that `TASKING` cannot detect deadlock if this feature is used within the application. This is due to the very nature of an interrupt driven task. It could be awakened at any time by the keyboard; therefore deadlock is not detectable.

### ***3.1.1.22 ENABLE\_MOUSE\_ACTIONS(...)***

This routine enables the specified MS-Mouse events to be received by the task which is (or will be) waiting to receive mouse actions. Essentially, this allows the task to communicate with `TASKING` the exact MS-Mouse features

that it is using. This routine can be called at any time to change the mouse actions that are of interest to the task.

### ***3.1.1.23 WAIT\_AND\_RECEIVE\_MOUSE\_ACTIONS(...)***

This routine receives MS-Mouse action parameters from the MS-Mouse driver and passes them to the waiting task. If none of the enabled mouse actions has occurred, then the task is blocked, waiting for one (or more) action to occur. This is another means of providing an application task the ability to perform I/O in a 'blocking' manner. Note that `TASKING` cannot detect deadlock if this feature is used within the application. This is due to the very nature of an interrupt driven task. It could be awakened at any time by the MS-Mouse action; therefore deadlock is not detectable.

### ***3.1.1.24 LOCK\_SCHEDULER(...)***

This routine allows a task to 'lock' the scheduler and prevent all further task rescheduling until it is unlocked (see section *3.1.1.25 UNLOCK\_SCHEDULER(...)*).

One use of this routine is when multiple tasks need to write to the display, and it is desired that the output from each task occurs without the output from other tasks intermixing with it. Note that during the time that the scheduler is 'locked' there will be no preemptive or cooperative task rescheduling. The application should avoid locking the scheduler if possible. If that is not feasible or practical for the application, then the time spent with the scheduler locked should be an absolute minimum.

It is imperative that each call to lock the scheduler have a corresponding call to unlock it. Successive (nested) calls to lock the scheduler are allowed, but the application should ensure that each scheduler lock request has a corresponding call to unlock it.

### **3.1.1.25 UNLOCK\_SCHEDULER(...)**

This procedure 'unlocks' the scheduler and allows task rescheduling to continue based on the previously existing configuration. Note that the scheduler must previously have been locked, or an 'Illegal Operation' error will occur (see section 3.1.1.24 *LOCK\_SCHEDULER(...)*).

## **3.1.2 Configuration Parameters**

TASKING provides a number of configurable parameters which help to match the RTOS to the real-time application. The following parameters must be initialized by the application before multitasking begins; therefore this initialization must occur in the application 'main' body or from the application INI file, which TASKING will read before the 'main' body executes. Once multitasking begins (i.e., when the 'main' body terminates), TASKING will ignore all changes to these parameters:

- Scheduling Policy - Task priorities can be either static throughout the execution of the application (default), or they can rotate on every context switch. Static priorities will ensure reliable, predictable system performance whereas rotating priorities will add a degree of uncertainty to the scheduling algorithm which may result in better performance for



some applications. For example, an application where 'fairness' is an important system feature may benefit from rotating priorities.

- Priority Inheritance - Task priority inheritance will be either enabled or disabled. Priority inheritance is effective in solving the problem of unbounded priority inversion. Priority inversion can occur when low and high priority tasks share a common resource and there exist intermediate priority tasks. If a low priority task 'locks' a resource and a high priority task requests that resource (i.e., attempts to 'lock' the resource), it is possible for an intermediate priority task to indefinitely prevent the low priority task from 'unlocking' the resource. This essentially creates the condition where an intermediate priority task can prevent a high priority task from executing (see *Figure 8 - Unbounded Priority Inversion*).

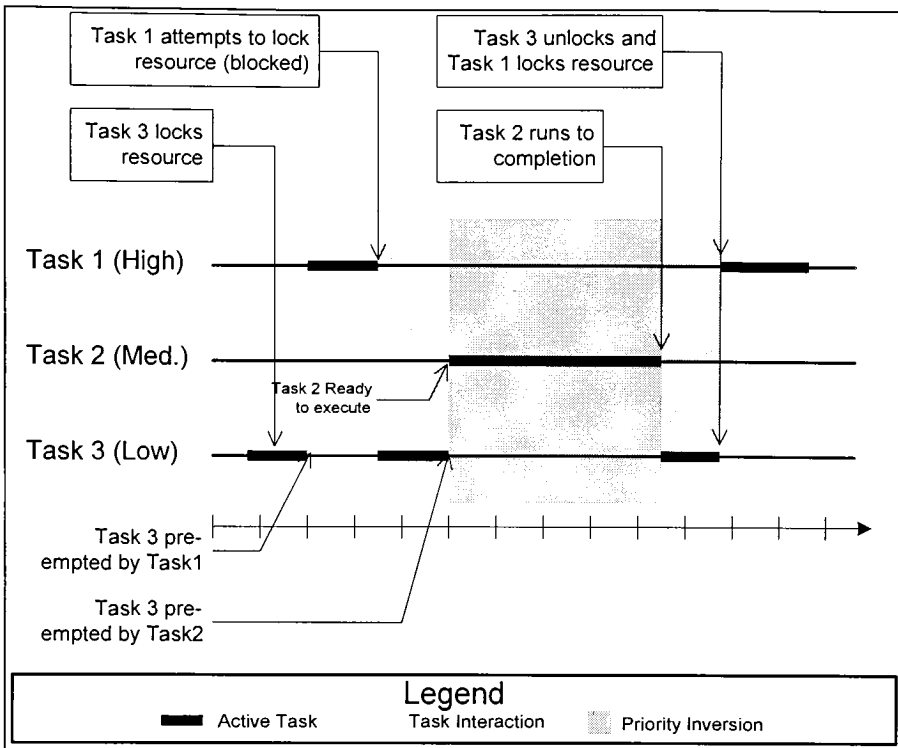


Figure 8 - Unbounded Priority Inversion

Priority inheritance prevents this situation from arising by temporarily raising the priority of a task with a locked resource to the level of the highest priority task attempting to lock that same resource for the duration of time that the lower priority task maintains control of the resource (see *Figure 9 - Priority Inheritance Protocol*).

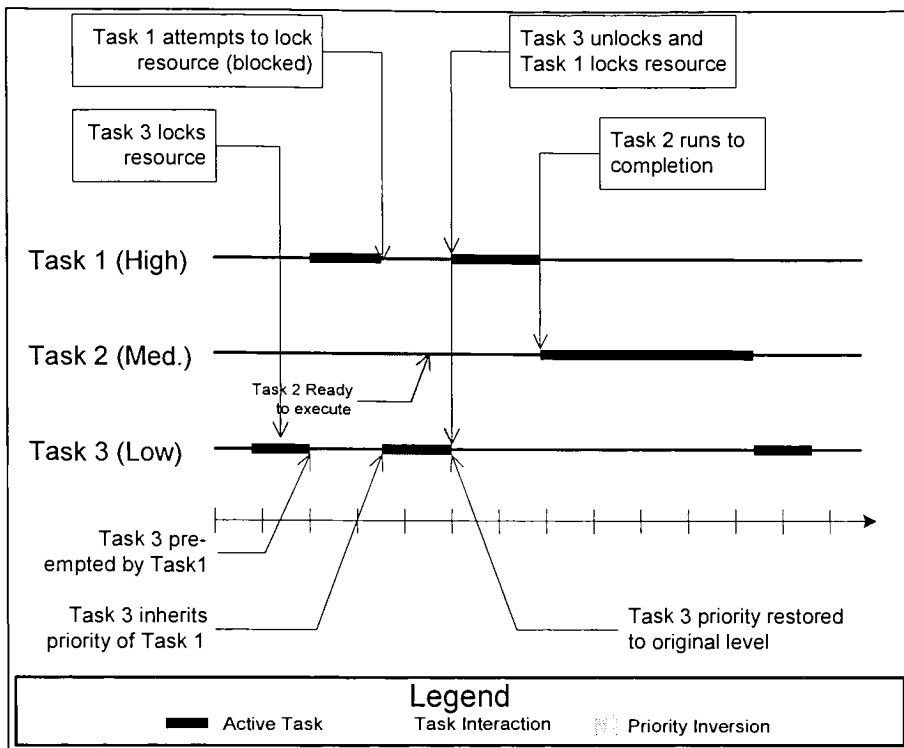


Figure 9 - Priority Inheritance Protocol

- **Tasking Model** - Both cooperative and preemptive multitasking models are supported. Cooperative multitasking only allows context switching when `TASKING` services are called. Cooperative multitasking requires that all tasks 'surrender' the CPU when they have completed their work (i.e., the tasks must cooperate by blocking to allow other tasks to execute, and there can be no continuously executing background tasks). This tasking model is often referred to as 'prioritized, run-to-completion'. In addition to allowing context switches when `TASKING` services are called, preemptive multitasking causes context switches periodically based on the timeslice parameter. The timeslice can range from 50  $\mu$ Sec. to 65 mSec. Preemptive multitasking removes the requirement that the tasks must cooperate in passing the CPU resource

from task to task. When preemptive multitasking is selected, `TASKING` will 'give' the CPU to a task for a single timeslice. After that, the CPU will be 'taken' from that task and 'given' to another task which is chosen based on the relative task priorities of all tasks which are ready to execute and the scheduling policy that is in effect. Note that under this multitasking model, continuously executing background tasks are allowed.

### 3.1.3 RTOS Performance

The most important performance parameter of an RTOS is that of overhead imposed by the RTOS on the application. This is the amount of processing that is performed by the RTOS to achieve its goals which in no way contributes to the completion of the application requirements. `TASKING` performance in this area is shown in the following figures. Note that although these measurements were made on a 66 MHz Pentium processor, a different x86-class processor would produce similar results, i.e., a faster processor would show a decrease in overhead proportional to its speed relative to a 66 MHz Pentium.

In order to fully understand the performance merits of `TASKING` there are two things that must be explained: how the measurements were made and why the results turned out as they did.

The measurements were made using two special programs to generate the data and another program to produce a graph, see *Appendix C Support Software Source Code* [`TASKS-A.PAS`, `TASKS-B.PAS` and `TSK-BNCH.PAS`]. First, a baseline measurement was made using a program which performs a set of

computations without using `TASKING` in any way [*TASKS-B.PAS*]. This produces a measurement of the best possible performance since any additional processing performed by `TASKING` results in performance which is degraded from this baseline measurement. Next, measurements are made using a program that performs the same set of computations but in this case `TASKING` is used to create multiple tasks each of which performs a subset of the computations [*TASKS-A.PAS*]. For these measurements `TASKING` is configured for preemptive multitasking with priority inheritance disabled (which is not relevant because none of the tasks interacts or shares resources). In addition, all tasks execute at the same priority level with static priorities. The point is to measure the impact of context switching on the time required for the application to perform all of the computations. This RTOS configuration achieves that goal by disabling or not using all other `TASKING` features. The only RTOS parameter that is allowed to vary is the target timeslice. The timeslice is varied from 1 mSec to 250 mSec; this range encompasses most reasonable timeslice values.

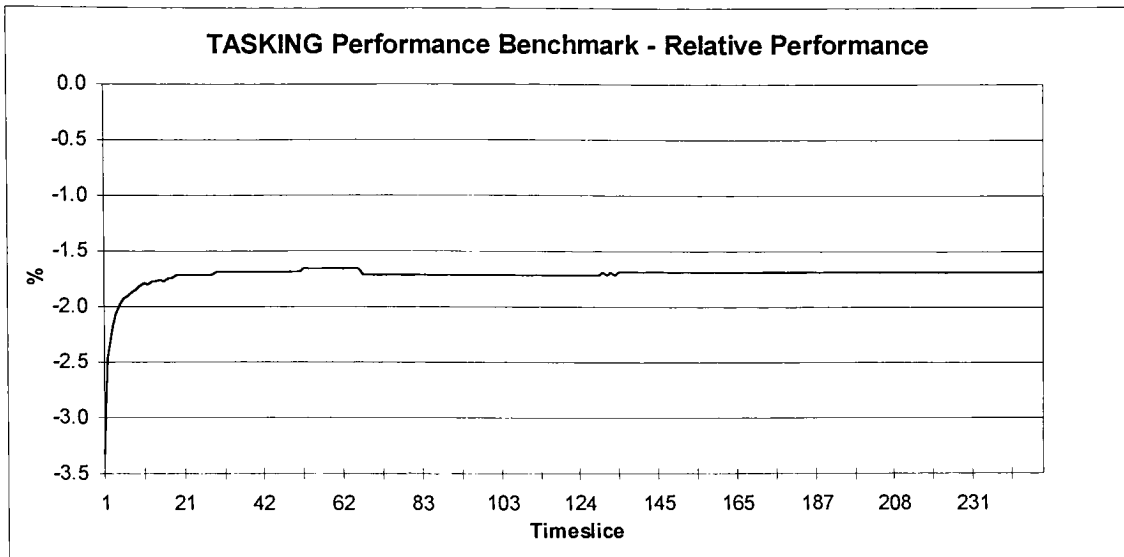
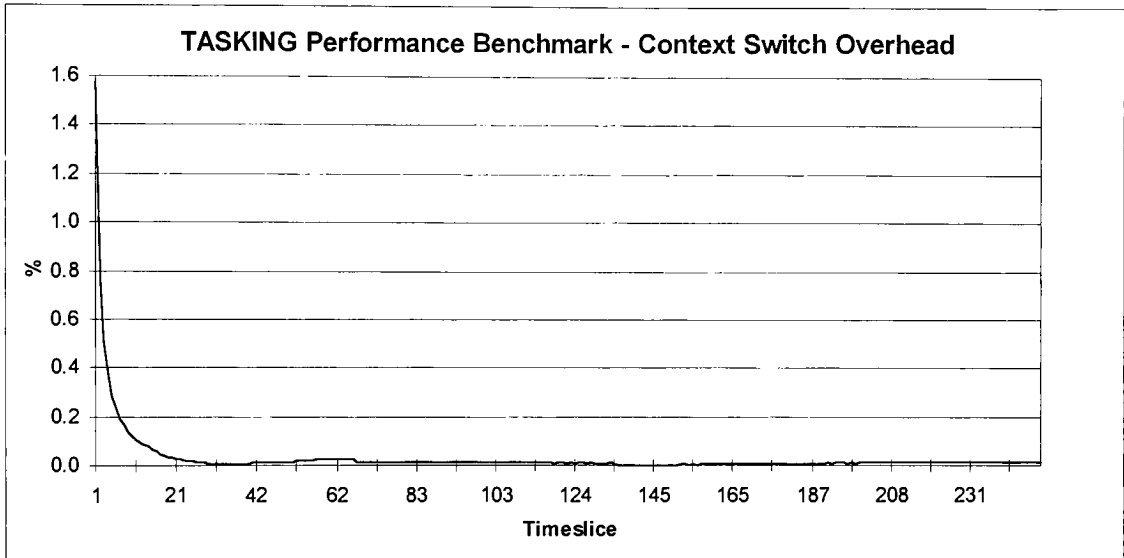


Figure 10 - TASKING Benchmark (Relative Performance)

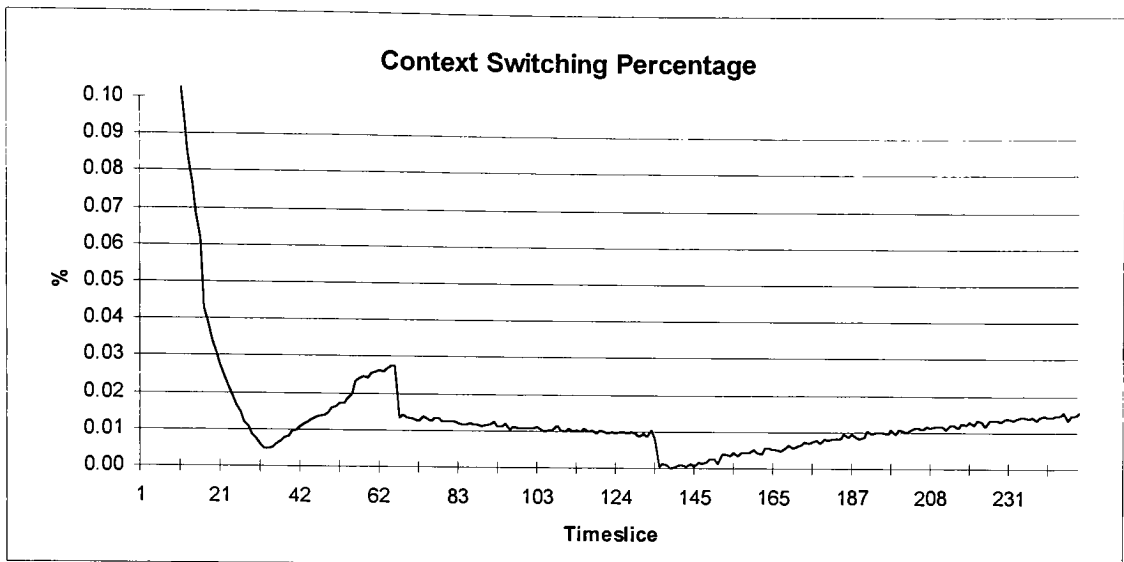
The results from these measurements are shown above in the *Figure 10 - TASKING Benchmark (Relative Performance)*. This graph shows that TASKING imposes a minimum of approximately 1.7% performance penalty on the application. What is most interesting is that after the timeslice value exceeded approximately 20 mSec, the performance degradation remained relatively constant. This implies that ~1.7% performance degradation is the best that TASKING can achieve. As the graph clearly shows, the overhead increases significantly for small timeslice values. The reason for this is that smaller timeslice values result in more context switches. As the following graphs show (*Figure 11* and *Figure 12*), the context switch overhead is the dominate factor in TASKING's performance only for small timeslice values. For larger timeslice values, the dominate (and constant) factor in TASKING's performance is TASKING's management of the system clock. That is, the handling of the system clock interrupt and the internal delay request checking performed during that interrupt service routine.

*Figure 11 - TASKING Benchmark (Context Switching Overhead)* below shows that context switching is not responsible for the performance degradation. Note that after the timeslice value exceeded approximately 20 mSec, the context switching overhead remained relatively constant at zero. It is only for small timeslice values that the context switch overhead is appreciable.



*Figure 11 TASKING Benchmark (Context Switching Overhead)*

As shown below in *Figure 12 - TASKING Benchmark (Context Switching Percentage)*, context switching time itself is negligible when compared to the timeslice value. The context switch time is only significant for small timeslice values. It is the presence of `TASKING` itself that impacts performance. This is a good indication that the performance degradation is relatively constant regardless of the application being used. This is a very important characteristic of a real-time operating system. Note that none of the values of *Figure 12* is zero, rather, some values are merely very small (or more likely beyond the measurement capabilities of the data collected).



*Figure 12 TASKING Benchmark (Context Switching Percentage)*

Analysis of the context switching time of TASKING over the same range of timeslice values as in *Figure 12* yields a value which never exceeds approximately 37  $\mu$ Sec. The context switching time does vary somewhat over this range but never beyond the accuracy of the measurements. That is, to within the accuracy possible with the data collected, the context switching time of TASKING remains relatively constant over this range of timeslice values. These two context switching time properties (relatively constant and guaranteed upper bound) are probably the most important RTOS characteristics for real-time applications. A constant context switch time allows accurate, verifiable timing analysis to be performed and an upper limit on the context switch time allows worse-case timing analysis to be performed. No system is actually ‘real-time’ if the context switch time varies excessively or is unbounded.

The performance degradation which occurs when using TASKING is a result of the monitoring of the system clock that TASKING must do in order to achieve preemptive multitasking. TASKING must change the clock interrupt



period of the main system clock hardware of the PC to a minimum of 1.000 kHz (1.000 mSec period) for timeslice values of 1 mSec or more. Because of this and the fact that the PC clock is originally set to 18.20 Hz (54.93 mSec period), TASKING performance is best at multiples of approximately 55 mSec. This is illustrated below, (see *Figure 13 - TASKING Benchmark (System Clock Impact)*). The graph shows the excess processing that TASKING must perform in order to simulate the MS-DOS clock timer. TASKING must determine the proper time that the original MS-DOS clock timer interrupt should be called so that MS-DOS will continue to keep the proper 'clock time'. Since MS-DOS assumes that its timer interrupt handler will be activated once every 54.93 mSec, in order for MS-DOS to keep proper time, TASKING must ensure that it does. It is this simulated interrupt which tends to add modulated irregularities to the TASKING performance curves. If MS-DOS was not used by the application, TASKING would not need to simulate the interrupt. That would stabilize the performance curves for TASKING but would not allow the application to use any MS-DOS services (such as screen and/or disk I/O). Since it is not practical to include all of those services into TASKING, the only alternative is to understand the performance measurement limitations and factor those limitations into the interpretation of the performance data.

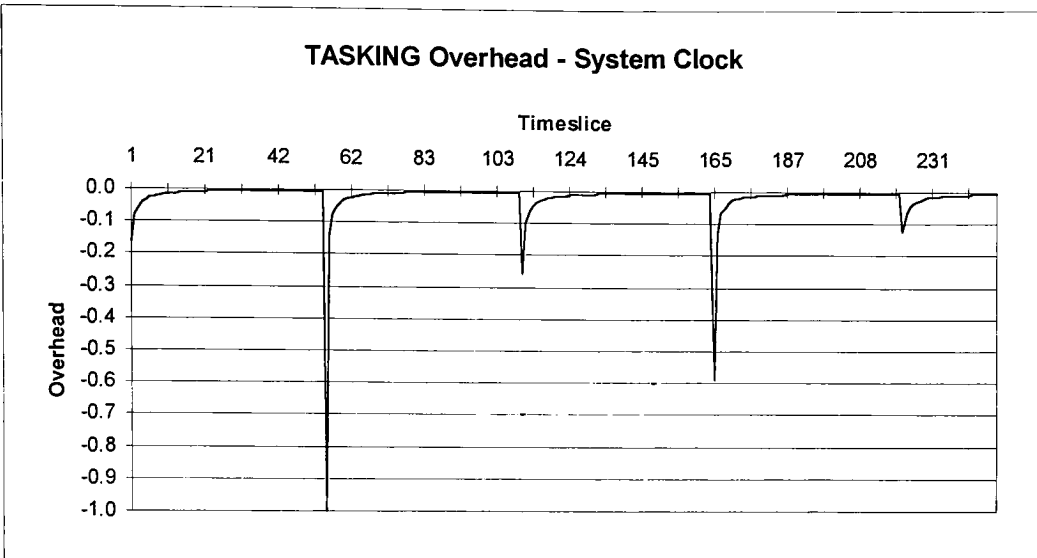


Figure 13 TASKING Benchmark (System Clock Impact)

In general, TASKING provides reliable, consistent performance. The only concern is that the application is in fact degraded by approximately 1.7%. Although this value is not excessive, it is also not negligible. Given that all of TASKING is written in Pascal and that there is no appreciable assembly language, these results are not surprising and certainly not unreasonable.

### 3.2 Genetic Algorithm

The GA is constructed in the 'standard' genetic algorithm manner (see *Figure 14 - Genetic Algorithm Process Flow*). The only exception to this is that the information within the GA will flow through data files created on the computer hard drive. This is necessary because the GA must be completely suspended in order to run the TASKING programs which are used to evaluate the individuals of the population. In order to achieve accurate and reproducible results, the GA must be run from a minimum PC configuration, i.e., no Terminate and Stay Resident (TSR) programs loaded. In addition, the GA must be run in a

pure MS-DOS<sup>®</sup> environment, i.e., not from within an MS-DOS<sup>®</sup> shell under MS- - -  
Windows<sup>®</sup>.

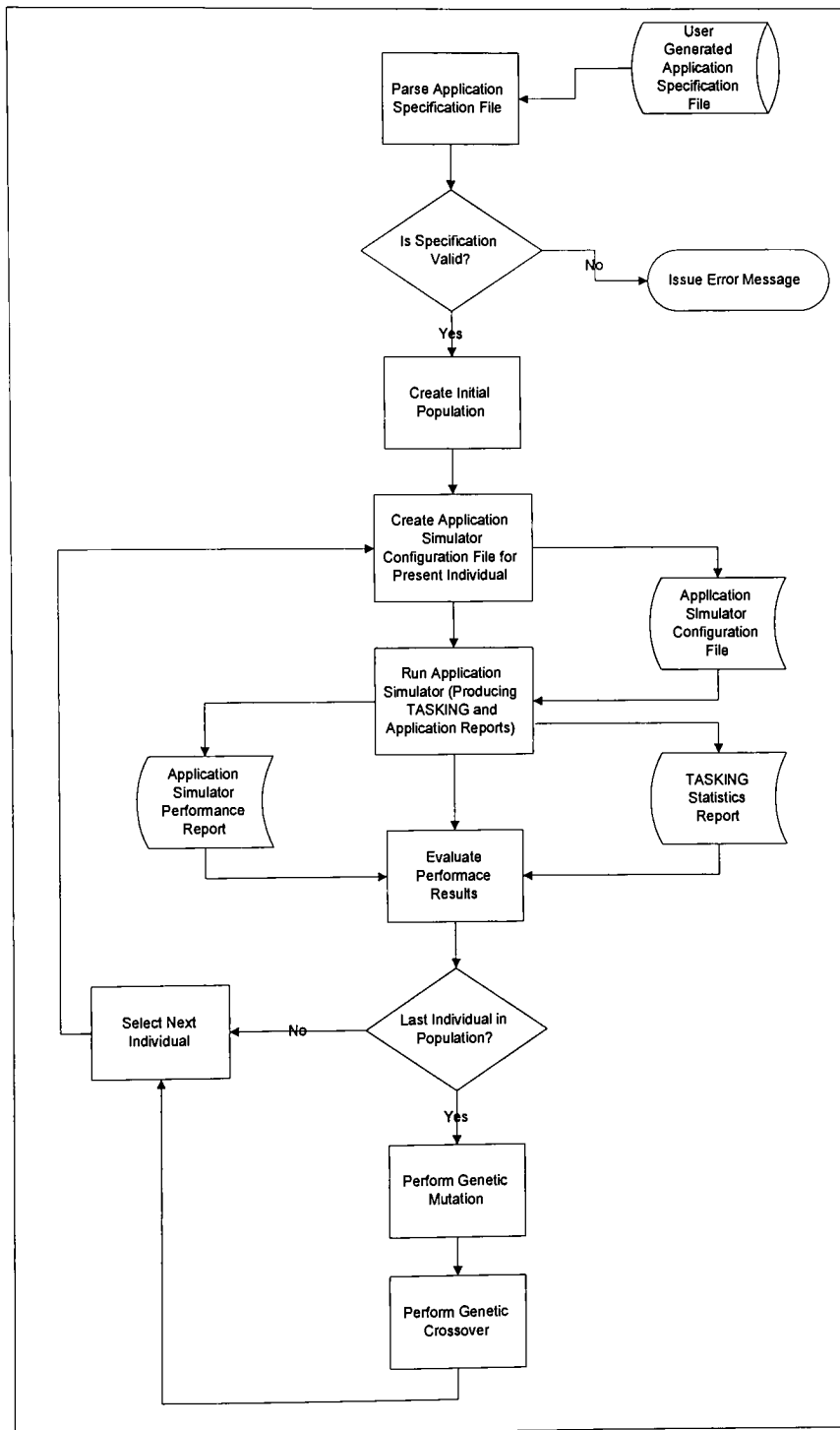


Figure 14 - Genetic Algorithm Process Flow

The most important aspect of the GA design is in the choice and/or design of the genetic operators: mutate(...) and crossover(...). These operators will affect the ability of the GA to converge to a solution, the rate of convergence and the probability of false convergence. Most of the research associated with this thesis has involved these operators. There is no shortage of opinions regarding the 'best' universal operators, but there has been nothing published regarding operators specific to this problem domain. A certain amount of experimentation is required in determining viable operators. The following sections detail the appropriate operators for this application and also describe the desirable characteristics of genetic operators.

Before the genetic operators can be designed, the genetic representation that will be used for the problem must be determined. The following section describes the genetic representation used for the GA.

### **3.2.1 Genetic Representation**

The genetic representation for the RTOS is representational in nature, that is, it is not a pure binary representation. Instead the representation will closely mirror the configuration parameters of the RTOS, see Section 3.1.2 *Configuration Parameters*. The following table lists the genes that comprise the genotype used for the RTOS configuration:

Gene	Possible Values
Tasking Model	Cooperative, Preemptive
Target Timeslice	50 - 65,535 $\mu$ Sec
Priority Inheritance Enabled	True, False
Priority Allocation	Static, Rotating
Initial Priority Assignment	Constant, Random, Rate Monotonic, Deadline Monotonic, Workload Monotonic

*Table 3 - RTOS Genotype*

Since the GA will not use a binary representation, it is important to note that there are five genes in the genotype above. All five are considered to be equal even though the actual representation length (bit length) of each gene is different. For the purposes of genetic mutation and crossover, they are considered to be of equal length. Therefore, each of the five has equal probabilities whenever a gene is to be operated upon in a random manner. This is considerably different from a binary representation GA where each bit is considered equal, and the length of the group of bits that comprise a gene has a significant impact on the relative probability of that gene becoming involved in mutation and/or crossover.

### 3.2.2 Mutation

The genetic mutation(...) operator for this GA simply mutates (which is a gene specific operation) all genes of the population with equal probability  $P_M$ . For example, suppose  $P_M=4\%$  (i.e., 0.04), the number of individuals  $N=225$ , and the number of genes per genotype  $M=5$ ; then for a typical generation there would be  $P_M \cdot N \cdot M = 0.04 \cdot 225 \cdot 5 = 45$  genetic mutations, on average.

For the representation used here, the mutation to be performed is a function of exactly which gene is to be mutated. As is the case for all

representational genotypes, this genetic algorithm mutation operator randomly selects a value from the allowable range for the gene to be mutated. In this way all genotypes are guaranteed to be made up of valid genes. This ensures that all individuals represent valid solutions to the problem at hand even if they have been involved in a genetic mutation. The allowable values for each gene are defined in *Table 3 - RTOS Genotype* above.

### 3.2.3 Crossover

The genetic crossover(...) operator for this GA is implemented using single-point crossover. First, two parents are selected at random from all above average individuals (i.e., individuals that have a fitness value greater than the average fitness for all individuals of the population). Then a crossover point is randomly selected. The crossover is performed as all genes up to the crossover point are taken from the first parent, and the remaining genes are taken from the second parent. The newly created individual replaces a below-average individual. Since there are five genes in the genotype, the crossover point is from one to four. This ensures that at least one gene is taken from each parent.

The results achieved when using single point crossover are somewhat dependent upon the ordering of the genes. Since the genes that are taken from the parents are contiguous, it is useful to group related genes together. For this application, the only genes that are related are the Tasking Model and Target Timeslice genes. As indicated in *Table 3 - RTOS Genotype* above, these two genes are adjacent within the genotype.

### **3.3 Problem Specification**

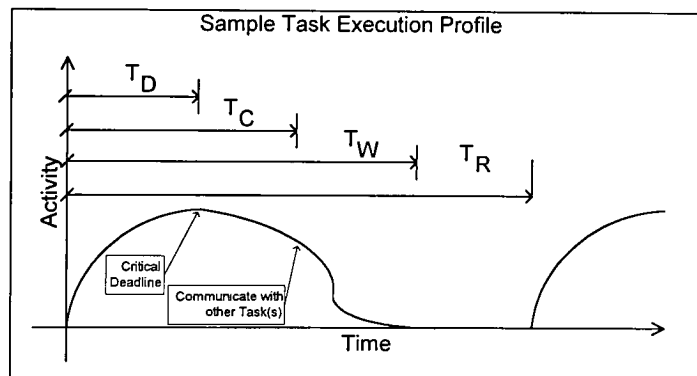
The real-time application must be specified in the most flexible manner possible in order to support an arbitrary real-time design. In order to do that, the GA reads the design information from a 'standard' data file. This approach allows the real-time application specification to be created independently of the GA. The following sections describe which aspects of the real-time application must be specified and how the specification must be created so that the GA can correctly process (parse) the information.

#### **3.3.1 Specification Parameters**

In order to completely specify the real-time application, the user will have to provide details about the characteristics of the tasks that make up the application. The only aspects of the application design that are germane to the GA are those that involve the tasks themselves and the task interaction. The details of what operations the tasks perform are irrelevant. The necessary task characteristics are as follows:

- Task Name - This will be a text string that will be used when performance results are reported and when communication is performed with other tasks.
- Period - This will be the amount of time between successively beginning executions of a periodic task. A value of zero is used to indicate that the task is non-periodic. Non-periodic tasks will only execute once. See *Figure 15 - Task Execution Profile, T<sub>R</sub>*.

- Deadline - This will be the amount of time between when a task begins execution and when it must reach a critical point in its execution. See *Figure 15 - Task Execution Profile,  $T_D$* .
- Deadline Hardness - This parameter will indicate the degree to which the previously specified deadline is 'hard'. A linear scale between 1 (soft) and 10 (hard) is used.
- Workload - This will be the amount of time between when a task begins executing and when it completes. A value of zero is not allowed. See *Figure 15 - Task Execution Profile,  $T_W$* .



*Figure 15 Task Execution Profile*

In addition, the user must specify the inter-task communication aspects of the real-time system. The following communication characteristics may also be specified (some applications and/or tasks will not require these features):

- Task Name - This will be the name of the task that is to be the recipient of the communication.



- Communication Type - This will be the `TASKING` inter-task communication type that is to be used (Message, Semaphore, Binary Semaphore, Event or Condition Variable).
- Communication Time - This will be the amount of time from when the task begins executing to when the communication should be performed with the other task. See *Figure 15 - Task Execution Profile*,  $T_C$ .

Up to two pairs of the above communication parameters can be specified. This will allow most complex inter-task communication architectures to be fully described.

### 3.3.2 Specification File Format

Since it is not the goal of this thesis to develop an application for creating the real-time application specification, a standard PC spreadsheet program will be used to create the file used to specify the design of the real-time system. Any spreadsheet can be used as long as the file is saved in the standard 'Comma Separated Value' format (all commercial spreadsheet applications support this format, as well as most commercial database applications). Text fields will be enclosed in double quotes; integer fields will be ASCII text but will not be enclosed in quotes; real numbers are not necessary and will not be supported.

Each line of the file will represent a single task specification, the length of which is limited to 255 characters. The order of the fields within each line of the spreadsheet must as specified in *Table 4 - Real-Time Application Specification* below.

Task Name	Period	Deadline	Hardness	Workload	Receive Resource #1	Receive Resource Type #1	Receive Resource Time #1	Signal Resource #1	Signal Resource Type #1	Signal Resource Time #1	Receive Resource #2	Receive Resource Type #2	Receive Resource Time #2	Signal Resource #2	Signal Resource Type #2	Signal Resource Time #2
Text	mSec	mSec		mSec	Text	Semaphore Cond_Var Event Message BSemaphore	mSec	Text	Semaphore Cond_Var Event Message BSemaphore	mSec	Text	Semaphore Cond_Var Event Message BSemaphore	mSec	Text	Semaphore Cond_Var Event Message BSemaphore	mSec
25 Char	0-1000	0-1000	1..10	0-1000	25 Char	10 Char	0-1000	25 Char	10 Char	0-1000	25 Char	10 Char	0-1000	25 Char	10 Char	0-1000

*Table 4 - Real-Time Application Specification*

## 4. Evolution Tool Set

This thesis will result in a set of PC-based tools which will compute near optimal RTOS parameters for any real-time application. The tools will provide a means of creating an application specification and performing the GA based analysis, using a custom RTOS - `TASKING`. Performing the analysis will include real-time graphical display of the progress of the GA 'population' and a report with the recommended near optimum RTOS configuration.

The set of project deliverables is as follows (detailed in the following sections):

- Custom RTOS (`TASKING`) to be used in evaluating RTOS parameters.
- Tool to analyze and evaluate an arbitrary real-time application using a specific set of RTOS parameters.
- Genetic Algorithm based tools to evolve a near optimal RTOS configuration for the real-time application.
- Set of real-time application specifications and analysis to verify proper operation of the tool set (i.e., problems with known solutions).

### 4.1 Custom RTOS (`TASKING`)

The RTOS used to evaluate the real-time application has already been described in detail (see section 3.1 *Real-Time Operating System (`TASKING`)*). The aspect of `TASKING` that has not been described is how it is instrumented to provide performance information. `TASKING` is capable of capturing and

reporting performance information about three different aspects of the execution of an application. These aspects are:

- Hardware Interrupts - includes information about which interrupts were serviced, how many times they were serviced and how frequently they were serviced (report filename: `IRQ.RPT`).
- Software Interrupts - includes information about which MS-DOS services were used by the application, how many times they were called and how frequently they were called (report filename: `SERVICES.RPT`).
- Task Statistics - includes information about the CPU and stack utilization of each task as well as the percentage of periodic events that failed to be signaled at the proper time (report filename: `TASKING.RPT`).

As it relates to the evaluation of an RTOS configuration for a real-time application, the task statistics information is the most useful. The following figure shows a sample report for the Dining Philosophers Problem (*Figure 16 - Sample TASKING.RPT File*).

```

1: Preemptive Multitasking Statistical Information
2: -Priority Inheritance Enabled
3: -Static Priorities
4:
5: Task Activity:
6: Context Switches          2002 (~640 per second)
7: Cooperative              361 (~115 per second)
8: Preemptive              1641 (~524 per second)
9: Target Time Slice        = 1500 usec
10: Achieved Time Slice     = ~1563 usec
11: Available CPU Bandwidth = 0.000 %
12: Periodic Event Faults  29.941 %
13: Tasks at Termination:
14:   Task State   Task      Stack Size Stack Used Stack Used CPU Used
15:   (Delayed Until) ID  Priority  (words)   (words)   (%)       (%)
16:
17:   Ready        1      9991     4000      110       2.8      58.4
18:   Ready        2      9986     4000      110       2.8      39.5
19:   Ready        3      9966     4000       31       0.8       0.0
20:   Running     4     10000     2000      687      34.3      2.1
21:
22: Maximum percent of stack used: 34.3%
23:
24: Heap Information (bytes):
25: Total Available:          100,000
26: Used for Application Stacks: 28,000
27: Available to Application:  72,000
28:
29: Execution Time:
30: Effective = 3.070 seconds
31: Absolute  = 3.13 seconds

```

Figure 16 - Sample TASKING.RPT File

The parameters of interest are on lines 11 and 12. The CPU utilization is determined by TASKING by recording the task that is executing when system timer interrupts occur. Since TASKING maintains a ‘null task’ which executes only when there are no application tasks which are ready to execute, TASKING can very easily (and accurately) determine the percentage of available CPU bandwidth. The percentage of periodic events that are missed by the application is also very easily determined. Since TASKING is in control of the system clock, whenever a periodic event is signaled the state of the event is examined. If the event is already signaled then the application will miss that event (i.e., the application was not able to finish its execution during the previous time period). These two parameters represent TASKING’s contribution to the RTOS evaluation by providing information usually available only to the operating system to the genetic algorithm.

## 4.2 Real-Time Application (RTOS-APP)

The real-time application produces a report which includes a measure of timeliness for various aspects of the application timing requirements. The real-time application monitors the system clock (provided by `TASKING`) to determine if the task's deadline and periodic execution occurred at the appropriate time. Since the measurements are made by the tasks themselves, the results are very accurate. `TASKING` is capable of signaling a periodic event at the correct time but the RTOS configuration and application design dictate whether the task will actually receive the signal at the correct time or not. The following figure shows a sample real-time application report file:

```
1: Individual Task Timeliness
2:   Task Name: TASK #1
3:   Periodicity: 213/213 (100.00%)
4:   Deadline:   14/214 (6.54%)
5:
6:   Task Name: TASK #2
7:   Periodicity: 141/141 (100.00%)
8:   Deadline:   68/141 (48.23%)
9:
10:  Task Name: TASK #3
11:  Periodicity: 0/0 (0.00%)
12:  Deadline:   0/0 (0.00%)
13:
14:
15: Average Task Timeliness
16: Periodicity: 66.67%
17: Deadline:   13.51%
18:
19: Application Timeliness = 40.09%
```

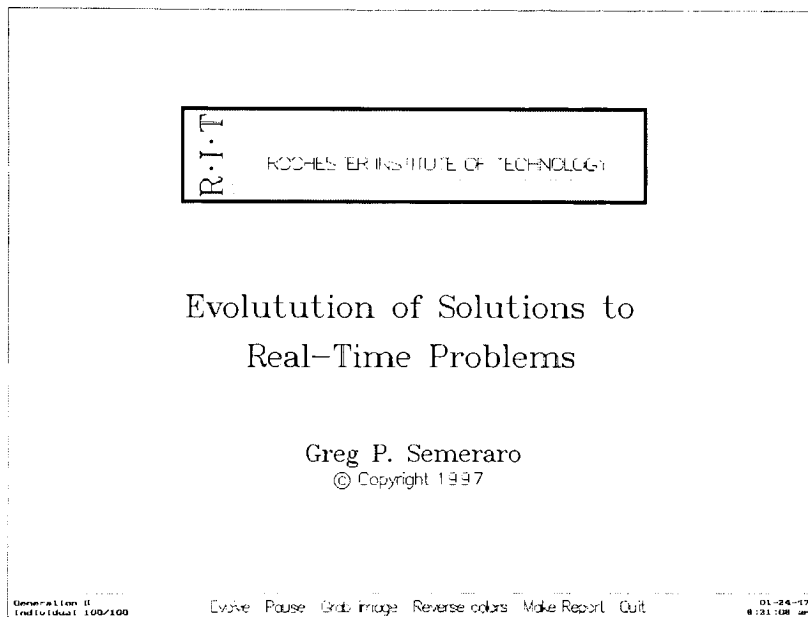
*Figure 17 - Sample RTOS-APP.RPT File*

The parameter of interest is shown on line 19, it is this value that is used to contribute to the evaluation of the RTOS configuration for the specific real-time application.

## 4.3 Genetic Algorithm (GRTOS-GA)

The GA performs the operations which result in evolving a near real-time RTOS configuration for the real-time application. The evaluation of the RTOS configuration is performed by parsing the `TASKING.RPT` and `RTOS-APP.RPT` files and combining the results. At the completion of the evolution process, the

GA produces a report (GA\_XX.RPT, where XX is a sequence number) which indicates what the top five RTOS configurations were. The startup screen for the tools is shown below (*Figure 18 - GRTOS-GA Startup Screen*). The tool accepts the following commands from the user (the underlined letter indicates the key required to activate the command, also note that the tool only responds to user commands at generation boundaries) :



*Figure 18 - GRTOS-GA Startup Screen*

- Evolve - Starts (and stops) the evolution process. After receiving this command the tool will switch between the startup screen (see *Figure 18*) and begin displaying the evolution statistics graphs.
- Pause - Pauses (and restarts) the evolution process at the next generation boundary.

- Grab image - Grabs the currently displayed image (in MS Windows Bit Map format) to the file `IMAGE_XX.BMP`, where `XX` is a sequence number. The file is saved in the directory specified in the INI file.
- Reverse colors - Reverses the color scheme of the display. Both the normal and reversed color schemes can be specified in the INI file.
- Make report - Makes a report file (`GA_XX.RPT`, where `XX` is a sequence number) in the same directory as the tool executable file. The report generated represents the current genetic algorithm statistics. Note that a report file is always generated when the program is terminated.
- Quit - Terminates the evolution tool (and produces a report file).



## 5. Test Suite Description

In order to evaluate the tool set it is necessary to use the tools to solve increasingly complex problems. As these problems are solved, confidence in the ability of the tools to provide correct solutions will increase. The real-time applications that were used to evaluate the tool set are described in the following sections and are of increasing complexity. Obviously, the ability of the tool to provide correct solutions to problems with known solutions is the first step in validating its applicability to arbitrary real-time applications.

A partial formal analysis of these problems is also possible and provides insight into the schedulability of the tasks that make up the problem. This analysis does not help in determining the solution but does determine if a solution exists which can satisfy the problem requirements. This schedulability test is defined in *Equation 5* below [Ref. 1, pg. 28]. If the inequality is met, then there is guaranteed to exist an RTOS configuration which will satisfy the problem requirements. If the inequality is not met then it is guaranteed that an RTOS configuration that satisfies the problem requirements does not exist.

$$\sum_{i=1}^N \frac{C_i}{T_i} \leq N \cdot (2^{\frac{1}{N}} - 1)$$

Where:  $N$  is the number of tasks

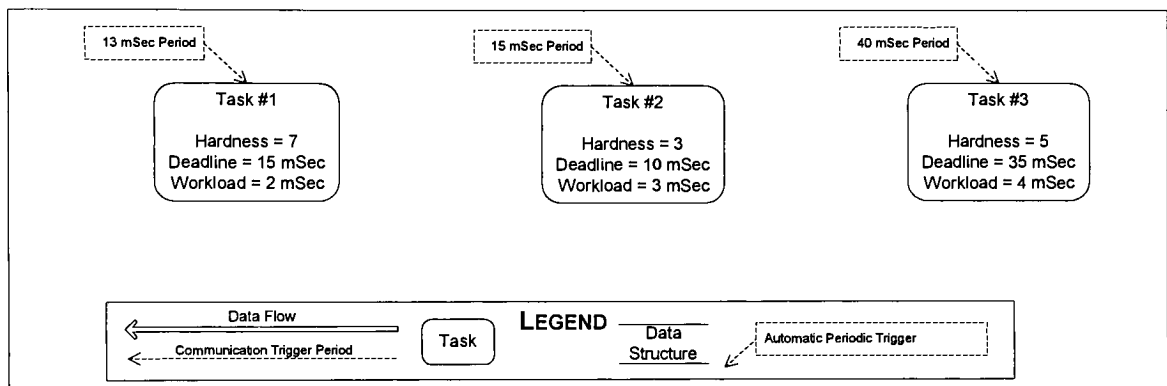
$C_i$  is the execution time of task  $i$

$T_i$  is the period of task  $i$

*Equation 5 - Task Schedulability*

## 5.1 Verification of Capabilities

The first step in verifying that the GA is capable of finding a reasonable solution to a real-time application is to use the GA to solve a problem with a known solution. The least complicated class of real-time applications is the class of purely independent, periodic tasks. The following figure (*Figure 19 - Real-Time Problem #1*) completely describes one such real-time application.



*Figure 19 - Real-Time Problem #1*

Note that none of the tasks interact in any way, this is an example of the simplest class of real-time problems. Applying the schedulability test of *Equation 5* to this problem yields the following results:

$$\frac{2}{13} + \frac{3}{15} + \frac{4}{40} \leq 3 \cdot (2^{\frac{1}{3}} - 1)$$

$$45.4\% \leq 77.9\% - \text{Passed}$$

*Equation 6 - Schedulability, Problem #1*

From this analysis it is clear that under all conditions, the tasks are schedulable, i.e., there is guaranteed to exist a task schedule (RTOS configuration) which will allow all of the tasks to meet all requirements at all times. The results from running `GRTOS-GA.EXE` are as follows:

Rank	Fitness	Tasking Model	Timeslice ( $\mu$ Sec)	Priority Inheritance	Priority Allocation	Priority Assignment
1	84.24	Preemptive	1273	Enabled	Rotating	Rate Monotonic
2	84.12	Preemptive	1227	Enabled	Static	Deadline Monotonic
3	83.74	Preemptive	1519	Disabled	Rotating	Random (674788598)
4	82.73	Preemptive	2560	Enabled	Rotating	Rate Monotonic
5	52.26	Cooperative	N/A (34282)	Disabled	Rotating	Rate Monotonic

*Table 5 - GRTOS-GA.EXE Results, Problem #1*

These results are not surprising, for purely independent tasks, an initial priority assignment using a rate monotonic algorithm produces an optimal task scheduling policy. Rate monotonic assignment theory does not provide any insight into the proper choice for the other RTOS configuration parameters. The results above (see *Table 5 - GRTOS-GA.EXE Results, Problem #1*) clearly show that a preemptive multitasking environment using priority inheritance with a timeslice of approximately 1.2 mSec provides the best results. It would be imprudent to read anything more into the results, recall that this tool provides guidelines for configuring the RTOS for the application. For example it appears that rotating priorities may provide some added benefit, although this result may not be conclusive.

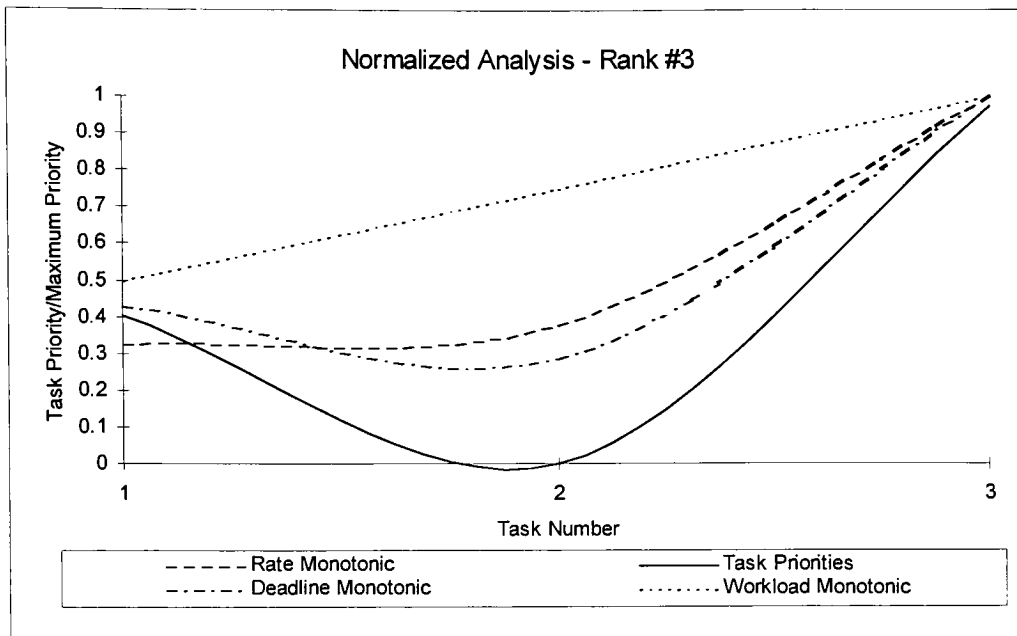


Figure 20 - Problem #1, 'Random' Assignment Analysis - Rank #3

Analyzing the priorities used in the random assignment above (ranked 3<sup>rd</sup>) shows that this random seed actually results in a (albeit non-proportional) deadline monotonic assignment (see *Figure 20 Problem #1, 'Random' Assignment Analysis - Rank #3*, above). It is evident from the fact that the 'Task Priorities' and 'Deadline Monotonic' curves possess the same general characteristics. This supports the conclusion that this 'random' task priority assignment produces priority assignments which are similar to a deadline monotonic assignment scheme.

## 5.2 Verification of Performance.

The next step in the verification process is to use the GA to find a solution to a real-time application for which a solution can be predicted but not proven. This type of problem was created by adding task dependencies to the basic

application. The advantage of using the GA to solve this type of problem is that it is relatively straight forward to see that the GA arrived at a reasonable solution.

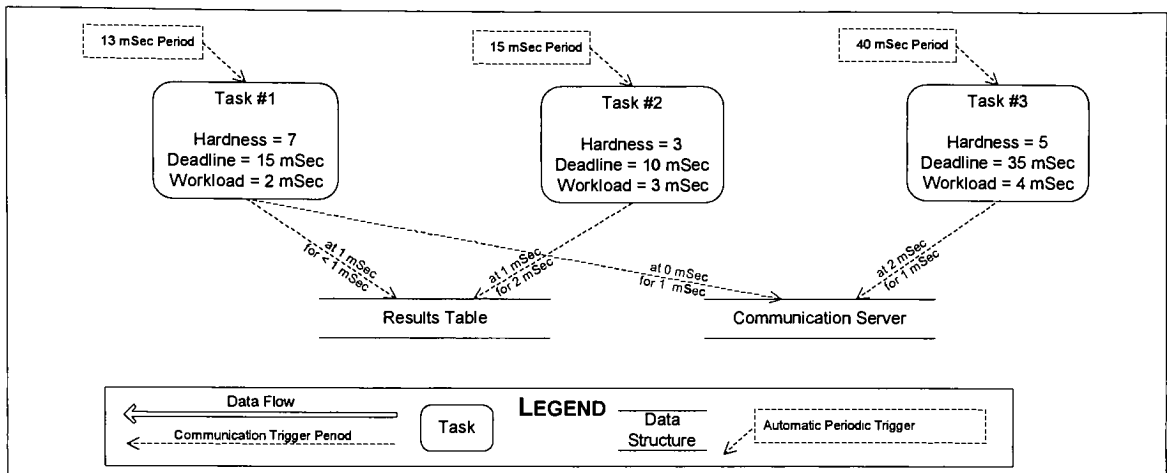


Figure 21 - Real-Time Problem #2

In the cases where task blocking can occur (as in this case) the schedulability test of Equation 5 must be extended to include blocking time in addition to execution time (in other words,  $C_i$  is replaced by  $C_i + B_i$ , where  $B_i$  is the total amount of time that all other tasks can lock resources used by the task). This analysis represents the worst case task schedule. The actual task schedule used may eliminate some (or all) of the blocking times. The following analysis shows the upper and lower bounds for the task schedulability:

$$\frac{2}{13} + \frac{3}{15} + \frac{4}{40} \leq 3 \cdot (2^{\frac{1}{3}} - 1)$$

Lower Bound = 45.4% ≤ 77.9% - Passed

$$\frac{2+3}{13} + \frac{3+1}{15} + \frac{4+1}{40} \leq 3 \cdot (2^{\frac{1}{3}} - 1)$$

Upper Bound = 77.6% ≤ 77.9% - Passed

Equation 7 - Schedulability, Problem #2

It is obvious that even under the worst case task scheduling (when all blocking time is included), the tasks are still schedulable, albeit barely. The results from running `GRTOS-GA.EXE` are as follows:

Rank	Fitness	Tasking Model	Timeslice ( $\mu$ Sec)	Priority Inheritance	Priority Allocation	Priority Assignment
1	76.68	Preemptive	1205	Enabled	Rotating	Rate Monotonic
2	75.81	Preemptive	1245	Enabled	Rotating	Rate Monotonic
3	74.41	Preemptive	1245	Disabled	Rotating	Rate Monotonic
4	72.78	Preemptive	1245	Enabled	Rotating	Rate Monotonic
5	54.67	Cooperative	N/A (33299)	Enabled	Static	Rate Monotonic

*Table 6 GRTOS-GA.EXE Results, Problem #2*

Again the results are not surprising, given that a schedule was possible under all conditions it is reasonable to expect that a rate monotonic assignment would perform best. It is interesting to note that the above results (see *Table 6 – GRTOS-GA.EXE Results, Problem #2*) clearly show results that are very similar to the previous problem (see *Table 5 – GRTOS-GA.EXE Results, Problem #1*). That is, a preemptive multitasking environment using priority inheritance with a timeslice of approximately 1.2 mSec provides the best results. These results point more strongly to the fact that rotating priorities may provide some added benefit, in fact that conclusion can clearly be made.

### **5.3 Test of Capabilities**

Now that the GA has been verified to operate correctly, the next test is a simple problem for which there is no obvious solution. This represents the first

opportunity to use the GA to find a solution for a real-time problem when the only other means of analysis is conjecture. The problem has been created by further augmenting the basic problem to include sporadic tasks modeled as periodic tasks.

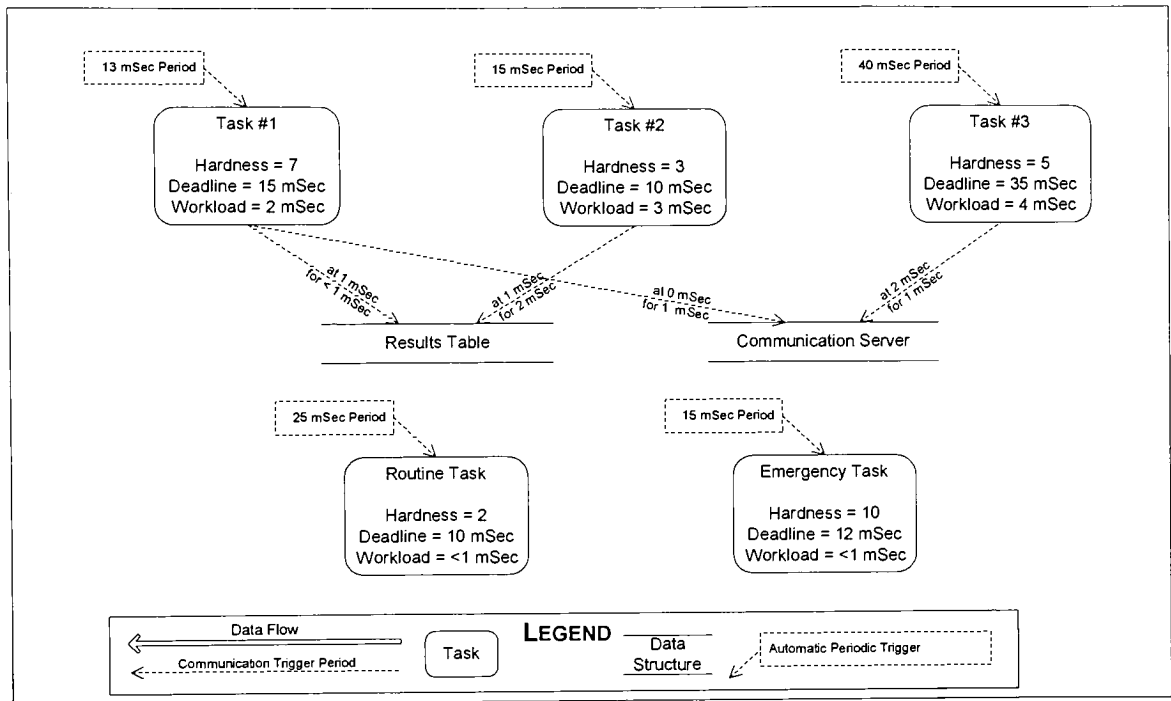


Figure 22 Real-Time Problem #3

As in the previous case, task blocking can occur, therefore the schedulability test of Equation 5 is again extended to include blocking time in addition to execution time. The following shows the analysis:

$$\frac{2}{13} + \frac{3}{15} + \frac{4}{40} + \frac{1}{25} + \frac{1}{15} \leq 5 \cdot (2^{\frac{1}{5}} - 1)$$

Lower Bound = 56.1% ≤ 74.3% - Passed

$$\frac{2+3}{13} + \frac{3+1}{15} + \frac{4+1}{40} + \frac{1+0}{25} + \frac{1+0}{15} \leq 5 \cdot (2^{\frac{1}{5}} - 1)$$

Upper Bound = 88.3% ≤ 74.3% - Failed

Equation 8 - Schedulability, Problem #3

As this analysis clearly shows, this problem (i.e. the set of tasks) may or may not be schedulable. Whether the tasks are schedulable or not is a function of the RTOS configuration itself. This is evident from the fact that the lower bound on the task schedulability is below the threshold and the upper bound is above it. This means that under the worst possible scenario (i.e., RTOS configuration) the tasks are not schedulable. The obvious conclusion is that it is likely that an RTOS configuration exists which results in some level of blocking which yields a CPU utilization which is below the allowable threshold.

The results from running `GRTOS-GA.EXE` are as follows:

Rank	Fitness	Tasking Model	Timeslice ( $\mu$ Sec)	Priority Inheritance	Priority Allocation	Priority Assignment
1	71.86	Preemptive	1005	Enabled	Rotating	Deadline Monotonic
2	71.01	Cooperative	N/A (26878)	Disabled	Static	Random (-1704415732)
3	69.85	Cooperative	N/A (24732)	Enabled	Static	Random (1366122493)
4	62.83	Preemptive	1879	Enabled	Rotating	Deadline Monotonic
5	61.61	Preemptive	44703	Disabled	Static	Random (712905428)

*Table 7 - GRTOS-GA.EXE Results, Problem #3*

Analyzing the characteristics of the ‘random’ priority assignments for this real-time application is very difficult. As the following graphs show (*Figure 23*, *Figure 24* and *Figure 25*), the task assignments that result from the random assignment algorithms cannot accurately be categorized.



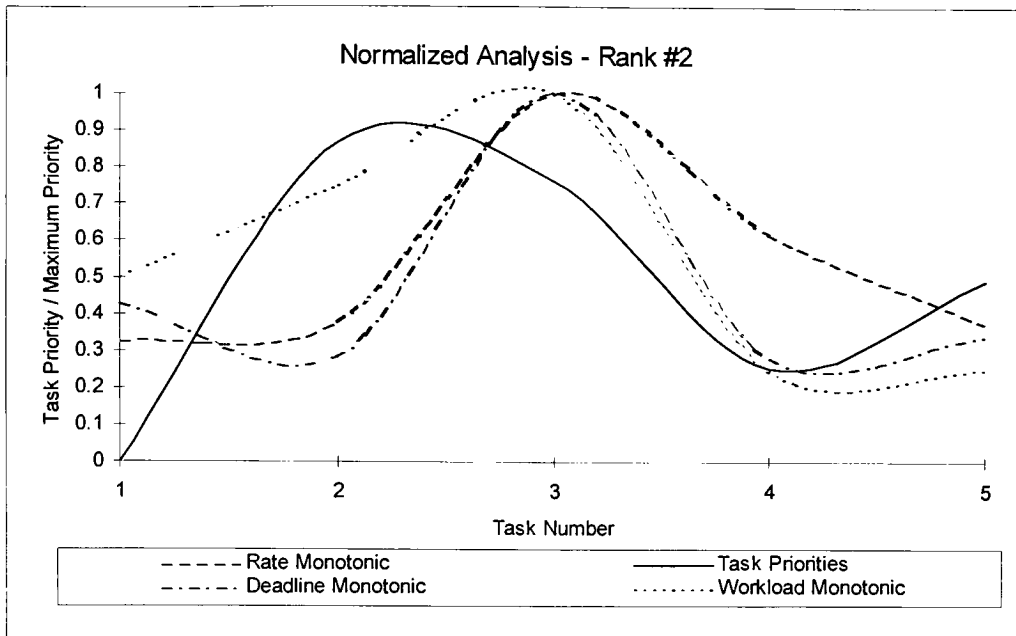


Figure 23 Problem #3, 'Random' Assignment Analysis - Rank #2

As can be clearly seen, the 'Task Priorities' curve does not resemble any of the other curves. This random task priority assignment does not share any general characteristics with the 'standard' task priority assignment algorithms. Therefore, this genotype cannot help in determining general trends for the priority assignment gene. The simple fact is that this random priority assignment outperformed most other priority assignment algorithms and cannot be described as anything other than random.

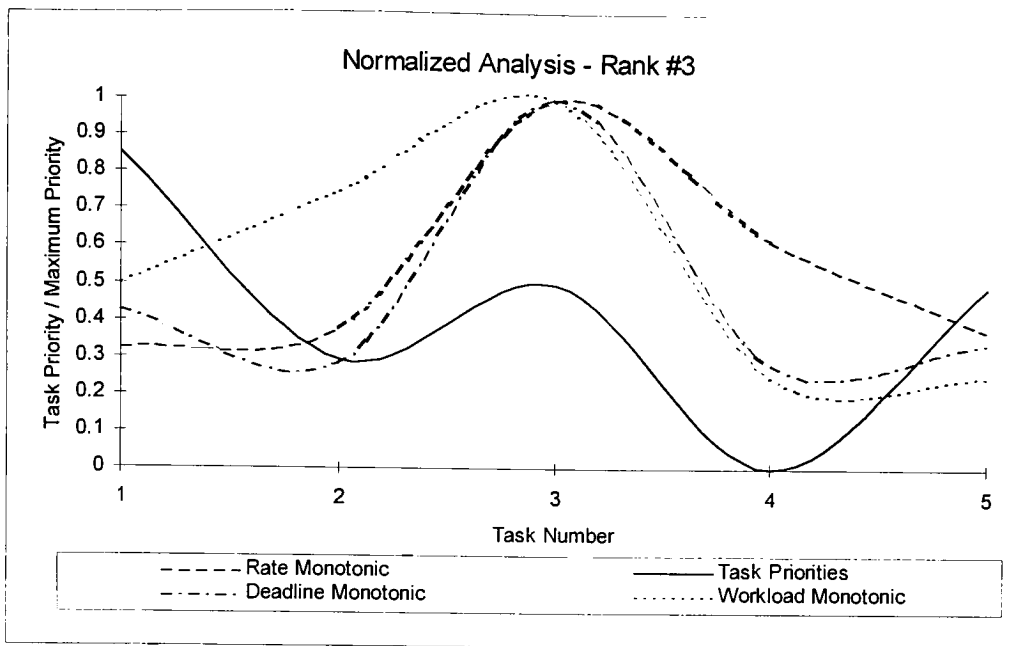


Figure 24 Problem #3, 'Random' Assignment Analysis - Rank #3

Although it is not obvious, the 'Task Priorities' curve above shares many common characteristics with the 'Deadline Monotonic' curve. The task priorities for this random assignment are more exaggerated than a pure deadline monotonic assignment but the relation between any two task priorities is the same. That is, for deadline monotonic assignment and the random assignment in this case, the following is true:

$$\forall \langle d_1, d_2 \rangle : d_1, d_2 \in [\text{SetOfDeadlineMonotonicPriorities}], \forall \langle r_1, r_2 \rangle : r_1, r_2 \in [\text{SetOfRandomPriorities}] \\ \longrightarrow [(d_1 < d_2 \longleftrightarrow r_1 < r_2) \wedge (d_1 > d_2 \longleftrightarrow r_1 > r_2) \wedge (d_1 = d_2 \longleftrightarrow r_1 = r_2)]$$

Equation 9 - Priority Assignment Characteristics

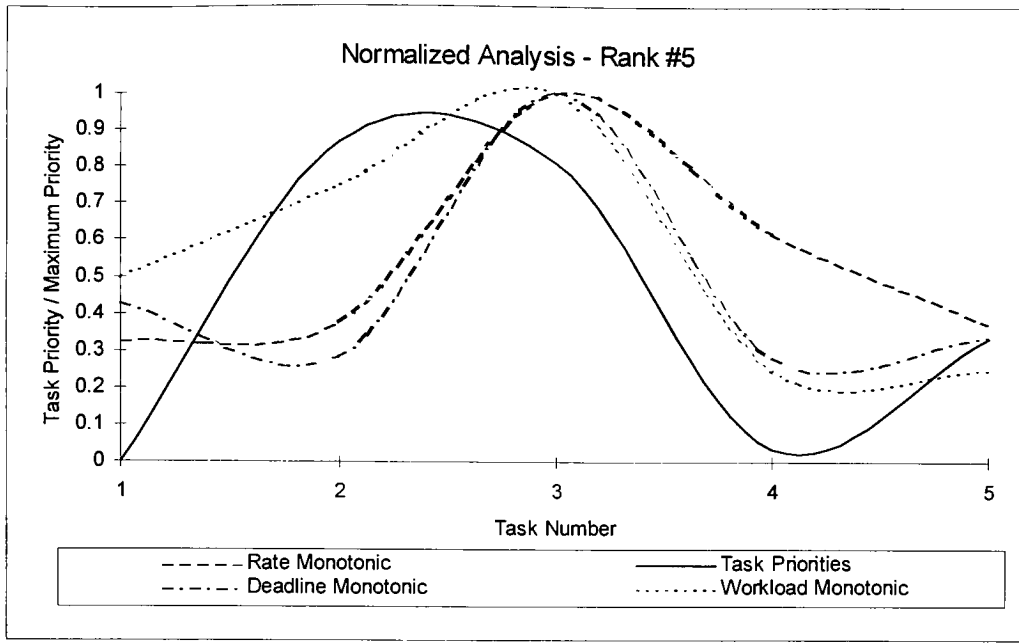


Figure 25 - Problem #3, 'Random' Assignment Analysis - Rank #5

Again, as can be clearly seen, the 'Task Priorities' curve does not entirely resemble any of the other curves. This random task priority assignment does share some general characteristics with the 'Workload Monotonic' curve but, that relationship is very weak. This weak association cannot accurately be used in determining general trends for the priority assignment gene.

From the results obtained it can be concluded that Deadline Monotonic task priority assignment yields the best results. This is, at least in part, surprising, since even worst case blocking analysis results in a nearly schedulable set of tasks (see *Equation 8 - Schedulability, Problem #3*). It is very promising that the genetic algorithm used here produced results that outperformed all others and would likely not have been derived by other means.

## **5.4 Test of Performance**

The last stage in evaluating the GA was to use it to find a solution to a practical, complex problem. The important aspect to be evaluated in this case is the ability of the GA to converge to a solution for a complicated problem. With a practical complex problem it is very important that the GA be tested to ensure that it will be capable of converging. The actual real-time application used for this phase of the testing involves the operation and control of a complex digital radio transmitter.

This embedded, real-time application encompasses the main control of the digital radio. The primary functions of the application are: provide the user interface to the transmitter; monitor the operation of the transmitter to prevent dangerous operating condition from arising; and perform the necessary control system functions to maintain the transmitter output power within a specified limit. To accomplish these functions the transmitter software is broken down into a number of distinct but interrelated tasks. The responsibility of each of the tasks is listed below:

- Keypad Scan - The operator is able to control the transmission parameters of the radio (e.g., frequency, modulation mode, audio input source selection, data input source selection, etc.). This task is responsible for providing the operator with a means of interacting with the radio. The operator input is via an alpha-numeric keypad. This task is responsible for determining which key(s) are pressed by the operator and translating those keys pressed into internal radio commands.

- Front Panel Display - The operator is informed as to the status of the transmitter operation at all times (e.g., output power level, input audio level, fault messages, etc.). This task is responsible for updating the front panel liquid crystal display (LCD) with the present state of the radio transmitter.
- Remote Control - The transmitter is capable of being remotely controlled via a serial interface. This task is responsible for accepting serial data which represents operator commands (with the same controllable parameters as from the operator keypad). It is also responsible for providing radio status over the serial interface by responding to operator queries (with the same status parameters as on the LCD display).
- Emergency Monitor - The transmitter operation is monitored to ensure that dangerous signal levels and/or temperatures are not building up within the radio. This task is responsible for monitoring the internal radio temperature, voltage and current levels to ensure continuous safe operation.
- Analog Sampler - In order to provide control information to the rest of the radio processing software, this task is responsible for reading all analog to digital converters and providing those measurements in useful units of measure to the rest of the radio control software.
- Power Control - The output power of the transmitter is kept within specified limit during transmission. This is accomplished regardless of

the level of input audio to the transmitter. This task is responsible for adjusting the input audio level and the transmitter power gain to ensure that the output of the transmitter stays within limits.

- System State Control - The transmitter hardware is complex and in many cases mechanical in nature, this task is responsible for the absolute control of the various motors and relays used to tune and adjust the transmitter.
- Keyline Processing - Once the operator has configured the transmitter with the desired transmission parameters the system can be 'keyed' to begin transmission. The keyline is typically connected to the radio handset (i.e., 'Push-To-Talk' keyline) and causes the handset audio to begin being modulated over the air by the radio. This task is responsible for monitoring the selected key source and translating the actuation and release of the keyline into the appropriate commands to the radio hardware to start and terminate transmission.

The system design of the transmitter software is given in *Figure 26 - Real-Time Problem #4*, below.

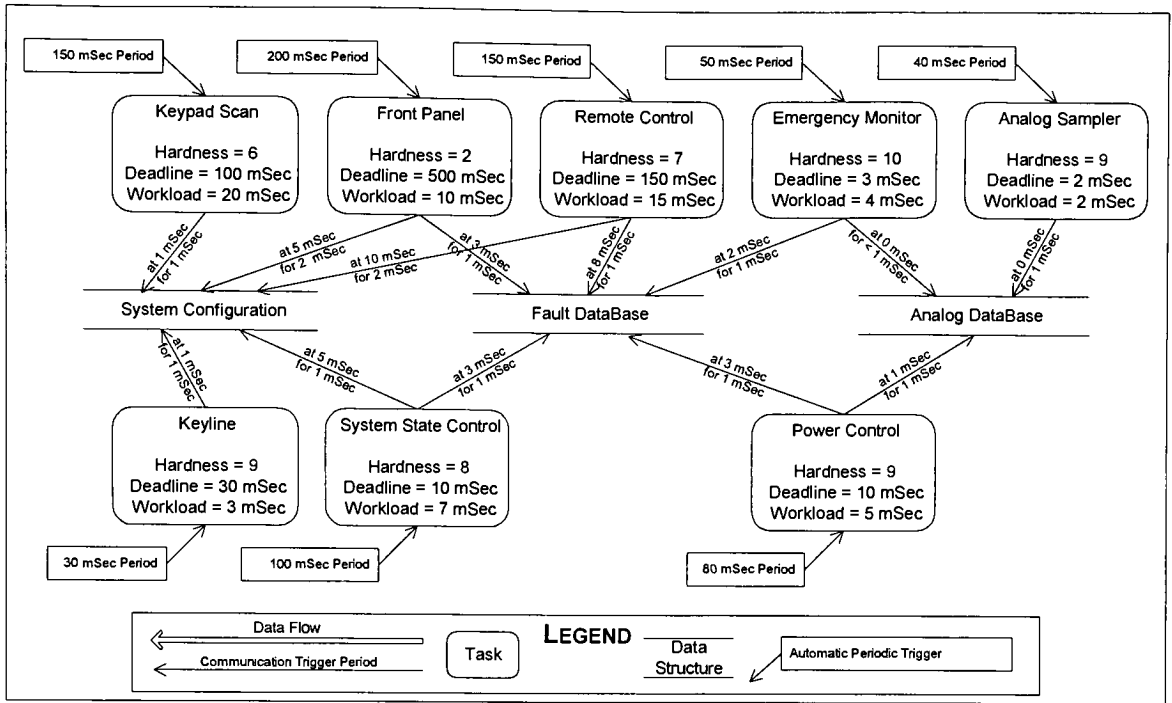


Figure 26 - Real-Time Problem #4

As in the previous case, task blocking can occur, therefore the schedulability test of Equation 5 is again extended to include blocking time in addition to execution time. The following shows the analysis:

$$\frac{20}{150} + \frac{10}{200} + \frac{15}{150} + \frac{4}{50} + \frac{2}{40} + \frac{5}{80} + \frac{7}{100} + \frac{3}{30} \leq 8 \cdot (2^{\frac{1}{8}} - 1)$$

Lower Bound = 64.6% ≤ 72.4% - Passed (barely)

$$\frac{20+6}{150} + \frac{10+9}{200} + \frac{15+9}{150} + \frac{4+6}{50} + \frac{2+2}{40} + \frac{5+6}{80} + \frac{7+10}{100} + \frac{3+6}{30} \leq 8 \cdot (2^{\frac{1}{8}} - 1)$$

Upper Bound = 133.6% ≤ 72.4% - Failed

Equation 10 - Schedulability, Problem #4

As this analysis clearly shows, this problem (i.e. the set of tasks) may be schedulable, but very likely is not. Whether the tasks are schedulable or not is a function of the RTOS configuration itself. Note that there is only a small margin between the utilization lower bound and the schedulability threshold. The obvious conclusion is that it is not likely that an RTOS configuration exists

which results in some level of blocking which yields a CPU utilization which is below the allowable threshold.

The results from running `GRTOS-GA.EXE` are as follows:

Rank	Fitness	Tasking Model	Timeslice ( $\mu$ Sec)	Priority Inheritance	Priority Allocation	Priority Assignment
1	76.02	Preemptive	1227	Disabled	Static	Deadline Monotonic
2	75.94	Preemptive	2001	Disabled	Static	Deadline Monotonic
3	75.69	Preemptive	21060	Enabled	Rotating	Random (-1675281081)
4	73.57	Preemptive	2166	Disabled	Rotating	Deadline Monotonic
5	72.71	Cooperative	N/A (47447)	Enabled	Rotating	Random (-4861458)

*Table 8 - GRTOS-GA.EXE Results, Problem #4*

Analyzing the characteristics of the ‘random’ priority assignments for this real-time application is somewhat difficult. As the following graphs show (*Figure 27* and *Figure 28*), one of the task assignments that result from the random assignment algorithms can be accurately be categorized, and the other cannot.



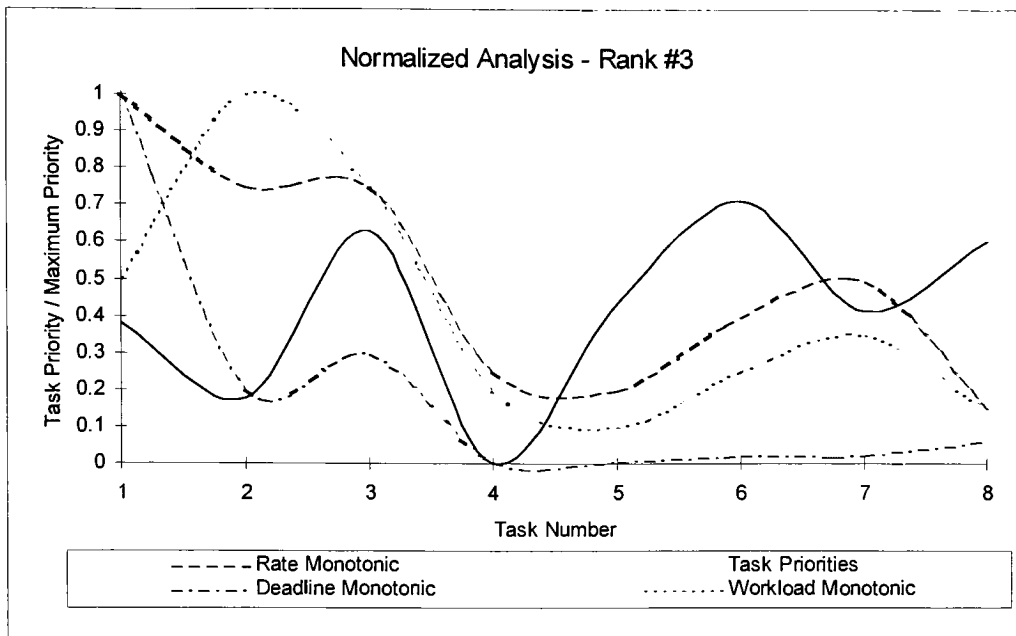


Figure 27 - Problem #4, 'Random' Assignment Analysis - Rank #3

Analyzing the priorities used in the random assignment above (ranked 3<sup>rd</sup>) shows that this random seed actually results in a (somewhat exaggerated) deadline monotonic assignment (see *Figure 27 - Problem #4, 'Random' Assignment Analysis - Rank #3*, above). Although not obvious, the 'Task Priorities' and 'Deadline Monotonic' curves possess the same general characteristics, again the relationships between task priorities are the same (see *Equation 9 - Priority Assignment Characteristics*). This supports the conclusion that this 'random' task priority assignment produces priority assignments which are similar to a deadline monotonic assignment scheme.

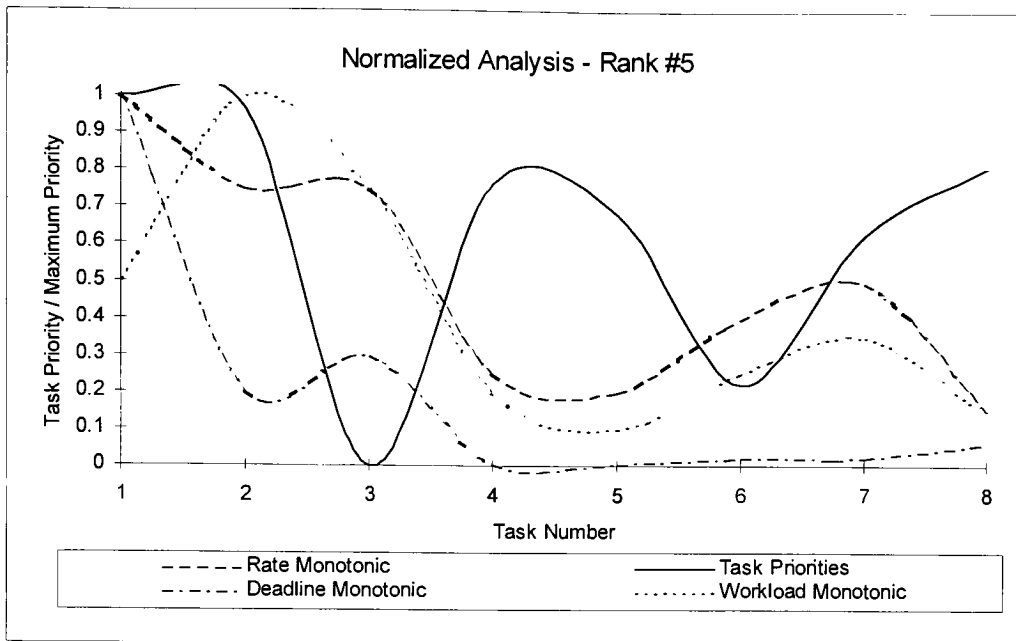


Figure 28 - Problem #4, 'Random' Assignment Analysis - Rank #5

As can be clearly seen, this 'Task Priorities' curve does not share any general characteristics with the 'standard' task priority assignment curves. There are simply no conclusions that can be drawn from this assignment with regard to priority assignment. The simple fact is that this random priority assignment outperformed many other priority assignment algorithms and cannot be described as anything other than random.

The results above (see *Table 8 - GRTOS-GA.EXE Results, Problem #4* and *Figure 27 - Problem #4, 'Random' Assignment Analysis - Rank #3*) clearly show that a preemptive multitasking environment, without priority inheritance, using deadline monotonic priority assignment and a timeslice of approximately 2 mSec provides the best results. The genetic algorithm produced these results for a problem specification which could not otherwise be solved.

## 6. Conclusions

From the results of the previous section it is reasonable to conclude that the genetic algorithm developed in this thesis can reliably be used to determine the ‘best’ real-time operating system configuration for an arbitrary real-time application. The tests performed verify that the GA converges to the correct solution in a case where the optimal solution is known. In addition, the GA converges to a reasonable solution in a case where the optimal solution can be predicted but not proven.

What is most interesting about the results obtained is that the task priority assignment algorithms which produced the best fitness for the problems chosen clearly fell into two categories. The problem specifications for ‘schedulable’ real-time problems resulted in rate monotonic assignment algorithms producing the best fitness values. This is not surprising since rate monotonic theory proves that this result produces an optimal solution. The problem specifications for ‘non-schedulable’ real-time problems resulted in deadline monotonic assignment algorithms producing the best fitness values. This is somewhat surprising but the data clearly supports this conclusion.

These results immediately raise the question: ‘Does deadline monotonic task priority assignment produce the best results when the application tasks are likely to be non-schedulable?’. This is a very interesting question that would, unfortunately, require further research to answer. The data obtained very clearly implies this conclusion but the lack of a large number of diverse real-time

application trials prevents that conclusion from being drawn. Opening this topic to additional research is probably the most significant result of this thesis.

There is an additional observation that must also be made. The most common industry practice when defining the RTOS configuration for a particular system is, by far, to use cooperative multitasking with uniform task priorities (this opinion is based on years of experience developing and studying fielded real-time systems). I believe that the reason for the choice of each of these configuration parameters is entirely different. As for the multitasking model, the reason that cooperative multitasking is often chosen is that cooperative multitasking systems are inherently simpler than preemptive multitasking systems. The argument for choosing uniform task priority assignment is that there is often little knowledge of the task execution profiles, when this is the case other priority assignment algorithms are not feasible.

The results of this thesis clearly show that preemptive multitasking is superior to cooperative multitasking in almost all situations, the added complexity of preemptive multitasking is more than outweighed by the performance improvement achieved. Also clearly evident from the results of this thesis is that uniform priority assignment always produces lower performance than any other assignment algorithm. Again, it is clear that analysis of the system design (i.e. task characteristics) is essential in order to determine the ‘best’ initial priority assignment algorithm for the real-time system.

## 7. Bibliography

- Ref. 1 J. T. Baldwin, Predicting and Estimating Real-Time Performance, Embedded Systems Programming, pg 30-41, February 1995.
- Ref. 2 S. Bodilsen, Scheduling Theory and Ada 9X, Embedded Systems Programming, pg 32-52, December 1994.
- Ref. 3 M. F. Bramlette and E. E. Bouchard, Genetic Algorithms in Parametric Design of Aircraft, Handbook of Genetic Algorithms, Van Nostrand Reinhold, 1991.
- Ref. 4 I. L. Bukatova and Y. V. Gulyaev, From Genetic Algorithms to Evolutionary Computer, Proceedings of the Fifth International Conference on Genetic Algorithms, pg 614-617, 1993.
- Ref. 5 C. Darwin, The Origin Of Species, Mentor Edition, NAL Penguin Publishing, 1958.
- Ref. 6 L. Davis, Handbook of Genetic Algorithms, Van Nostrand Reinhold, 1991.
- Ref. 7 K. DeJong and W. Spears, On the State of Evolutionary Computation, Proceedings of the Fifth International Conference on Genetic Algorithms, pg 618-623, 1993.
- Ref. 8 K. Ellison, Scheduling Algorithms for Systems with Hard Deadlines, Embedded Systems Programming, pg 24-31, April 1995.
- Ref. 9 C. Heitmeyer and D. Mandrioli, Formal Methods for Real-Time Computing: An Overview, Formal Methods for Real-Time Computing, John Wiley & Sons, pg 1-32, 1996.
- Ref. 10 E. D. Jensen, Eliminating the Hard/Soft Real-Time Dichotomy, Embedded Systems Programming, pg 28-33, October 1994.
- Ref. 11 I. Lee, H. Ben-Abdallah and J-Y. Choi, A Process Algebraic Method for the Specification and Analysis of Real-Time Systems, Formal Methods for Real-Time Computing, John Wiley & Sons, pg 167-194, 1996.

- Ref. 12 Z. Michalewicz, Genetic Algorithms + Data Structures = Evolution Programs, Springer-Verlag, 1992.
- Ref. 13 M. Milenkovic, Operating Systems: Concepts and Design, McGraw-Hill, Inc., 1987.
- Ref. 14 A. K. Mok, D. A. Stuart, and F. Jahanian, Specification and Analysis of Real-Time Systems: Modechart Language and Toolset, Formal Methods for Real-Time Computing, John Wiley & Sons, pg 33-54, 1996.
- Ref. 15 R. Obenza, Guaranteeing Real-Time Performance Using RMA, Embedded Systems Programming, pg 26-40, May 1994.
- Ref. 16 A. V. Oppenheim and A. S. Willsky, Signals and Systems, Prentice-Hall, 1983.
- Ref. 17 N. J. Radcliff and F. A. W. George, A Study in Set Recombination Proceedings of the Fifth International Conference on Genetic Algorithms, pg 23-30, 1993.
- Ref. 18 A. Tanenbaum, Modern Operating Systems, Prentice-Hall, 1992.
- Ref. 19 D. M. Tate and A. E. Smith, Expected Allele Coverage and the Role of Mutation in Genetic Algorithms, Proceedings of the Fifth International Conference on Genetic Algorithms, pg 31-37, 1993.



```

HOURS      : 0.23;
MINUTES    : 0.59;
SECONDS    : 0.59;
MILLISECONDS : 0.999;
end; (* TIME *)

.....
(* This is the single data structure that must be configured for each
 * application task.
 * .....
TASK_ATTRIBUTES = record
.....
(* This priority is subject to the 'Priority Scheduling Policy'
 * established in the USER_CONFIGURATION data structure. Note that
 * higher numbers are higher priority.
 * .....
PRIORITY : USER_PRIORITIES;
.....
(* This is the number of words that the task requires for stack
 * space. There is no easy way to determine this value accurately.
 * Turning on the TASK_STATISTICS is the best method of determining
 * stack requirements. Note that if the application 'locks up' it
 * is very likely that one (or more) of the task stacks is not big
 * enough.
 * .....
STACK_WORDS_NEEDED : word;
.....
(* If an application does not assign error handlers then TASKING
 * will use internal default handlers which provide some diagnostics
 * and attempt to allow the application to continue running. It is
 * unlikely that the application will be able to run in the presence
 * of errors, but all attempts are made to do so. If the default
 * error handlers are to be used then the corresponding
 * ERROR_HANDLERS element should be set to 'nil'.
 * .....
ERROR_HANDLERS : array (SYSTEM_ERRORS) of HANDLER_PROCEDURES;
end; (* TASK_ATTRIBUTES *)

.....
(*
 * Counting 'Dijkstra' Semaphores
 * .....
(* All semaphores used by the application must be variables declared
 * (or dynamically created) of this type. 'Dijkstra' semaphores can
 * retain signals because information is maintained so that TASKING
 * knows how many signals are held by the semaphores at all times.
 * .....
SEMAPHORE = 0..$FFFF;

.....
(*
 * Software Events
 * .....
(* All events used by the application must be variables declared (or
 * dynamically created) of this type. Events can retain only a single
 * signal. All signals to be signaled all subsequent signals
 * are lost (i.e., the event remains signaled).
 * .....
EVENT = (UNSIGNALLED, SIGNALLED);

.....
(*
 * Binary Semaphores
 * .....
(* All binary semaphores used by the application must be variables
 * declared (or dynamically created) of this type. Binary semaphores
 * are used to ensure mutually exclusive access to a resource which is
 * shared between tasks. It is logically an error to signal a
 * signaled binary semaphore, TASKING will detect this error as an
 * illegal operation.
 * .....
BINARY_SEMAPHORE = 0..1;

.....
(*
 * Condition Variables
 * .....
(* All condition variables used by the application must be variables
 * declared (or dynamically created) of this type. It is not likely
 * that applications will use condition variables (they are most
 * useful in the implementation of Hoare Monitors). Condition
 * variables do not have any memory associated with them, only if a
 * task is already waiting on a condition variable when it is signaled
 * will that task become unblocked. If a condition variable is
 * signaled and there is no task waiting on it then the signal is lost.
 * .....
CONDITION_VARIABLE = (CONDITION_VARIABLES_DO_NOT_HAVE_VALUES);

.....
(*
 * Memory Management Operations
 * .....
procedure GETMEM
.....
(* This procedure replaces the standard Turbo Pascal memory allocation
 * procedure of the same name. Since TASKING must perform all memory
 * allocation and de-allocation, this routine 'must' be used instead of the
 * system.getmem() procedure. If 'type P : T;' is declared, then the
 * correct usage is 'getmem(pointer(P), sizeof(T))'.
 * .....
procedure FREEMEM
.....
(* This procedure replaces the standard Turbo Pascal memory allocation
 * procedure of the same name. Since TASKING must perform all memory
 * allocation and de-allocation, this routine 'must' be used instead of the
 * system.freemem() procedure. If 'type P : T;' is declared, then the
 * correct usage is 'freemem(pointer(P), sizeof(T))'.
 * .....
.....
(*
 * General Task Management Operations
 * .....
procedure CREATE
.....
(* This routine creates all system data structures for a parallel
 * executing task. The task is made ready to run and will be available for
 * execution. Actual execution will occur when the scheduler activates the
 * task (i.e. precise execution time is unknown and depends on availability
 * of system resources and system load). In addition NO tasks will be
 * scheduled until the application 'main' has terminated.
 * .....
procedure DESTROY
.....
(* This routine terminates and (if appropriate) releases all of the task's
 * resources to be used by other tasks and then removes all evidence of the
 * task ever having existed. Tasks may only 'destroy' other tasks, not
 * themselves. This operation should only be performed when the full
 * consequences of task destruction are known and the result of the task
 * .....
procedure SUSPEND
.....
(TASK : TASK_IDS);
.....
(* This routine causes a previously running or ready task to be blocked
 * unconditionally. The only way that a task suspended by this call can run
 * again is if another task RESUME()s it.
 * .....
procedure RESUME
.....
(TASK : TASK_IDS);
.....
(* This routine causes a previously suspended task to be made ready. The
 * task is not necessarily the next to run, that is based on the priorities
 * .....
function GET_MILLISECOND_TICKS : longint;
.....
(* This routine returns the number of milliseconds that have elapsed since
 * the program started executing.
 * .....
procedure WAIT_FOR_DELAY
.....
(DURATION : TIME);
.....
(* This system call causes the currently running task to become blocked
 * for the specified length of time. There is no guarantee that the task
 * will begin executing when the time expires, the only guarantee is that
 * the task will become ready to execute at that time.
 * .....
procedure PREEMPTABLE_DELAY
.....
(DURATION : TIME);
.....
(* This system call causes the currently running task to be delayed by the
 * specified length of time.
 * .....
procedure CHANGE_PRIORITY
.....
(PRIORITY : USER_PRIORITIES);
.....
(* This routine simply changes the priority of the currently running task
 * and re-schedules the running task (if there is a new highest
 * priority task ready to run).
 * .....
.....
(*
 * Counting 'Dijkstra' Semaphore Operations
 * .....
procedure SIGNAL_SEMAPHORE
.....
(* This routine performs an Up() operation on the specified semaphore. If
 * there are tasks waiting on the semaphore then one is awakened. If there
 * are none waiting then the semaphore is incremented.
 * .....
(*
 * Binary Semaphore Operations
 * .....
procedure WAIT_ON_SEMAPHORE
.....
(* This routine performs a Down() operation on the specified semaphore. If
 * the semaphore is zero then the running task is suspended. If it is
 * non-zero then the semaphore is decremented.
 * .....
.....
(*
 * Software Event Operations
 * .....
procedure SIGNAL_EVENT
.....
(* This routine signals the event specified. If there are tasks waiting on
 * the event then one task is made ready. If there are no tasks waiting on
 * the event then the signal is saved, only one signal is maintained no
 * matter how many times the event is signaled.
 * .....
procedure BROADCAST_EVENT
.....
(* This routine signals the event specified. If there are tasks waiting on
 * the event then they are ALL made ready. If there are no tasks waiting on
 * the event then the signal is saved and the next SINGLE task to wait on
 * the event will consume the event. Only one signal is maintained no
 * matter how many times the event is signaled.
 * .....
procedure START_PERIODIC_EVENT
.....
(* This routine starts the periodic signalling (by TASKING) of the event
 * specified at the specified interval. This routine cannot cause a task to
 * be suspended but a re-schedule may occur.
 * .....
procedure STOP_PERIODIC_EVENT
.....
(* This routine stops the periodic signalling (by TASKING) of the event
 * specified. This routine cannot cause a task to be suspended but a
 * re-schedule may occur.
 * .....
procedure WAIT_ON_EVENT
.....
(* This routine causes the calling task to be suspended waiting for the
 * specified event to occur. If the event has already occurred then the
 * calling task will immediately become ready, although a re-schedule may
 * occur.
 * .....
.....
(*
 * Binary Semaphore Operations
 * .....
procedure SIGNAL_BINARY_SEMAPHORE
.....
(* This routine performs an Up() operation on the specified semaphore. If
 * there are tasks waiting on the semaphore then one is awakened. If there
 * are none waiting then the semaphore is set.
 * .....
procedure WAIT_ON_BINARY_SEMAPHORE
.....

```



```

|var SEM : BINARY_SEMAPHORE;
|-----|
|* This routine performs a Down| operation on the specified semaphore. |
|* If the semaphore is locked then the running task is suspended. If it is |
|* non-zero then the semaphore is cleared. |
|*-----|
|*-----|
|* Condition Variable Operations |
|*-----|
procedure SIGNAL_CONDITION_VARIABLE
|-----|
|var C_VAR : CONDITION_VARIABLE;
|-----|
|* This routine signals the condition variable specified, if there is a |
|* task waiting on the condition variable then it is made ready for one of |
|* the multiple waiting tasks as soon as another task signals the |
|* condition variable then the signal is lost. |
|*-----|
procedure WAIT_ON_CONDITION_VARIABLE
|-----|
|var C_VAR : CDNDITION_VARIABLE;
|-----|
|* This routine causes the calling task to become blocked. Because |
|* condition variables are not stored by the system, the calling task is |
|* guaranteed to be blocked as soon as another task signals the |
|* condition variable the waiting task will be made ready for one of the |
|* multiple waiting tasks will be made ready. |
|*-----|
|*-----|
|* Message Passing Operations |
|*-----|
procedure SEND_MESSAGE_TO
|RECIPIENT_TASK : TASK_ID; XMIT_MESSAGE_PTR : pointer;
|-----|
|* This routine sends the message (actually the pointer) to the task |
|* by placing it in the task's mailbox (which is automatically created if |
|* necessary). Although the calling task cannot block while sending the |
|* message it is possible for a reschedule to occur. Typically messages are |
|* dynamically allocated before being sent and the receiver is responsible |
|* for disposing of them after they have been used. |
|*-----|
procedure RECEIVE_MESSAGE
|-----|
|var RCV_MESSAGE_PTR : pointer;
|-----|
|* This routine retrieves a message (actually a pointer to a message) from |
|* the caller's mailbox (which is created if necessary), if there are no |
|* messages in the mailbox then nil is returned. The task cannot become |
|* blocked by calling this routine but a reschedule is possible. Typically |
|* the message will be disposed of after it has been used by the receiving |
|* task. |
|*-----|
procedure WAIT_AND_RECEIVE_MESSAGE
|-----|
|var RCV_MESSAGE_PTR : pointer;
|-----|
|* This routine retrieves a message (actually a pointer to a message) from |
|* the caller's mailbox (which is created if necessary), if there are no |
|* messages in the mailbox then the task is blocked waiting for the mailbox |
|* to become non-empty, because of this the return pointer can never be nil. |
|* Typically the message will be disposed of after it has been used by the |
|* receiving task. |
|*-----|
|*-----|
|* MS-Mouse Handling Operations |
|*-----|
procedure ENABLE_MOUSE_ACTIONS
|-----|
|MOUSE_ACTIONS : word;
|-----|
|* This routine enables the specified MS-Mouse events to be received by |
|* the task which is (or will be) waiting to receive mouse actions. |
|*-----|
procedure WAIT_AND_RECEIVE_MOUSE_ACTIONS
|-----|
|var MOUSE_INFO : MOUSE_PARAMETERS;
|-----|
|* If an MS-Mouse handler is to be created then communication between the |
|* mouse driver and the 'mouse handling' task is done through a variable of |
|* type MS_MOUSE.MOUSE_PARAMETERS (which can be declared or dynamically |
|* created). |
|* This routine receives MS-Mouse action parameters from the MS-Mouse |
|* driver. If none of the enabled mouse actions has occurred then the task |
|* is blocked waiting for one (or more) action to occur. |
|*-----|
|*-----|
|* Keyboard Handling Operations |
|*-----|
function WAIT_ON_READKEY : char;
|-----|
|* This routine provides a means for an application task to perform I/O in |
|* a 'blocking' manner. This means that while there is no input from the |
|* keyboard the calling task is suspended, as soon as input arrives from the |
|* keyboard the caller is awakened and given the ASCII code of the key that |
|* was pressed by the user (i.e., a blocking equivalent of 'readkey'). |
|* An added benefit of using blocked keyboard input is that this version |
|* of 'readkey' is able to detect (and pass back to the application) the |
|* extended function keys F11 & F12 (as well as shift, alt & ctrl versions). |
|* The Turbo Pascal 'readkey' function is unable to recognize F11 & F12. |
|*-----|
|* WARNING: |
|* Some add-on utilities that extended the size of the keyboard buffer |
|* will actually 'steal' the keys from the blocked task. All such |
|* utilities should be disabled when using this blocking I/O feature of |
|* TASKING. |
|* Using blocking I/O causes the application to be unable to distinguish |
|* between multiple characters arriving from the pressing of an ANSI.SYS |
|* re-programmed key on the keyboard from the actual re-programmed key. |
|* That is, if a keyboard key is re-programmed using ANSI.SYS then the |
|* application will never be able to see that key, it will only see the |
|* re-programmed value whenever the key is pressed (the key translation |
|* by ANSI.SYS occurs before the 'key' is detectable by software, although |
|* the standard Turbo Pascal 'readkey' is able to see the 'base key'). |
|* This anomaly can create synchronization problems between the keyboard |
|* and TASKING if the key has been reprogrammed as multiple keys (which is |
|* usually the case). Under these conditions the keyboard buffer can end |
|* up with more characters in it than are reported to the application. |
|* This anomaly can be avoided by disabling the keyboard buffer re-programming |
|* ANSI.SYS re-programming should be disabled when using this blocking I/O |
|* feature of TASKING. |
|*-----|
|*-----|
|* Scheduler Handling Operations |
|*-----|
procedure LOCK_SCHEDULER;
|-----|
|* This procedure 'locks' the scheduler and prevents all further task |
|* rescheduling until UNLOCK_SCHEDULER is called. |
|* One use of this routine is when multiple tasks need to write to the |
|* display and it is desired that the output from each task occur without |
|* the output from other tasks being intermixed with it. During the time |
|* that the scheduler is 'locked' there will be 'no' preemptive/cooperative |
|* task rescheduling. The application should avoid 'locking' the scheduler |
|* if possible, if that is not feasible then the time spent with the |
|* scheduler 'locked' should be an absolute minimum. |
|* It is imperative the 'each' call to LOCK_SCHEDULER have a corresponding |
|* call to UNLOCK_SCHEDULER, successive nested calls to LOCK_SCHEDULER are |
|* allowed as long as each has a corresponding call to UNLOCK_SCHEDULER. |
|*-----|
procedure UNLOCK_SCHEDULER;
|-----|
|* This procedure 'unlocks' the scheduler and allows task rescheduling to |
|* continue based on the previously existing configuration. Note that |
|* LOCK_SCHEDULER must previously have been called. |
|*-----|

```

## 8.2 TASKING.PAS

```
unit TASKING;
(* Extends Turbo Pascal by adding a preemptive or cooperative tasking *)
(* model to the language. The multitasking extensions are fully supported *)
(* while running under MS-DOS. *)
(*-----*)
(* compiler options (Ver. 7.0) *)
(*$A+ Word Alignment *)
(*$B- Short Circuit Boolean Evaluation *)
(*$D+ Debug Code Generation ON (Sort of) *)
(* Requires /V option to TPC to activate *)
(*$L+ Local Debug symbols ON (Sort of) *)
(* Requires /V option to TPC to activate *)
(*$F- Far calls only as necessary *)
(*$G- Generic 80x86 code only *)
(*$I- I/O Checking OFF *)
(*$N- Do Not allow 80x87 instructions *)
(*$O- Overlays NOT allowed *)
(*$P- Standard 'string' parameters *)
(*$Q- Overflow Checking OFF *)
(*$R- Range Checking OFF *)
(*$S- Stack Checking OFF *)
(*$T- Do Not Force Typed '@' references *)
(*$V- Var-string Checking OFF *)
(*$W- Disable Extended syntax *)
(*-----*)

(*$I TASKING.INT - Filename of the unit interface (i.e., 'public' stuff) *)
implementation
uses dos;

const
(*-----*)
(* This is the ID of the TASKING task which is executed when there is *)
(* nothing else to execute (this is reserved and cannot be assigned to *)
(* user tasks). *)
(*-----*)
NULL_TASK_ID = low(TASK_IDS);

(*-----*)
(* This is used to document the code where the task ID parameter to a *)
(* procedure is needed by the compiler but not used by the code *)
(* because of the context of the call. *)
(*-----*)
ANY_TASK = NULL_TASK_ID;

const
(*-----*)
(* These priority declarations allow TASKING to create tasks that are *)
(* guaranteed to have higher and lower priorities than user tasks. *)
(* The obvious use of this is the null task, which has a priority *)
(* lower than all user tasks, and tasks which are executing in a *)
(* scheduler-locked code region which have a priority higher than all *)
(* others. *)
(*-----*)
HIGHEST_PRIORITY = succ(high(USER_PRIORITIES));
LOWEST_PRIORITY = pred(low(USER_PRIORITIES));

type
(*-----*)
(* General task priority type which includes both user priorities and *)
(* TASKING reserved priorities. *)
(*-----*)
PRIORITIES = LOWEST_PRIORITY..HIGHEST_PRIORITY;

const
(*-----*)
(* This variable keeps track of the highest priority when dynamic *)
(* priority rotation is enabled (i.e., dynamic priorities). *)
(*-----*)
HIGHEST_DYNAMIC_PRIORITY : PRIORITIES = HIGHEST_PRIORITY;

type
(*-----*)
(* This list enumerates the allowable priority comparisons (used to *)
(* determine the next highest priority task to execute). *)
(*-----*)
COMPARISONS = (GREATER_THAN, LESS_THAN, EQUAL_TO);

const
(*-----*)
(* This is used to document the code where the address of the block *)
(* parameter to a procedure is needed by the compiler but not used by *)
(* the code because of the context of the call. *)
(*-----*)
THERE_IS_NO_BLOCK_TO_MATCH = nil;

(*-----*)
(* These are the hardware specific constants necessary to program the *)
(* system clock chip (i8254) for TASKING and MS-DOS operation. *)
(*-----*)
TIMER_CLOCK_FREQUENCY = 1193180;
MSDOS_TIMER_0_VALUE = $FF;
ONE_MSEC_TIMER_0_VALUE = round(TIMER_CLOCK_FREQUENCY * 0.001);
DEFAULT_MSDOS_CLOCK_PERIOD = MSDOS_TIMER_0_VALUE / TIMER_CLOCK_FREQUENCY;
TIMER_CONTROL_PORT = 43;
TIMER_0_DATA_PORT = 340;
SYSTEM_TIMER_CONTROL_WORD = 334;

(*-----*)
(* Allows TASKING application to run over a month boundary and report *)
(* the correct execution time. *)
(*-----*)
DAYS_IN : array[1..12] of integer = (31,28,31,30,31,30,31,31,30,31,30,31);

(*-----*)
(* This is used as a special delay value which indicates that the *)
(* task is no longer (or never was) waiting for a delay to expire. *)
(*-----*)
NO_TIME_DELAY = -1;

(*-----*)
(* These values are used to compute the number of milliseconds that a *)
(* task must 'wait' when it requests a WAIT_FOR_DELAY(). *)
(*-----*)
MILLISECONDS_PER_SECOND = 1000;
MILLISECONDS_PER_MINUTE = MILLISECONDS_PER_SECOND * 60;
MILLISECONDS_PER_HOUR = MILLISECONDS_PER_MINUTE * 60;
MILLISECONDS_PER_DAY = MILLISECONDS_PER_HOUR * 24;

(*-----*)
(* This is an interrupt that is vectored from its original vector *)
(* to allow TASKING to monitor/control software interrupts. MS-DOS *)
(* allows user applications to use interrupts 0xFO-0xFF. This allows *)
(* TASKING to prevent preempting MS-DOS (which is guaranteed to crash *)
(* if reentered). *)
(*-----*)
REVECTORED_INTERRUPT_NUMBER = $FF;

(*-----*)
(* This allows the statistics functions to determine how much of the *)
(* task stack was used by the stack and how much is wasted. *)
(*-----*)
DEFAULT_STACK_VALUE = $A5A5;

(*-----*)
(* 80x86 FLAGS register bit definitions used to test/set/clear *)
(* individual bits in the register. *)
(*-----*)
CARRY_FLAG = $0001; PARITY_FLAG = $0004; AUXILIARY_FLAG = $0010;
ZERO_FLAG = $0040; SIGN_FLAG = $0080; TRAP_FLAG = $0100;
INTERRUPT_FLAG = $0200; DIRECTION_FLAG = $0400; OVERFLOW_FLAG = $0800;

(*-----*)
(* MS-DOS, system and hardware specific interrupt declarations (i.e., *)
(* interrupt vector table entries). This information was derived from *)
(* a number of different sources and represents the most comprehensive *)
(* PC interrupt usage description that could be compiled. *)
(*-----*)
DIVIDE_BY_ZERO_INTERRUPT_NUMBER = $00;
SINGLE_STEP_INTERRUPT_NUMBER = $01;
NMI_INTERRUPT_NUMBER = $02;
BREAKPOINT_INTERRUPT_NUMBER = $03;
OVERFLOW_INTERRUPT_NUMBER = $04;
PRINT_SCREEN_INTERRUPT_NUMBER = $05;
INVALID_OPCODE_INTERRUPT_NUMBER = $06;
SYSTEM_TIMER_INTERRUPT_NUMBER = $08;
DISK_SERVICES_INTERRUPT_NUMBER = $13;
SERIAL_COMM_INTERRUPT_NUMBER = $14;
CASSETTE_IO_INTERRUPT_NUMBER = $15;
KEYBOARD_IO_INTERRUPT_NUMBER = $16;
PRINTER_IO_INTERRUPT_NUMBER = $17;
BOOTSTRAP_INTERRUPT_NUMBER = $1A;
TIME_OF_DAY_INTERRUPT_NUMBER = $1B;
KEYBOARD_BREAK_INTERRUPT_NUMBER = $1C;
CLOCK_INTERRUPT_NUMBER = $1D;
VIDEO_PARAMS_POINTER = $1E;
FLOPPY_DISK_PARAMS_POINTER = $1F;
VIDEO_GRAPHICS_POINTER = $20;
PROGRAM_TERMINATION_INTERRUPT_NUMBER = $21;
MSDOS_BIOS_FUNCTION_INTERRUPT_NUMBER = $22;
PROGRAM_TERMINATION_ADDRESS = $23;
CONTROL_C_ADDRESS = $24;
CRITICAL_ERROR_ADDRESS = $25;
ABSOLUTE_DISK_READ_INTERRUPT_NUMBER = $26;
ABSOLUTE_DISK_WRITE_INTERRUPT_NUMBER = $27;
TSR_INTERRUPT_NUMBER = $28;
IDLE_HANDLER_INTERRUPT_NUMBER = $29;
MSDOS_TTY_HANDLER_INTERRUPT_NUMBER = $2A;
MSDOS_NETWORK_INTERRUPT_NUMBER = $2B;
BATCH_EXECUTE_INTERRUPT_NUMBER = $2C;
MULTIPLY_INTERRUPT_NUMBER = $2D;
CPM_JUMP_PART_1_POINTER = $30;
CPM_JUMP_PART_2_POINTER = $31;
MS_MOUSE_INTERRUPT_NUMBER = $33;
DISKETTE_REQUEST_INTERRUPT_NUMBER = $34;
FIXED_DISK_PARAMS_1_POINTER = $35;
EGA_GRAPHICS_POINTER = $36;
FIXED_DISK_PARAMS_2_POINTER = $37;
USER_ALARM_ADDRESS = $38;
ROM_BIOS_ALARM_INTERRUPT_NUMBER = $39;
NETWORK_FUNCTIONS_INTERRUPT_NUMBER = $3A;
REVECTORED_19H_INTERRUPT_NUMBER = $3B;
NETWORK_USE_INTERRUPT_NUMBER = $3C;
LIM_DMA_INTERRUPT_NUMBER = $3D;
REAL_TIME_CLOCK_INTERRUPT_NUMBER = $3E;
IRQ_2_REDIRECT_INTERRUPT_NUMBER = $3F;
IRQ_0_INTERRUPT_NUMBER = $40;
IRQ_1_INTERRUPT_NUMBER = $41;
IRQ_15_INTERRUPT_NUMBER = $4F;

(*-----*)
(* This set of interrupts are supposedly used by BASIC. There is no *)
(* other documentation to indicate which are actually interrupt *)
(* handlers, in fact examining a 'standard' PC vector table shows that *)
(* most of these are actually nil (i.e., 0000:0000). *)
(*-----*)
BASIC_INTERRUPT_NUMBERS = [$80..$FF];

(*-----*)
(* These interrupts are not used by MS-DOS and documented as usable *)
(* by MS-DOS application programs. *)
(*-----*)
PROGRAM_USABLE_INTERRUPT_NUMBERS = [$60..$66, $E1..$FF];

(*-----*)
(* Members of this set of interrupts are only generated by the 80x86 *)
(* hardware in response to abnormal program operation. *)
(*-----*)
SYSTEM_INTERRUPT_NUMBERS =
[
DIVIDE_BY_ZERO_INTERRUPT_NUMBER,
SINGLE_STEP_INTERRUPT_NUMBER,
NMI_INTERRUPT_NUMBER,
BREAKPOINT_INTERRUPT_NUMBER,
OVERFLOW_INTERRUPT_NUMBER,
INVALID_OPCODE_INTERRUPT_NUMBER,
];

(*-----*)
(* This set of interrupts comprises the Operating Service handlers of *)
(* MS-DOS. This includes 'add-on' services provided by device drivers *)
(* and/or TSR programs. They are all software interrupts which can *)
(* only be accessed by the 80x86 machine instruction 'INT n'. *)
(*-----*)
OS_SERVICES_INTERRUPT_NUMBERS =
[
PRINT_SCREEN_INTERRUPT_NUMBER,
VIDEO_SERVICES_INTERRUPT_NUMBER,
EQUIPMENT_CHECK_INTERRUPT_NUMBER,
MEMORY_SIZE_INTERRUPT_NUMBER,
DISK_SERVICES_INTERRUPT_NUMBER,
SERIAL_COMM_INTERRUPT_NUMBER,
CASSETTE_IO_INTERRUPT_NUMBER,
KEYBOARD_IO_INTERRUPT_NUMBER,
PRINTER_IO_INTERRUPT_NUMBER,
TIME_OF_DAY_INTERRUPT_NUMBER,
KEYBOARD_BREAK_INTERRUPT_NUMBER,
CLOCK_INTERRUPT_NUMBER,
PROGRAM_TERMINATION_INTERRUPT_NUMBER,
MSDOS_BIOS_FUNCTION_INTERRUPT_NUMBER,
ABSOLUTE_DISK_READ_INTERRUPT_NUMBER,
ABSOLUTE_DISK_WRITE_INTERRUPT_NUMBER,
TSR_INTERRUPT_NUMBER,
IDLE_HANDLER_INTERRUPT_NUMBER,
MSDOS_TTY_HANDLER_INTERRUPT_NUMBER,
MSDOS_NETWORK_INTERRUPT_NUMBER,
BATCH_EXECUTE_INTERRUPT_NUMBER,
MULTIPLY_INTERRUPT_NUMBER,
MS_MOUSE_INTERRUPT_NUMBER,
MS_MOUSE_INTERRUPT_NUMBER,
DISKETTE_REQUEST_INTERRUPT_NUMBER,
ROM_BIOS_ALARM_INTERRUPT_NUMBER,
];
```

```

NETWORK_FUNCTIONS_INTERRUPT_NUMBER,
REVECTORED_19H_INTERRUPT_NUMBER,
NETWORK_USE_INTERRUPT_NUMBER,
LIM_EMS_INTERRUPT_NUMBER
);

(* The members of this set of interrupts are only documented as *)
(* 'RESERVED'. The use of any of these interrupts is therefore *)
(* forbidden (in order to maintain 100% compatibility). Unfortunately *)
(* some programs use interrupts from this set. An example of this *)
(* is DEC PathWorks Network driver which uses $68. Note that this *)
(* unauthorized use can cause incompatibilities with TASKING (and *)
(* other programs). *)
RESERVED_INTERRUPT_NUMBERS =
(
    $07, $2B..$2D, $32, $34..$3F, $42, $44..$45,
    $47..$49, $4B..$5B, $5D..$5F, $68..$6E, $78..$7F
);

(* The members of this set of interrupts are used in an incompatible *)
(* manner. These incompatibilities were discovered by testing, there *)
(* are no documents which describe in any way in which the interrupts are *)
(* used. These interrupts are supposed to be used for other purposes *)
(* but are 'taken over' by the applications listed below and used in a *)
(* manner that is not standard and is not completely understood. *)
INCOMPATIBLE_INTERRUPT_NUMBERS =
(
    (*-----*)
    (* DEC PathWorks Network Driver Software *)
    (* It appears that as part of handling this interrupt the driver *)
    (* modifies the location of the next interrupt goes to a *)
    (* different location in the driver. This is definitely not a good *)
    (* approach and TASKING cannot deal with it. Because TASKING has *)
    (* taken over the interrupt and the passes control to the actual *)
    (* handler, if the handler then wants the next interrupt to go *)
    (* somewhere else TASKING has no way of intercepting it. The only *)
    (* alternative is to let the driver 'fend for itself' and hope that *)
    (* it can handle reentrancy or that reentrancy never occurs. *)
    $68
    (*-----*)
);

(* The 18259A Programmable Interrupt Controller (PIC) chip must be *)
(* directly accessed by TASKING in order to support preemptive multi- *)
(* tasking. These declarations provide the necessary PC hardware *)
(* specific information. *)
NON_SPECIFIC_END_OF_INTERRUPT = $20;
PROGRAMMABLE_INTERRUPT_CONTROLLER_PORT = $20;

(* Areas of MS-DOS Data area used to support blocked keyboard input. *)
DOS_DATA_SEGMENT = $0040;
KYBD_BUFFER_HEAD_PTR_OFS = $001A;
KYBD_BUFFER_TAIL_PTR_OFS = $001C;
KYBD_BUFFER_BEGIN_PTR_OFS = $0080;
KYBD_BUFFER_END_PTR_OFS = $0082;

(* General text processing control characters. *)
CR = chr($0D); (* Carriage Return *)
LF = chr($0A); (* Line Feed *)
NULL_CHR = chr($00); (* Marks the beginning of an extended key *)

type
(* In order to allow efficient allocation/deallocation of TASK IDs, *)
(* TASKING maintains a list of sets of TASK_IDS. Since each set can *)
(* hold 256 elements, this list defines the number (and range) of such *)
(* sets. Note that '0' must be the start of the range. *)
TASK_ID_SETS = 0..(MAXIMUM_NUMBER_OF_TASKS div 256 + 1);

(* When tasks use message passing then this data structure is used to *)
(* manage the message lists within each mailbox. The size of the *)
(* mailbox is dependent upon system heap space only because the *)
(* messages are managed by this linked list structure. *)
MESSAGE_PTR = ^MESSAGE;
MESSAGE = record
    INFO_ADDRESS : pointer;
    NEXT : ^MESSAGE_PTR;
end; (* MESSAGE *)

(* When tasks use message passing then this data structure is used to *)
(* maintain the list of mailboxes. Each mailbox must have an owner, *)
(* i.e. the task that will receive the mail and each task can have at *)
(* most one mailbox. Messages are managed using semaphores. The *)
(* number of tasks receiving mail is dependent upon heap space only *)
(* because the mailboxes are managed by this linked list structure. *)
(* For performance improvements this could be a binary tree sorted by *)
(* the OWNER field. Performance will only become an issue with very *)
(* large applications with many message passing tasks. Note that *)
(* performance will only degrade if the number of tasks that receive *)
(* messages is large NOT if the number of messages being passed is *)
(* large. *)
MAILBOX_PTR = ^MAILBOX;
MAILBOX = record
    OWNER : TASK_IDS;
    SEM : SEMAPHORE;
    CONTENTS : MESSAGE_PTR;
    NEXT : MAILBOX_PTR;
end; (* MAILBOX *)

(* The task context is saved into and restored from this data *)
(* structure. It is basically a representation of the 8086 registers. *)
(* Remember all applications under DOS actually run on an 8086 no *)
(* matter what CPU is used in the PC. *)
CONTEXT_REGISTERS = record
    AX : word; BX : word; CX : word; DX : word;
    BP : word; SI : word; DI : word;
    ES : word; DS : word;
    SS : word; SP : word;
    CS : word; IP : word;
    FLAGS : word;
end; (* CONTEXT_REGISTERS *)

(* Basic pointer to a semaphore, used to find the semaphore that the *)
(* tasks are waiting on. Internally semaphores are distinguished by *)
(* their addresses. *)
SEMAPHORE_PTR = ^SEMAPHORE;

(* This data structure is used to maintain the information about what *)
(* (and if) a task is waiting for. Since everything except time is *)
(* implemented as a semaphore a task waiting for a mail message (for

```

```
(* used by the application (used when priority inheritance is enabled).*)
LOCKED_RESOURCE_LIST_PTR : LOCKED_RESOURCE_PTR = nil;

(* These variables are used when TASKING statistics are gathered to *)
(* provide the user with various system parameters which can be of *)
(* interest to the TASKING application developer. *)
TASKING_STATISTICS : record
START_TIME : real;
NUMBER_OF_PERIODIC_EVENTS : longint;
NUMBER_OF_PERIODIC_EVENTS_MISSED : longint;
CONTEXT_SWITCHES : longint;
COOPERATIVE_CONTEXT_SWITCHES : longint;
TOTAL_AVAILABLE_HEAP : longint;
SERVICE_CALLS : array($00..$FF) of longint;
HARDWARE_INTERRUPTS : array($00..$0F) of longint;
end (* TASKING_STATISTICS *) =

START TIME : 0.0;
NUMBER OF PERIODIC EVENTS : 0;
NUMBER OF PERIODIC EVENTS MISSED : 0;
CONTEXT SWITCHES : 0;
COOPERATIVE CONTEXT SWITCHES : 0;
TOTAL AVAILABLE HEAP : 0;
SERVICE CALLS : array($00..$0F) of 0;
HARDWARE INTERRUPTS : array($00..$0F) of 0;

(* This variable keeps track of the nesting level of the scheduler *)
(* 'locks' that have occurred, its value can not be negative. *)
IN_PREEMPTABLE_REGION : integer = 0;

(* Default Error Handlers (Preview) *)
(* $F* Handlers must be 'fair' because they are used as procedural variables *)
var
(* This file is created (in the same directory as the application *)
(* executable if LOG_ERRORS_TO_FILE is true) and contains a log of all *)
(* errors and warnings detected by TASKING. *)
ERROR_LOG : text;

procedure TASK_ALREADY_ACTIVE_ERROR_HANDLER (TASK_ID : TASK_IDS); forward;
procedure INSUFFICIENT_RESOURCES_ERROR_HANDLER (TASK_ID : TASK_IDS); forward;
procedure TASK_ALREADY_SUSPENDED_ERROR_HANDLER (TASK_ID : TASK_IDS); forward;
procedure ILLEGAL_TASK_ID_ERROR_HANDLER (TASK_ID : TASK_IDS); forward;
procedure ILLEGAL_OPERATION_ERROR_HANDLER (TASK_ID : TASK_IDS); forward;
procedure DEADLOCK_DETECTED_ERROR_HANDLER; forward;

(* End of Default Error Handlers (Preview) *)
(* $F- *)

(* Null Task Definition *)
procedure QUEUE_AND_RESCHEDULE (TASK : TASK_IDS; RESULTING_STATE TASK_STATES); forward;

procedure NULL_TASK;
(* This task is used only to provide a place for the CPU to execute when *)
(* it has absolutely nothing else to do. All that this task tries to do is *)
(* pass control to tasks which are ready to run by continuously rescheduling *)
(* itself, since the null task has the lowest possible priority, all tasks *)
(* which are READY to run will be scheduled as soon as possible. *)
(* When it is not possible to pass control to another task (i.e., there *)
(* are no other tasks in the application is terminated. *)
begin (* NULL_TASK *)
repeat
asm shl end;
(* If there are no other user tasks which are still active or *)
(* blocked, then the only remaining task is the null task. If the *)
(* null task is the only one left for terminate everything, else *)
(* pass control to some other task (or at least attempt to). *)
if (READY_QUEUE.HEAD = nil) and (READY_QUEUE.TAIL = nil) and
(Delayed_Queue_Head = nil) and (Delayed_Queue_Tail = nil)
then
if (Blocked_Queue_Head = nil) and (Blocked_Queue_Tail = nil)
then
(* Shut down everything (via the exit procedure chain) *)
halt(0)
else
(* Deadlock is defined as the condition where all tasks that *)
(* exist are blocked waiting for something to be signaled by *)
(* another task. Since all tasks are blocked (and will not be *)
(* scheduled to run) the action necessary to unblock any of *)
(* the tasks cannot be satisfied. In TASKING, this deadlock *)
(* condition exists when there are no READY or DELAYED tasks *)
(* and all existing tasks are BLOCKED on things 'other than' *)
(* keyboard or mouse events (which can be signaled as a result *)
(* of interrupts). *)
if (USER_KEYBOARD_SEM_PTR = nil) and
(USER_MOUSE_EVNT_PTR = nil)
then
DEADLOCK_DETECTED_ERROR_HANDLER
else
(* for the dangling 'else' *)
else
(* Continuously attempt to surrender the CPU to another task. *)
QUEUE_AND_RESCHEDULE(NULL_TASK_ID, READY);
until false;
end; (* NULL_TASK *)

const NULL_TASK_STACK_SIZE = 511;
var NULL_TASK_STACK : array[0..NULL_TASK_STACK_SIZE] of word;
const NULL : TASK_CONTROL_BLOCK =
( TASK_ID : NULL_TASK_ID;
clock_tick : 0;
PRIORITY : LOWEST_PRIORITY;
STACK_PTR : addr(NULL_TASK_STACK);
STACK_SIZE : sizeof(NULL_TASK_STACK);
CONTEXT :
( AX : $0000; BX : $0000; CX : $0000; DX : $0000;
BP : $0000; SI : $0000; DI : $0000;
ES : $0000; DS : seg(NULL_TASK_STACK);
SS : seg(NULL_TASK_STACK[NUL_TASK_STACK_SIZE]);
SP : ofs(NULL_TASK_STACK[NUL_TASK_STACK_SIZE]);
CS : seg(NULL_TASK); IP : ofs(NULL_TASK);
FLAGS : INTERRUPT_FLAG
);
WAITING_FOR : nil;
ERROR_HANDLER :
( TASK_ALREADY_ACTIVE_ERROR_HANDLER,
INSUFFICIENT_RESOURCES_ERROR_HANDLER,
TASK_IS_NOT_ACTIVE_ERROR_HANDLER,
TASK_ALREADY_SUSPENDED_ERROR_HANDLER,
ILLEGAL_TASK_ID_ERROR_HANDLER,
ILLEGAL_OPERATION_ERROR_HANDLER
);
NEXT : nil;
PREVIOUS : nil;
);
(* End of Null Task Definition *)
const
(* This variable always points to the currently running task *)
(* control block. There can only be one task which is running at any *)
(* point in time. This variable will never be nil. *)
RUNNING_TASK_PTR : TASK_CONTROL_BLOCK_PTR = @NULL;

(* This variable contains a pointer to the beginning of the list of *)
(* mailboxes, of course until a message is sent or a receive is *)
(* attempted there are no mailboxes. *)
MAILBOX_LIST : MAILBOX_PTR = nil;

(* This variable is used to hold the number of milliseconds which *)
(* elapsed since the program started running. It is the basis for all *)
(* task delay calculations. *)
MILLISECOND_TICK : longint = 0;

(* This variable is used to prevent timer interrupts from occurring *)
(* until all processing for the present interrupt has completed. This *)
(* ensures that TASKING internal data structures are not compromised *)
(* due to re-entrant time (this can only occur when *)
(* interrupt handling takes longer than the clock timer period). *)
TIMER_NSOI_NEEDS_TO_BE_DONE : boolean = false;

(* This variable is used to allow TASKING to know when it is 'fair' *)
(* to preempt a running task. This ensures that all tasks get their *)
(* full timeslice every time (of course this is only used when the *)
(* tasking mode is preemptive). *)
PREEMPTIVE_TIMESLICE : real = 0.0; (* Macroseconds *)

(* This is used to account for the possible inaccuracy between the *)
(* target timeslice and the time period that the clock is capable of *)
(* providing (due to integer count values). *)
95_PERCENT = 0.95;

(* This counter is used to prevent the TASKING application from *)
(* terminating if there are child programs still executing. Failure *)
(* to do this will almost certainly crash MS-DOS because a child will *)
(* live past its parent, MS-DOS never expects this and internal MS-DOS *)
(* data structures and memory allocation schemes do not support it. *)
NUMBER_OF_CHILD_PROGRAMS : integer = 0;

(* This flag is set once the application 'main' has terminated. This *)
(* ensures that a task created in 'main' does not cause a reschedule *)
(* (rescheduling 'main' is meaningless because it is not a task). *)
ALLOW_RESCHEDULE_IN_CREATE : boolean = false;

(* When TASKING is started there are no mouse events to report (hence *)
(* the initialization values). This variable is used to transfer *)
(* MS-Mouse parameters from the TASKING mouse handler to the user task *)
(* (if and only if a mouse task has been installed). *)
MOUSE_INFO_FROM_HANDLER : MS_MOUSE.MOUSE_PARAMETERS =
( ACTIVITY_MASK : 0;
BUTTON_STATES : 0;
VERTICAL_TEXT_POSITION : 0;
HORIZONTAL_TEXT_POSITION : 0;
);

(* This set of interrupts are 'all' revector to the same interrupt *)
(* handler which determines the actual interrupt number, disables *)
(* preemption and then dispatches to the actual interrupt handler. *)
(* Only 'software' interrupts can be included in this set. *)
PROTECTED_INTERRUPTS :
OS_SERVICES_INTERRUPT_NUMBERS (* System Software Interrupts *)
+ PROGRAM_USABLE_INTERRUPT_NUMBERS (* User " " *)
- SYSTEM_INTERRUPT_NUMBERS (* System Hardware " *)
+ RESERVED_INTERRUPT_NUMBERS (* Unused " *)
+ BASIC_INTERRUPT_NUMBERS (* BASIC " *)
- INCOMPATIBLE_INTERRUPT_NUMBERS (* Non-standard usage " *)
- (REVECTORED_INTERRUPT_NUMBER); (* TASKING Generic Interrupt *)

var
(* When TASKING is started there are no TASK_IDs allocated so the set *)
(* of allocated task identifiers must be the empty set. The task *)
(* identifier of the null task is not included in this set because it *)
(* cannot be allocated to user tasks. *)
ALLOCATED_TASK_IDS : array[TASK_ID_SETS] of set of byte;

(* Loop control variable to initialize allocated TASK_ID sets. *)
TASK_ID_SET : TASK_ID_SETS;

(* This variable contains the address of the next exit procedure in *)
(* the exit call chain. This value is restored when TASKING *)
(* terminates. *)
SAVE_EXIT : pointer;

```

```

(* Pointer to original clock interrupt handler, used to allow MS-DOS *)
(* to see clock ticks at 18.2 Hz rate (initialized at startup). *)
CLOCK_VECTOR : pointer;

(* This variable contains the original address of the interrupt that *)
(* TASKING 'takes over', it is saved here during initialization so *)
(* that it can be restored upon termination. *)
OLD_REVECTORED_INTERRUPT : HANDLERS;

(* This variable contains information used to manage interrupts which *)
(* originate from hardware sources. This is necessary because add-on *)
(* card driver software may (or may not) support preemption, this *)
(* guarantees that it won't be preempted in case reentrancy will cause *)
(* the driver to crash. *)
HARDWARE_HANDLER : array[D..15] of HANDLERS;

(* This variable contains the addresses of the next MS-DOS function *)
(* dispatcher in the call chains. These are restored when TASKING *)
(* terminates. *)
INTERRUPT_VECTOR_HANDLER : array[$DD..$FF] of HANDLERS;

(* These variables define TASKING's view of clock timer ticks which *)
(* TASKING calculates based on the requested timeslice value. (Don't *)
(* worry...before termination it is restored back to what MS-DOS *)
(* expects it to be). *)
TIMER_PERIOD : real; (* Microseconds *)
NEW_TIMER_D_VALUE : word;

(* The TASKING configuration is copied from the default (user) values *)
(* just prior to beginning multi-tasking, this variable holds the *)
(* configuration parameters used by TASKING. *)
ACTUAL_TASKING_CONFIGURATION : CONFIGURATION;

(* Non-User functions *)

function COMPARE_PRIORITIES
(* This routine compares the relative priorities passed in based on the *)
(* highest dynamic priority 'at this time'. This is necessary so that a *)
(* priority rotation scheme can be used. Note that for static priorities *)
(* this routine will also return meaningful results (although direct *)
(* comparisons are possible). *)
var
(* Converted priorities which can then be directly compared for *)
(* relative value. *)
NORMALIZED_A, NORMALIZED_B : integer;
begin (* COMPARE_PRIORITIES *)
  case PRIORITY_A of
    LOWEST_PRIORITY : NORMALIZED_A := -maxint;
    HIGHEST_PRIORITY : NORMALIZED_A := +maxint;
  else
    if PRIORITY_A > HIGHEST_DYNAMIC_PRIORITY
    then
      NORMALIZED_A := integer(PRIORITY_A) mod
        integer(HIGHEST_DYNAMIC_PRIORITY)
    else if PRIORITY_A = HIGHEST_DYNAMIC_PRIORITY then
      NORMALIZED_A := integer(HIGHEST_DYNAMIC_PRIORITY)
    else (* PRIORITY_A < HIGHEST_DYNAMIC_PRIORITY *)
      NORMALIZED_A := PRIORITY_A * (HIGHEST_DYNAMIC_PRIORITY - 1 -
        integer(HIGHEST_DYNAMIC_PRIORITY));
  end; (* case...of *)
  case PRIORITY_B of
    LOWEST_PRIORITY : NORMALIZED_B := -maxint;
    HIGHEST_PRIORITY : NORMALIZED_B := +maxint;
  else
    if PRIORITY_A > HIGHEST_DYNAMIC_PRIORITY
    then
      NORMALIZED_B := integer(PRIORITY_B) mod
        integer(HIGHEST_DYNAMIC_PRIORITY)
    else if PRIORITY_B = HIGHEST_DYNAMIC_PRIORITY then
      NORMALIZED_B := integer(HIGHEST_DYNAMIC_PRIORITY)
    else (* PRIORITY_B < HIGHEST_DYNAMIC_PRIORITY *)
      NORMALIZED_B := PRIORITY_B * (HIGHEST_DYNAMIC_PRIORITY - 1 -
        integer(HIGHEST_DYNAMIC_PRIORITY));
  end; (* case...of *)
  case COMPARISON of
    GREATER_THAN : COMPARE_PRIORITIES := NORMALIZED_A > NORMALIZED_B;
    LESS_THAN : COMPARE_PRIORITIES := NORMALIZED_A < NORMALIZED_B;
    EQUAL_TO : COMPARE_PRIORITIES := NORMALIZED_A = NORMALIZED_B;
  end; (* case...of *)
end; (* COMPARE_PRIORITIES *)

procedure INSERT_INTRO
(* This routine inserts the specified task control block into the queue *)
(* specified using the task priority to insert the task into the correct *)
(* place in the priority queue. *)
var
(* Task Control Block pointer which is 'walked' through the queue *)
(* in search of the proper insertion point. *)
WALKING_PTR : TASK_CONTROL_BLOCK_PTR;
(* Temporary pointer needed to perform insertion between two task *)
(* control blocks. *)
TEMP_PTR : TASK_CONTROL_BLOCK_PTR;
begin (* INSERT_INTRO *)
  asm
    pushf
    cli
  end; (* asm *)
  (* Find place in queue to insert into *)
  while (WALKING_PTR := QUEUE.TAIL;
    while (WALKING_PTR <> nil) and
      COMPARE_PRIORITIES(NEW_TCB_PTR^.PRIORITY, GREATER_THAN,
        WALKING_PTR^.PRIORITY) do
    WALKING_PTR := WALKING_PTR^.PREVIOUS;
  end; (* Insert element *)
  if WALKING_PTR = nil
  then (* This task is highest priority *)
    if QUEUE.HEAD = nil
    then (* the list is empty *)
      begin
        QUEUE.HEAD := NEW_TCB_PTR;
        QUEUE.HEAD^.PREVIOUS := nil;
        QUEUE.TAIL := NEW_TCB_PTR;
        QUEUE.TAIL^.NEXT := nil;
      end (* if...then *)
    else
      begin
        NEW_TCB_PTR^.NEXT := QUEUE.HEAD;
        QUEUE.HEAD^.PREVIOUS := NEW_TCB_PTR;
        QUEUE.HEAD := NEW_TCB_PTR;
        QUEUE.HEAD^.PREVIOUS := nil;
      end (* if...then...else *)
    else
      if WALKING_PTR = QUEUE.TAIL
      then (* This task is lowest priority *)
        begin
          QUEUE.TAIL^.NEXT := NEW_TCB_PTR;
          NEW_TCB_PTR^.PREVIOUS := QUEUE.TAIL;
          QUEUE.TAIL := NEW_TCB_PTR;
          QUEUE.TAIL^.NEXT := nil;
        end (* if...then *)
      else (* This task is somewhere in the list *)
        begin
          TEMP_PTR := WALKING_PTR^.NEXT;
          WALKING_PTR^.NEXT := NEW_TCB_PTR;
          NEW_TCB_PTR^.PREVIOUS := WALKING_PTR;
          NEW_TCB_PTR^.NEXT := TEMP_PTR;
          TEMP_PTR^.PREVIOUS := NEW_TCB_PTR;
          end; (* if...then...else *)
        asm popf end;
      end; (* INSERT_INTRO *)

function REMOVE_FROM_HEAD
(* This routine removes one task control block from the head of the *)
(* specified queue and returns that pointer. Note that this works even if *)
(* the queue is empty because then nil is returned. *)
var
(* Temporary pointer necessary for knowing the value of the 'head' *)
(* before the removal was performed. *)
TEMP_PTR : TASK_CONTROL_BLOCK_PTR;
begin (* REMOVE_FROM_HEAD *)
  asm
    pushf
    cli
  end; (* asm *)
  TEMP_PTR := QUEUE.HEAD;
  if QUEUE.HEAD <> QUEUE.TAIL
  then
    begin
      QUEUE.HEAD := QUEUE.HEAD^.NEXT;
      QUEUE.HEAD^.PREVIOUS := nil;
    end (* if...then *)
  else
    begin
      QUEUE.HEAD := nil;
      QUEUE.TAIL := nil;
    end; (* if...then *)
  if TEMP_PTR <> nil then
    begin
      TEMP_PTR^.NEXT := nil;
      TEMP_PTR^.PREVIOUS := nil;
    end; (* if...then *)
  REMOVE_FROM_HEAD := TEMP_PTR;
  asm popf end;
end; (* REMOVE_FROM_HEAD *)

function FIND_AND_REMOVE
(* This routine searches the specified queue for the task control block *)
(* which matches the search criteria implied by the reason that the task is *)
(* being removed. If a matching task control block is found then it is *)
(* returned, nil is returned if no match is found in the specified queue. *)
var
(* Task Control Block pointer which is 'walked' through the queue *)
(* in search of the block which is to be removed. *)
WALKING_PTR : TASK_CONTROL_BLOCK_PTR;
begin (* FIND_AND_REMOVE *)
  asm
    pushf
    cli
  end; (* asm *)
  WALKING_PTR := QUEUE.HEAD;
  case REASON of
    UNCONDITIONALLY :
      while (WALKING_PTR <> nil) and (WALKING_PTR^.TASK_ID <> TASK) do
        WALKING_PTR := WALKING_PTR^.NEXT;
      BECAUSE THE BLOCK WAS SIGNALLED :
        while (WALKING_PTR <> nil) and
          ((WALKING_PTR^.WAITING_FOR = nil) or
            ((WALKING_PTR^.WAITING_FOR <> nil) and
              (WALKING_PTR^.WAITING_FOR^.SEM_PTR <> SEM_PTR)))
          do
            WALKING_PTR := WALKING_PTR^.NEXT;
          BECAUSE THE DELAY HAS COMPLETED :
            while (WALKING_PTR <> nil) and
              ((WALKING_PTR^.WAITING_FOR <> nil) and
                ((WALKING_PTR^.WAITING_FOR^.ABSOLUTE_TIME = NO_TIME_DELAY) or
                  (WALKING_PTR^.WAITING_FOR^.ABSOLUTE_TIME > MILLISECOND_TICKS)))
              do
                WALKING_PTR := WALKING_PTR^.NEXT;
            end; (* case...of *)
  if (REASON <> BECAUSE THE BLOCK WAS SIGNALLED) and
    (WALKING_PTR <> nil)
  ) or
    (REASON = BECAUSE THE BLOCK WAS SIGNALLED) and
    (WALKING_PTR^.WAITING_FOR <> nil) and
    (WALKING_PTR^.WAITING_FOR^.SEM_PTR = SEM_PTR)
  then
    begin
      if (WALKING_PTR^.PREVIOUS <> nil) and (WALKING_PTR^.NEXT <> nil)
      then
        begin
          WALKING_PTR^.PREVIOUS^.NEXT := WALKING_PTR^.NEXT;

```

```

WALKING_PTR^.NEXT^.PREVIOUS := WALKING_PTR^.PREVIOUS;
end if ..then *
else if !WALKING_PTR^.PREVIOUS := nil; and
    !WALKING_PTR^.NEXT <> nil;
then
begin
    QUEUE.HEAD := QUEUE.HEAD^.NEXT;
    QUEUE.HEAD^.PREVIOUS := nil;
end if ..then *
else if !WALKING_PTR^.PREVIOUS <> nil; and
    !WALKING_PTR^.NEXT := nil;
then
begin
    QUEUE.TAIL := QUEUE.TAIL^.PREVIOUS;
    QUEUE.TAIL^.NEXT := nil;
end if ..then *
else
begin
    QUEUE.HEAD := nil;
    QUEUE.TAIL := nil;
end if ..then ..else *
WALKING_PTR^.NEXT := nil;
WALKING_PTR^.PREVIOUS := nil;
if REASON <> UNCONDITIONALLY then
begin
    WALKING_PTR^.WAITING_FOR^.ABSOLUTE_TIME := NO_TIME_DELAY;
    WALKING_PTR^.WAITING_FOR^.SEM_PTR := nil;
end if ..then *
FIND_AND_REMOVE := WALKING_PTR;
end if ..then *
else
    FIND_AND_REMOVE := nil;
asm popf end;
end; !* FIND_AND_REMOVE *

function FIND_OR_CREATE_MBOX_FOR
.....
[TASK : TASK_IDS] : MAILBOX_PTR;
.....
!* This routine searches the system mailbox list and returns a pointer to *
!* the mailbox owned by the specified task. If the task does not yet own a *
!* mailbox then one is created and added to the mailbox list. *
.....
var
.....
!* Pointer to mailbox structures which is "walked" through the list *
!* in search of the matching owner. *
.....
WALKING_PTR : MAILBOX_PTR;
.....
!* Temporary return value [function name cannot be used on right *
!* hand side of an assignment statement]. *
.....
TEMP_PTR : MAILBOX_PTR;
.....
begin !* FIND_OR_CREATE_MBOX_FOR *!
asm
    pushf
    cli
end; !* asm *!
WALKING_PTR := MAILBOX_LIST;
while !WALKING_PTR <> nil; and !WALKING_PTR^.OWNER <> TASK; do
    WALKING_PTR := WALKING_PTR^.NEXT;
if WALKING_PTR^.OWNER = TASK
then
    TEMP_PTR := WALKING_PTR
else
begin
    if WALKING_PTR = nil
    then
        new[WALKING_PTR];
    else
begin
        new[WALKING_PTR^.NEXT];
        WALKING_PTR := WALKING_PTR^.NEXT;
        end; !* if ..then ..else *!
        TEMP_PTR := MAILBOX_LIST;
        MAILBOX_LIST := WALKING_PTR;
        MAILBOX_LIST^.NEXT := TEMP_PTR;
        TEMP_PTR := MAILBOX_LIST;
        with TEMP_PTR do
            begin
                OWNER := TASK;
                SEM := 0;
                CONTENTS := nil;
                end; !* with ..do *!
            end; !* if ..then ..else *!
        FIND_OR_CREATE_MBOX_FOR := TEMP_PTR;
    asm popf end;
end; !* FIND_OR_CREATE_MBOX_FOR *!

procedure REMOVE_MBOX_OF
.....
[TASK : TASK_IDS];
.....
!* This routine searches the system mailbox list and removes the mailbox *
!* owned by the specified task. If the task does not yet own a mailbox then *
!* nothing is done. *
.....
var
.....
!* Pointer to mailbox structures which is "walked" through the list *
!* in search of the matching owner. *
.....
WALKING_PTR : MAILBOX_PTR;
.....
!* Temporary pointer used to re-arrange mailbox list links. *
.....
TEMP_PTR : MAILBOX_PTR;
.....
begin !* REMOVE_MBOX_OF *!
asm
    pushf
    cli
end; !* asm *!
WALKING_PTR := MAILBOX_LIST;
while !WALKING_PTR <> nil; and !WALKING_PTR^.OWNER <> TASK; do
    WALKING_PTR := WALKING_PTR^.NEXT;
if WALKING_PTR^.OWNER = TASK then
begin
    TEMP_PTR := WALKING_PTR;
    !* Patch link past element to delete *!
    WALKING_PTR := MAILBOX_LIST;
    while WALKING_PTR <> TEMP_PTR do
        WALKING_PTR := WALKING_PTR^.NEXT;
        WALKING_PTR^.NEXT := TEMP_PTR^.NEXT;
    !* Free memory allocated to deleted mailbox *!
    dispose[TEMP_PTR];
    end; !* if ..then *!
    asm popf end;
end; !* REMOVE_MBOX_OF *!

procedure CLOCK_INTERCEPTOR; interrupt;
.....
!* This routine checks to see if the DELAYED queue is empty and if it is *
.....
const
.....
!* This counter is used to determine when the MS-DOS clock *
!* interrupt handler should be called so that MS-DOS thinks that the *
!* timer is running at a 55 msec period. *
.....
MSDOS_CLOCK_PERIOD : real = DEFAULT_MSDOS_CLOCK_PERIOD [sec] * 1e6 / 2;
.....
!* This counter is used to generate TASKING millisecond clock ticks *
.....
MILLISECOND_PERIOD : real = 0.0; [usec];
.....
var
.....
!* Task Control Block pointer used to release all tasks that have *
!* been delayed for the proper amount of time. *
.....
DELAY_COMPLETED_TASK_PTR : TASK_CONTROL_BLOCK_PTR;
.....
!* Pointer to the list of periodic events [possibly nil]. *
.....
WALKING_PTR : PERIODIC_EVENTS_PTR;
.....
!* Pointer to Task Control Block of waiting task to be made ready. *
.....
WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;
.....
begin !* CLOCK_INTERCEPTOR *!
    LOCK_SCHEDULER;
    if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
        begin
            inc[TASKING_STATISTICS.HARDWARE_INTERRUPTS[0]];
            inc[RUNNING_TASK_PTR^.CLOCK_TICKS];
            end; !* if ..then *!
    !* Perform all local processing *!
    MILLISECOND_PERIOD := MILLISECOND_PERIOD + TIMER_PERIOD;
    if MILLISECOND_PERIOD >= !1000 [usec] * 95_PERCENT; then
        begin
            MILLISECOND_PERIOD := MILLISECOND_PERIOD - !1000.0 [usec];
            inc[MILLISECOND_TICKS];
            end; !* if ..then *!
    if DELAYED_QUEUE.HEAD <> nil then
        begin
            DELAY_COMPLETED_TASK_PTR := FIND_AND_REMOVE
            !
            ANY_TASK,
            !From the; DELAYED QUEUE,
            BECAUSE THE DELAY HAS COMPLETED,
            THERE IS NO_BLOCK_TO MATCH
            !
            !
            if DELAY_COMPLETED_TASK_PTR <> nil then
                INSERT_INTO_READY_QUEUE, DELAY_COMPLETED_TASK_PTR;
            end; !* if ..then *!
    !* Find the place in the list to insert the event *!
    WALKING_PTR := PERIODIC_EVENTS_LIST_PTR;
    while WALKING_PTR <> nil do
        begin
            if !MILLISECOND_TICKS mod WALKING_PTR^.PERIOD; = 0 then
                with WALKING_PTR do
                    begin
                        if EVENT_PTR = SIGNALLED then
                            inc[TASKING_STATISTICS.NUMBER_OF_PERIODIC_EVENTS_MISSED];
                            WAITING_TASK_PTR := FIND_AND_REMOVE
                            !
                            ANY_TASK,
                            !From the; BLOCKED QUEUE,
                            BECAUSE THE BLOCK WAS SIGNALLED,
                            !At; addr;EVENT_PTR!
                            !
                            if WAITING_TASK_PTR = nil
                                then
                                    EVENT_PTR := SIGNALLED
                                else
                                    INSERT_INTO_READY_QUEUE, WAITING_TASK_PTR;
                                    inc[TASKING_STATISTICS.NUMBER_OF_PERIODIC_EVENTS];
                                    end; !* with ..do *!
                                    WALKING_PTR := WALKING_PTR^.NEXT;
                                    end; !* while ..do *!
        !* Call original interrupt handler once every 55 msec! *!
        MSDOS_CLOCK_PERIOD := MSDOS_CLOCK_PERIOD + TIMER_PERIOD;
        if MSDOS_CLOCK_PERIOD >= !DEFAULT_MSDOS_CLOCK_PERIOD * 95_PERCENT;
        then
            begin
                MSDOS_CLOCK_PERIOD := MSDOS_CLOCK_PERIOD -
                    DEFAULT_MSDOS_CLOCK_PERIOD * 1e6;
            asm
                pushf
                call dword ptr ds:CLOCK_VECTOR
                cli
            end; !* asm *!
            end; !* if ..then *!
        else
            TIMER_NSEOI_NEEDS_TO_BE_DONE := true;
            UNLOCK_SCHEDULER;
        if !ACTUAL_TASKING_CONFIGURATION.TASKING_MODEL = PREEMPTIVE; and
            !IN_PREEMPTABLE_REGION = 0; then
            begin
                PREEMPTIVE_TIMESLICE := PREEMPTIVE_TIMESLICE + TIMER_PERIOD;
                if PREEMPTIVE_TIMESLICE >=
                    ACTUAL_TASKING_CONFIGURATION.TARGET_TIMESLICE
                then
                    REQUEST_QUEUE_AND_RESCHEDULE[RUNNING_TASK_PTR^.TASK_ID, READY];
                    end; !* if ..then *!
            if TIMER_NSEOI_NEEDS_TO_BE_DONE then
                begin
                    !* Issue NSEOI to PIC *!
                    port[PROGRAMMABLE_INTERRUPT_CONTROLLER_PORT] :=
                        NON_SPECIFIC_END_OF_INTERRUPT;
                    TIMER_NSEOI_NEEDS_TO_BE_DONE := false;
                    end; !* if ..then *!
                end; !* CLOCK_INTERCEPTOR *!
        procedure KEYBOARD_INTERCEPTOR; interrupt;
        .....
        !* This routine checks to see if the keyboard interrupt that was just *
        !* detected corresponds to an incoming character or just some type of *
        !* keyboard activity [like pressing the SHIFT, ALT, CTRL or 'lock' key]. If *
        !* the interrupt was because of a key pressed then the controlling semaphore *
        !* is signaled if and only if the semaphore was created as a result of a *
        !* task waiting for a key to be pressed. *
        .....
        const
        .....
        !* Pointer to pointer to the head of the keyboard buffer. *
        .....
        HEAD_PTR = ptr[IDOS_DATA_SEGMENT, KYBD_BUFFER_HEAD_PTR_OFFSET];

```

```

(* Pointer to pointer to the tail of the keyboard buffer. *)
TAIL_PTR = ptr(DOS_DATA_SEGMENT, KYBD_BUFFER_TAIL_PTR_OFS);

(* Pointer to the beginning of the keyboard buffer (not the head. *)
BUFFER_START = ptr(DOS_DATA_SEGMENT, KYBD_BUFFER_BEGIN_PTR_OFS);

(* Pointer to the end of the keyboard buffer (not the tail. *)
BUFFER_END = ptr(DOS_DATA_SEGMENT, KYBD_BUFFER_END_PTR_OFS);

var
(* Pointer to Task Control Block which is waiting on the keyboard *)
(* semaphore (may very well be nil). *)
WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

(* Offset into the buffer of the tail "before" the interrupt is *)
(* processed by the MS-DOS handler (used to determine if a character *)
(* was received). *)
TAIL : word;

(* The number of characters received into the buffer, probably one *)
(* for each interrupt but could be more. *)
NUM_CHARS : integer;

(* Pointer to original keyboard handler. *)
VECTOR : pointer;

begin (* KEYBOARD INTERCEPTOR *)
VECTOR := HARDWARE_HANDLER($1).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($1));

TAIL := word(TAIL_PTR^);

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

if (TAIL < word(TAIL_PTR^)) and (USER_KEYBOARD_SEM_PTR <> nil) then
begin (* tail changed, therefore a character was received *)
WAITING_TASK_PTR := FIND_AND_REMOVE
(
ANY_TASK,
(From the) BLOCKED_QUEUE,
BECAUSE THE BLOCK WAS SIGNALLED,
(AT) addr(USER_KEYBOARD_SEM_PTR^
);
if TAIL < word(TAIL_PTR^)
then
NUM_CHARS := word(TAIL_PTR^) - TAIL
else
NUM_CHARS := (word(BUFFER_END^) - TAIL) +
(word(TAIL_PTR^) - word(BUFFER_START^));
(* Convert bytes in buffer to characters *)
NUM_CHARS := NUM_CHARS div 2;
if WAITING_TASK_PTR = nil
then
inc(USER_KEYBOARD_SEM_PTR^)
else
INSERT INTO(READY_QUEUE, WAITING_TASK_PTR;
inc(USER_KEYBOARD_SEM_PTR^, NUM_CHARS - 1);
end; (* if...then *)
end; (* KEYBOARD INTERCEPTOR *)

procedure IRQ_2_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)
(* passes control to the actual interrupt routine. *)

var
(* Pointer to original hardware interrupt handler. *)
VECTOR : pointer;

begin (* IRQ_2 INTERCEPTOR *)
LOCK_SCHEDULER;

VECTOR := HARDWARE_HANDLER($2).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($2));

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

UNLOCK_SCHEDULER;
end; (* IRQ_2_INTERCEPTOR *)

procedure IRQ_3_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)
(* passes control to the actual interrupt routine. *)

var
(* Pointer to original hardware interrupt handler. *)
VECTOR : pointer;

begin (* IRQ_3 INTERCEPTOR *)
LOCK_SCHEDULER;

VECTOR := HARDWARE_HANDLER($3).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($3));

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

UNLOCK_SCHEDULER;
end; (* IRQ_3_INTERCEPTOR *)

procedure IRQ_4_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)
(* passes control to the actual interrupt routine. *)

var
(* Pointer to original hardware interrupt handler. *)
VECTOR : pointer;

begin (* IRQ_4 INTERCEPTOR *)
LOCK_SCHEDULER;

VECTOR := HARDWARE_HANDLER($4).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($4));

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

UNLOCK_SCHEDULER;
end; (* IRQ_4_INTERCEPTOR *)

procedure IRQ_5_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)
(* passes control to the actual interrupt routine. *)

var
(* Pointer to original hardware interrupt handler. *)
VECTOR : pointer;

begin (* IRQ_5 INTERCEPTOR *)
LOCK_SCHEDULER;

VECTOR := HARDWARE_HANDLER($5).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($5));

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

UNLOCK_SCHEDULER;
end; (* IRQ_5_INTERCEPTOR *)

procedure IRQ_6_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)
(* passes control to the actual interrupt routine. *)

var
(* Pointer to original hardware interrupt handler. *)
VECTOR : pointer;

begin (* IRQ_6 INTERCEPTOR *)
LOCK_SCHEDULER;

VECTOR := HARDWARE_HANDLER($6).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($6));

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

UNLOCK_SCHEDULER;
end; (* IRQ_6_INTERCEPTOR *)

procedure IRQ_7_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)
(* passes control to the actual interrupt routine. *)

var
(* Pointer to original hardware interrupt handler. *)
VECTOR : pointer;

begin (* IRQ_7 INTERCEPTOR *)
LOCK_SCHEDULER;

VECTOR := HARDWARE_HANDLER($7).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($7));

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

UNLOCK_SCHEDULER;
end; (* IRQ_7_INTERCEPTOR *)

procedure IRQ_8_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)
(* passes control to the actual interrupt routine. *)

var
(* Pointer to original hardware interrupt handler. *)
VECTOR : pointer;

begin (* IRQ_8 INTERCEPTOR *)
LOCK_SCHEDULER;

VECTOR := HARDWARE_HANDLER($8).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($8));

asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)

UNLOCK_SCHEDULER;
end; (* IRQ_8_INTERCEPTOR *)

procedure IRQ_9_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and *)

```

```

(* passes control to the actual interrupt routine.
..... *)
var (.....)
(* Pointer to original hardware interrupt handler.
..... *)
VECTOR : pointer;
begin (* IRQ_9_INTERCEPTOR *)
LOCK_SCHEDULER;
VECTOR := HARDWARE_HANDLER($9).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($9));
asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)
UNLOCK_SCHEDULER;
end; (* IRQ_9_INTERCEPTOR *)
procedure IRQ_A_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and
* passes control to the actual interrupt routine.
..... *)
var (.....)
(* Pointer to original hardware interrupt handler.
..... *)
VECTOR : pointer;
begin (* IRQ_A_INTERCEPTOR *)
LOCK_SCHEDULER;
VECTOR := HARDWARE_HANDLER($A).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($A));
asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)
UNLOCK_SCHEDULER;
end; (* IRQ_A_INTERCEPTOR *)
procedure IRQ_B_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and
* passes control to the actual interrupt routine.
..... *)
var (.....)
(* Pointer to original hardware interrupt handler.
..... *)
VECTOR : pointer;
begin (* IRQ_B_INTERCEPTOR *)
LOCK_SCHEDULER;
VECTOR := HARDWARE_HANDLER($B).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($B));
asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)
UNLOCK_SCHEDULER;
end; (* IRQ_B_INTERCEPTOR *)
procedure IRQ_C_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and
* passes control to the actual interrupt routine.
..... *)
var (.....)
(* Pointer to original hardware interrupt handler.
..... *)
VECTOR : pointer;
begin (* IRQ_C_INTERCEPTOR *)
LOCK_SCHEDULER;
VECTOR := HARDWARE_HANDLER($C).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($C));
asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)
UNLOCK_SCHEDULER;
end; (* IRQ_C_INTERCEPTOR *)
procedure IRQ_D_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and
* passes control to the actual interrupt routine.
..... *)
var (.....)
(* Pointer to original hardware interrupt handler.
..... *)
VECTOR : pointer;
begin (* IRQ_D_INTERCEPTOR *)
LOCK_SCHEDULER;
VECTOR := HARDWARE_HANDLER($D).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($D));
asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)
UNLOCK_SCHEDULER;
end; (* IRQ_D_INTERCEPTOR *)
procedure IRQ_E_INTERCEPTOR; interrupt;
..... *)
(* Hardware interrupt handler, ensures that preemption is disabled and
* passes control to the actual interrupt routine.
..... *)
var (.....)
(* Pointer to original hardware interrupt handler.
..... *)
VECTOR : pointer;
begin (* IRQ_E_INTERCEPTOR *)
LOCK_SCHEDULER;
VECTOR := HARDWARE_HANDLER($E).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($E));
asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)
UNLOCK_SCHEDULER;
end; (* IRQ_E_INTERCEPTOR *)
procedure IRQ_F_INTERCEPTOR; interrupt;
(* Hardware interrupt handler, ensures that preemption is disabled and
* passes control to the actual interrupt routine.
..... *)
var (.....)
(* Pointer to original hardware interrupt handler.
..... *)
VECTOR : pointer;
begin (* IRQ_F_INTERCEPTOR *)
LOCK_SCHEDULER;
VECTOR := HARDWARE_HANDLER($F).VECTOR;
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.HARDWARE_INTERRUPTS($F));
asm
(* Call original interrupt handler *)
pushf
call dword ptr ss:VECTOR
cli
end; (* asm *)
UNLOCK_SCHEDULER;
end; (* IRQ_F_INTERCEPTOR *)
procedure SW_INTERRUPT_INTERCEPTOR
[.....]
[FLAGS, CS, IP, AX, BX, CX, DX, SI, DI, DS, ES, BP : word;
interrupt;
(* This routine is used to intercept *all* software interrupts (accessed *)
(* via 'INT mn' instruction). By inspecting the actual opcode before the *)
(* return address, the interrupt number is determined and control is *)
(* transferred to that handler. This is done to ensure that all MS-DOS (and *)
(* add-on) software is not preempted (i.e. reentered).
..... *)
const (* Hexadecimal value for Intel 80x86 'INT' instruction.
..... *)
SW_INT_OPCODE = $CD;
(* Hexadecimal value for MS-DOS 'exec()' function.
..... *)
EXEC_FUNCTION_CODE = $4B00;
var (.....)
(* Intel 80x86 register variables used to pass original register *)
(* values to software interrupt handler.
..... *)
REGS : registers;
(* Pointer to position in interrupt vector table where the original *)
(* handler address must be installed prior to passing control to the *)
(* handler (via software interrupt).
..... *)
REVECTOR : pointer absolute $0000:(REVECTORED_INTERRUPT_NUMBER shl 2);
(* Flag that indicates that an call to the MS-DOS function 'exec' *)
(* is being performed, this must be re-entrant to allow parallel *)
(* execution of another program.
..... *)
MSDOS_EXEC_FUNCTION_CALL : boolean;
(* Value to hold the software interrupt number being called.
..... *)
INTERRUPT_NUMBER : byte;
begin (* SW_INTERRUPT_INTERCEPTOR *)
asm
pushf
cli
end; (* asm *)
if mem[CS, _IP - 2] = SW_INT_OPCODE
then
begin
INTERRUPT_NUMBER := mem[CS, _IP - 1];
REVECTOR := INTERRUPT_VECTOR_HANDLER(INTERRUPT_NUMBER).VECTOR;
(* Gather statistics? *)
if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS then
inc(TASKING_STATISTICS.SERVICE_CALLS(INTERRUPT_NUMBER));
(* Allow rescheduling during interrupt? *)
MSDOS_EXEC_FUNCTION_CALL :=
(INTERRUPT_NUMBER = MSDOS_BIOS_FUNCTION_INTERRUPT_NUMBER) and
[AX = EXEC_FUNCTION_CODE];
if MSDOS_EXEC_FUNCTION_CALL
then
inc(NUMBER_OF_CHILD_PROGRAMS)
else
LOCK_SCHEDULER;
with REGS do
begin
AX := AX; BX := BX; CX := CX; DX := DX;
BP := BP; SI := SI; DI := DI;
DS := DS; ES := ES;
FLAGS := FLAGS;
Intf(REVECTORED_INTERRUPT_NUMBER, REGS);
asm cli end;
AX := AX; BX := BX; CX := CX; DX := DX;
BP := BP; SI := SI; DI := DI;
DS := DS; ES := ES;
_FLAGS := FLAGS or INTERRUPT_FLAG;

```



```

end; (* with...do *)
if MSDOS_EXEC_FUNCTION_CALL
then
  dec(NUMBER_OF_CHILD_PROGRAMS)
else
  UNLOCK_SCHEDULER;
end (* if...then *)
else
  with RUNNING_TASK_PTR^ do ERROR_HANDLER(ILLEGAL_OPERATION)(TASK_ID);
asm popf end;
end; (* SW_INTERRUPT_INTERCEPTOR *)

procedure DISPATCH_TASK;
(*-----*)
(* This routine causes the task specified by the RUNNING_TASK_PTR task *)
(* control block to begin running. *)
(*-----*)
const (*-----*)
(* Temporary variable used to switch contexts (all registers!) *)
(* NOTE: MUST be constant to force it into the DSeg! *)
(*-----*)
VALUE : word = 0;

begin (* DISPATCH_TASK *)
asm cli end;
(* Reset preemptive timeslice counter *)
PREEMPTIVE_TIMESLICE := PREEMPTIVE_TIMESLICE -
  ACTUAL_TASKING_CONFIGURATION.TARGET_TIMESLICE;

(* Dispatch a new task *)

(* Switch Stacks *)
VALUE := RUNNING_TASK_PTR^.CONTEXT.SS;
asm mov bx,VALUE end;

VALUE := RUNNING_TASK_PTR^.CONTEXT.SP;
asm
  mov sp,VALUE
  mov ss,bx
end; (* asm *)

(* Setup new context, enabling interrupts *)
VALUE := RUNNING_TASK_PTR^.CONTEXT.FLAGS or INTERRUPT_FLAG;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.CS;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.IP;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.AX;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.BX;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.CX;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.DX;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.BP;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.SI;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.DI;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.EI;
asm push VALUE end;
VALUE := RUNNING_TASK_PTR^.CONTEXT.DS;
asm push VALUE end;

if TIMER_NSEOI_NEEDS_TO_BE_DONE then
begin
  (* Issue NSEOI to PIC *)
  port(PROGRAMMABLE_INTERRUPT_CONTROLLER_PORT) :=
    NON_SPECIFIED_END_OF_INTERRUPT;
  TIMER_NSEOI_NEEDS_TO_BE_DONE := false;
end; (* if...then *)

(* Switch Contexts *)
asm
  pop ds
  pop es
  pop di
  pop si
  pop bp
  pop dx
  pop cx
  pop bx
  pop ax
  iret
end; (* asm *)
end; (* DISPATCH_TASK *)

procedure REQUEUE_AND_RESCHEDULE
(*-----*)
(* TASK : TASK_IDS; RESULTING_STATE : TASK_STATES; *)
(*-----*)
(* This routine is used to preempt a task (which must either be running *)
(* or ready) and put it into the specified queue (either blocked or ready). *)
(* This is useful for 'blocking' a task or for allowing another ready task *)
(* to run (effectively used for context switching). *)
(*-----*)

label (*-----*)
(* Label within procedure for the return (i.e. 'Dispatch') address *)
(* of a task which is going from RUNNING to some other state. *)
(*-----*)
REQUEUE_AND_RESCHEDULE_CONTINUE;

const (*-----*)
(* Speed vs Storage trade-off. The queues are kept in an array so *)
(* that no run-time decisions (case...of) are needed to access the *)
(* proper queue. *)
(*-----*)
QUEUE_OF : array (TASK_STATES) of ^QUEUES =
  (@BLOCKED_QUEUE, @DELAYED_QUEUE, @READY_QUEUE);

var (*-----*)
(* Temporary storage used to determine the contents of the 80x86 *)
(* registers. *)
(*-----*)
VALUE : word;

(*-----*)
(* Temporary pointer to a Task Control Block which 'was/is' running *)
(* but 'is/will be' rescheduled. *)
(*-----*)
RESCHEDULED_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

begin (* REQUEUE_AND_RESCHEDULE *)
asm
  pushf
  cli
end; (* asm *)

(* Don't reschedule a task in a 'locked' region *)
if COMPARE_PRIORITIES(RUNNING_TASK_PTR^.PRIORITY, EQUAL_TO,
  HIGHEST_PRIORITY) then
begin
  asm popf end;
  exit;
end; (* if...then *)

PREVIOUSLY_RUNNING_TASK_ID := RUNNING_TASK_PTR^.TASK_ID;
if RUNNING_TASK_PTR^.TASK_ID = TASK
then
  begin
    (* Save context of the running task *)
    asm
      push ax
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.AX := VALUE;
    asm
      push bx
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.BX := VALUE;
    asm
      push cx
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.CX := VALUE;
    asm
      push dx
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.DX := VALUE;
    asm
      push bp
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.BP := VALUE;
    asm
      push si
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.SI := VALUE;
    asm
      push di
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.DI := VALUE;
    asm
      push es
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.ES := VALUE;
    asm
      push ds
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.DS := VALUE;
    RUNNING_TASK_PTR^.CONTEXT.SS := aseg;
    RUNNING_TASK_PTR^.CONTEXT.SP := aptr;
    asm
      mov ax,seg REQUEUE_AND_RESCHEDULE_CONTINUE
      push ax
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.CS := VALUE;
    asm
      mov ax,offset REQUEUE_AND_RESCHEDULE_CONTINUE
      push ax
      pop VALUE
    end; (* asm *)
    RUNNING_TASK_PTR^.CONTEXT.IP := VALUE;
    RESCHEDULED_TASK_PTR := RUNNING_TASK_PTR;
    end (* if...then *)
  else (* The task must be in the Ready Queue *)
    RESCHEDULED_TASK_PTR := FIND_AND_REMOVE
      (
        TASK,
        (From the) READY_QUEUE,
        UNCONDITIONALLY,
        THERE_IS_NO_BLOCK_TO_MATCH
      );
  if RESCHEDULED_TASK_PTR <> nil
  then
    begin
      RUNNING_TASK_PTR := REMOVE_FROM_HEAD(READY_QUEUE);
      if (RUNNING_TASK_PTR <> nil) and
        (RUNNING_TASK_PTR^.TASK_ID = NULL TASK ID) then
        INSERT_INTO(READY_QUEUE, RUNNING_TASK_PTR);
      INSERT_INTO(QUEUE_OF(RESULTING_STATE)^, RESCHEDULED_TASK_PTR);

      if (RUNNING_TASK_PTR = nil) or
        (RUNNING_TASK_PTR^.TASK_ID = NULL TASK ID) then
        RUNNING_TASK_PTR := REMOVE_FROM_HEAD(READY_QUEUE);
      if (ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS) and
        (RUNNING_TASK_PTR^.TASK_ID <> PREVIOUSLY_RUNNING_TASK_ID) then
        begin
          inc(TASKING_STATISTICS.CONTEXT_SWITCHES);
          if PREEMPTIVE_TIMESLICE <
            ACTUAL_TASKING_CONFIGURATION.TARGET_TIMESLICE
          then
            inc(TASKING_STATISTICS.COOPERATIVE_CONTEXT_SWITCHES);
          end; (* if...then *)

          if ACTUAL_TASKING_CONFIGURATION.PRIORITY_SCHEDULING_POLICY =
            ROTATING_PRIORITIES then
            begin
              if HIGHEST_DYNAMIC_PRIORITY = low(USER_PRIORITIES)
              then
                HIGHEST_DYNAMIC_PRIORITY := high(USER_PRIORITIES)
              else
                HIGHEST_DYNAMIC_PRIORITY := HIGHEST_DYNAMIC_PRIORITY - 1;
              end; (* if...then *)

              DISPATCH_TASK;
            end (* if...then *)
          else
            with RUNNING_TASK_PTR^ do ERROR_HANDLER(TASK_IS_NOT_ACTIVE)(TASK_ID);
          REQUEUE_AND_RESCHEDULE_CONTINUE;
        end; (* REQUEUE_AND_RESCHEDULE *)
      end; (* REQUEUE_AND_RESCHEDULE *)
    end;
  end;
end;

procedure MOUSE_EVENT_HANDLER
(*-----*)
(* (MOUSE : MOUSE_PARAMETER); far; *)
(*-----*)
(* This routine is called by the MS-Mouse driver in response to registered *)
(* mouse actions. The TASKING mouse event is signaled (or the waiting task *)
(* is made ready, whichever is appropriate). *)
(*-----*)

var (*-----*)
(* Pointer to Task Control Block which is waiting on the MS Mouse *)
(* semaphore (may very well be nil). *)
(*-----*)
WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

```

```

begin (* MOUSE_EVENT_HANDLER *)
asm
  pushf
  cli
end: (* asm *)
MOUSE_INFO_FROM_HANDLER := MOUSE;
if USER_MOUSE_EVENT_PTR <> nil then
begin
  WAITING_TASK_PTR := FIND_AND_REMOVE
  (
    ANY_TASK,
    (from the) BLOCKED_QUEUE,
    BECAUSE_THE_BLOCK_WAS_SIGNALLED,
    (AT) addr(USER_MOUSE_EVENT_PTR)
  );
  if WAITING_TASK_PTR = nil
  then
    USER_MOUSE_EVENT_PTR := SIGNALLED
  else
    INSERT_INT0(READY_QUEUE, WAITING_TASK_PTR);
  end: (* if...then *)
asm popf end;
end: (* MOUSE_EVENT_HANDLER *)

procedure CHECK AND SETUP_BLOCK
(
  DELAY_UNTIL_ABSOLUTE_TIME : longint; SEMAPHORE_ADDRESS : pointer;
)
(* This routine is used to establish the currently running task as being *)
(* blocked on a semaphore (for any number of reasons) or as being delayed *)
(* for some time period. These functions are integrated because many of the *)
(* tasking operations require this service. *)
begin (* CHECK_AND_SETUP_BLOCK *)
asm
  pushf
  cli
end: (* asm *)
if (RUNNING_TASK_PTR^.WAITING_FOR <> nil) and
(RUNNING_TASK_PTR^.WAITING_FOR^.SEMAPTR <> nil)
then
begin
  with RUNNING_TASK_PTR^ do
    ERROR_HANDLER(TASK_ALREADY_SUSPENDED)(TASK_ID);
    REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
  end (* if...then *)
else
begin
  if RUNNING_TASK_PTR^.WAITING_FOR = nil then
    new(RUNNING_TASK_PTR^.WAITING_FOR);
  with RUNNING_TASK_PTR^.WAITING_FOR^ do
    begin
      ABSOLUTE_TIME := DELAY_UNTIL_ABSOLUTE_TIME;
      SEMPTR := SEMAPHORE_ADDRESS;
      if SEMPTR <> nil
      then
        REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, BLOCKED)
      else
        if ABSOLUTE_TIME <> NO_TIME_DELAY
        then
          REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, DELAYED)
        else
          with RUNNING_TASK_PTR^ do
            ERROR_HANDLER(ILLEGAL_OPERATION)(TASK_ID);
            REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
          end: (* with...do *)
        end: (* if...then...else *)
      asm popf end;
    end: (* CHECK_AND_SETUP_BLOCK *)
  end: (* if...then *)
end: (* CHECK_AND_SETUP_BLOCK *)

procedure CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST
(
  SEMAPHORE_ADDRESS : pointer;
)
(* This routine is used (when priority inheritance is enabled) to check to *)
(* see if the priority of a task needs to be promoted as a result of a task *)
(* attempting to lock a resource that is already locked. If the resource is *)
(* being locked for the first time then the resource is added to the linked *)
(* list of resources. *)
var
  (* Pointer to locked resource that is 'walked' through the list of *)
  (* all locked resources. *)
  WALKING_PTR : LOCKED_RESOURCE_PTR;
  (* Pointer to Task Control Block for the task that is having its *)
  (* priority promoted (as a result of priority inheritance). *)
  PROMOTED_TASK_PTR : TASK_CONTROL_BLOCK_PTR;
begin (* CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST *)
asm
  pushf
  cli
end: (* asm *)
WALKING_PTR := LOCKED_RESOURCE_LIST_PTR;
while (WALKING_PTR <> nil) and
(WALKING_PTR^.SEMAPTR <> SEMAPHORE_ADDRESS) do
  WALKING_PTR := WALKING_PTR^.NEXT;
if (WALKING_PTR <> nil) and
(WALKING_PTR^.SEMAPTR = SEMAPHORE_ADDRESS)
then
begin
  if WALKING_PTR^.OWNER = nil
  then
    with WALKING_PTR^ do
      begin
        OWNER := RUNNING_TASK_PTR;
        PRIORITY := RUNNING_TASK_PTR^.PRIORITY;
      end (* with...do *)
    else
      (* Promote priority? *)
      if COMPARE_PRIORITIES(WALKING_PTR^.OWNER^.PRIORITY, LESS_THAN,
        RUNNING_TASK_PTR^.PRIORITY) then
        begin
          PROMOTED_TASK_PTR := FIND_AND_REMOVE
          (
            WALKING_PTR^.OWNER^.TASK_ID,
            (from the) READY_QUEUE,
            UNCONDITIONALLY,
            THERE_IS_NO_BLOCK_TO_MATCH
          );
          PROMOTED_TASK_PTR^.PRIORITY := RUNNING_TASK_PTR^.PRIORITY;
          INSERT_INT0(READY_QUEUE, PROMOTED_TASK_PTR);
        end: (* if...then *)
        end: (* if...then...else *)
      end: (* if...then *)
    else
      (* Find place to insert list element *)
      WALKING_PTR := LOCKED_RESOURCE_LIST_PTR;
      if WALKING_PTR = nil
      then
        begin
          new(WALKING_PTR);
          LOCKED_RESOURCE_LIST_PTR := WALKING_PTR;
        end (* if...then *)
      else
        while WALKING_PTR^.NEXT <> nil do
          WALKING_PTR := WALKING_PTR^.NEXT;
        end: (* if...then...else *)
        (* Add new locked resource *)
        with WALKING_PTR^ do
          begin
            OWNER := RUNNING_TASK_PTR;
            PRIORITY := RUNNING_TASK_PTR^.PRIORITY;
            SEMAPHORE_PTR := SEMAPHORE_ADDRESS;
            NEXT := nil;
          end: (* with...do *)
        end: (* if...then...else *)
      asm popf end;
    end: (* CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST *)
  end: (* CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST *)
end: (* CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST *)

procedure FIND_AND_REMOVE_LOCKED_RESOURCE
(
  SEMAPHORE_ADDRESS : pointer;
)
(* This routine is used to remove the ownership of the resource from any *)
(* task. This is necessary after the resource has been 'unlocked' and is *)
(* therefore no longer 'owned'. *)
var
  (* Pointer to locked resource that is 'walked' through the list of *)
  (* all locked resources. *)
  WALKING_PTR : LOCKED_RESOURCE_PTR;
begin (* FIND_AND_REMOVE_LOCKED_RESOURCE *)
asm
  pushf
  cli
end: (* asm *)
WALKING_PTR := LOCKED_RESOURCE_LIST_PTR;
while (WALKING_PTR <> nil) and
(WALKING_PTR^.SEMAPTR <> SEMAPHORE_ADDRESS) do
  WALKING_PTR := WALKING_PTR^.NEXT;
if (WALKING_PTR <> nil) and (WALKING_PTR^.OWNER <> nil) and
(WALKING_PTR^.SEMAPTR = SEMAPHORE_ADDRESS) then
begin
  if COMPARE_PRIORITIES(RUNNING_TASK_PTR^.PRIORITY, GREATER_THAN,
    WALKING_PTR^.OWNER^.PRIORITY) and
    (WALKING_PTR^.OWNER := RUNNING_TASK_PTR) then
    RUNNING_TASK_PTR^.PRIORITY := WALKING_PTR^.PRIORITY;
  WALKING_PTR^.OWNER := nil;
end: (* if...then *)
asm popf end;
end: (* FIND_AND_REMOVE_LOCKED_RESOURCE *)

(*-----*)
(* End of non-user functions *)
(*-----*)
(* User functions *)
(*-----*)

procedure GETMEM
(
  PTR : pointer; SIZE : word;
)
(* This procedure replaces the standard Turbo Pascal memory allocation *)
(* procedure by the same name. Since TASKING must perform all memory *)
(* allocation and de-allocation, this routine "must" be used instead of the *)
(* system.getmem() procedure. If 'type P : ^T;' is declared, then the *)
(* correct usage is 'getmem(pointer(P), sizeof(T));'. Within TASKING, the *)
(* system.new() procedure can also be used because it is always used when *)
(* interrupts are disabled. *)
begin (* GETMEM *)
asm
  pushf
  cli
end: (* asm *)
system.getmem(PTR, SIZE);
asm popf end;
end: (* GETMEM *)

procedure FREEMEM
(
  PTR : pointer; SIZE : word;
)
(* This procedure replaces the standard Turbo Pascal memory allocation *)
(* procedure by the same name. Since TASKING must perform all memory *)
(* allocation and de-allocation, this routine "must" be used instead of the *)
(* system.freemem() procedure. If 'type P : ^T;' is declared, then the *)
(* correct usage is 'freemem(pointer(P), sizeof(T));'. Within TASKING, the *)
(* system.dispose() procedure can also be used because it is always used *)
(* when interrupts are disabled. *)
begin (* FREEMEM *)
asm
  pushf
  cli
end: (* asm *)
system.freemem(PTR, SIZE);
asm popf end;
end: (* FREEMEM *)

procedure CREATE
(
  TASK : TASK_IDS; ATTR : TASK_ATTRIBUTES; ENTRY_POINT : TASK_TYPE;
)
(* This routine creates all system data structures for a parallel *)
(* executing task. The task is made ready to run and will be available for *)
(* execution. Actual execution will occur when the scheduler activates the *)
(* task (i.e. precise execution time is unknown and depends on availability *)
(* of system resources and system load. *)
const
  (* When a task implicitly terminates there must be room for calls *)
  (* to routines necessary for proper TASKING clean-up. This constant *)
  (* must be larger than the maximum bytes needed for these clean-up *)
  (* routines to execute properly. *)
  STACK_HEADROOM = 48;
var
  (* Pointer to the to-be-created Task Control Block. *)
  NEW_TCB_PTR : TASK_CONTROL_BLOCK_PTR;

```

```

(* Pointer to the top of the task stack (as opposed to the first *)
(* byte of the stack which is used to reference the Turbo Pascal *)
(* variable returned when the memory is allocated. *)
TOP_OF_STACK_PTR : pointer;

(* Loop control variable of possible errors used to initialize the *)
(* task specific error handlers. *)
ERROR : SYSTEM_ERRORS;

function NEW_TASK_ID : TASK_IDS;
(* This function determines what the next available task identifier is, *)
(* marks it as allocated and returns it to the caller. *)
var
  (* Temporary variable used to hold the potential new task ID. *)
  THE_NEW_TASK_ID : TASK_IDS;
  (* Loop control variable to search arrays of TASK_ID sets. *)
  TASK_ID_SET : TASK_ID_SETS;
begin
  (* NEW_TASK_ID *)
  TASK_ID_SET := low(TASK_ID_SETS);
  THE_NEW_TASK_ID := succ(NULL_TASK_ID);
  while (THE_NEW_TASK_ID in ALLOCATED_TASK_IDS(TASK_ID_SET)) and
    (THE_NEW_TASK_ID <= high(byte)) and
    (TASK_ID_SET <= high(TASK_ID_SETS)) do
    begin
      inc(THE_NEW_TASK_ID);
      if (THE_NEW_TASK_ID > high(byte)) and
        (TASK_ID_SET < high(TASK_ID_SETS)) then
        begin
          TASK_ID_SET := succ(TASK_ID_SET);
          THE_NEW_TASK_ID := low(TASK_IDS);
        end;
      (* if...then *)
    end;
    (* while...do *)
  if not (THE_NEW_TASK_ID in ALLOCATED_TASK_IDS(TASK_ID_SET))
  then
    begin
      ALLOCATED_TASK_IDS[TASK_ID_SET] :=
        ALLOCATED_TASK_IDS[TASK_ID_SET] + (THE_NEW_TASK_ID);
      NEW_TASK_ID := (TASK_ID_SET shl 8) or THE_NEW_TASK_ID;
    end;
  else
    with RUNNING_TASK_PTR do
      ERROR_HANDLER[INSUFFICIENT_RESOURCES](TASK_ID);
  end;
  (* NEW_TASK_ID *)

function INC_PTR
(* (THE_PTR : pointer; INCREMENT : word) : pointer; *)
(* This routine merely increments a double word pointer by the value *)
(* specified. *)
var
  (* Temporary variables to hold intermediate result. *)
  SEGMENT, OFFSET : word;
  SUM_OFS : longint;
begin
  (* INC_PTR *)
  SUM_OFS := ofs(THE_PTR) + INCREMENT;
  SEGMENT := seg(THE_PTR);
  OFFSET := SUM_OFS and $0000FFFF;
  if (SUM_OFS and $FFFF0000) <> 0 then
    (* Overflowed 64k boundary by 1 (because INCREMENT is a word!) *)
    inc(SEGMENT, 1600);
  INC_PTR := ptr(SEGMENT, OFFSET);
end;
(* INC_PTR *)

begin
  (* CREATE *)
  asm
    pushf
    cll
  end;
  (* asm *)
  TASK := NEW_TASK_ID;
  new(NEW_TCB_PTR);
  NEW_TCB_PTR^.STACK_SIZE := ATTR.STACK_WORDS_NEEDED * 2 (* Bytes/word *);
  with NEW_TCB_PTR do
    begin
      getmem(STACK_PTR, STACK_SIZE);
      (* Get pointer to Top-Of-Stack *)
      TOP_OF_STACK_PTR := INC_PTR(STACK_PTR, STACK_SIZE-12-STACK_HEADROOM);
      TASK_ID := TASK;
      CLOCK_TICKS := 0;
      PRIORITY := ATTR.PRIORITY;
      WAITING_FOR := nil;
      NEXT := nil;
      PREVIOUS := nil;
      for ERROR := low(SYSTEM_ERRORS) to high(SYSTEM_ERRORS) do
        if addr(ATTR.ERROR_HANDLERS[ERROR]) = nil
        then
          ERROR_HANDLER[ERROR] := NULL.ERROR_HANDLER[ERROR]
        else
          ERROR_HANDLER[ERROR] := ATTR.ERROR_HANDLERS[ERROR];
      with CONTEXT do
        begin
          AX := $0000; BX := $0000; CX := $0000; DX := $0000;
          BP := $0000; SI := $0000; DI := $0000; ES := $0000;
          DS := dseg;
          SS := seg(TOP_OF_STACK_PTR);
          SP := (ofs(TOP_OF_STACK_PTR) + 1) and $FFFF;
          CS := seg(ENTRY_POINT); IP := ofs(ENTRY_POINT);
          FLAGS := INTERRUPT_FLAG;
        end;
        (* with...do *)
        (* Clear the stack *)
        fillchar(STACK_PTR, STACK_SIZE-1, byte(DEFAULT_STACK_VALUE));
        (* Install task identifier and priority onto stack *)
        mems[CONTEXT.SS + CONTEXT.SP + 4] := PRIORITY;
        mems[CONTEXT.SS + CONTEXT.SP + 6] := TASK_ID;
        (* Install implied call to DESTROY() onto task stack *)
        mems[CONTEXT.SS + CONTEXT.SP + 0] := ofs(DESTROY);
        mems[CONTEXT.SS + CONTEXT.SP + 2] := seg(DESTROY);
        mems[CONTEXT.SS + CONTEXT.SP + 12] := TASK_ID;
      end;
      (* with...do *)
      INSERT_INTO(READY_QUEUE, NEW_TCB_PTR);
    end;
  if ALLOW_RESCHEDULE IN CREATE then
    REQUEST_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
  asm popf end;
end;
(* CREATE *)

procedure DESTROY
(* (TASK : TASK_IDS); *)
(* This routine terminates (if appropriate) releases all of the task's *)
(* resources to be used by other tasks and then removes all evidence of the *)
(* task ever having existed. An implied call to this routine is setup at *)
(* the time that a task is created (by setting up the stack to return here). *)
(* This implied call ensures that a task that reaches its 'end' statement is *)
(* properly terminated and does not become an orphan (or a zombie). *)
const
  (* Pointer to the Task Control Block that is being destroyed. *)
  (* NOTE: MUST be constant to force it into the DSeg *)
  DESTROYED_TASK_PTR : TASK_CONTROL_BLOCK_PTR = nil;
  (* Flag to indicate that the task that is being destroyed 'was/is' *)
  (* running, this has stack access consequences. *)
  (* NOTE: MUST be constant to force it into the DSeg *)
  DESTROYED_TASK_WAS_RUNNING : boolean = false;
  (* Temporary variable used to switch stacks (SS and SP). *)
  (* NOTE: MUST be constant to force it into the DSeg *)
  VALUE : word = 0;
procedure FREE_TASK_ID
(* (TASK_ID : TASK_IDS); *)
(* This routine merely puts a task identifier back into the free list so *)
(* NOTE: MUST be constant to force it into the DSeg *)
begin
  (* FREE_TASK_ID *)
  ALLOCATED_TASK_IDS[hi(TASK_ID) shr 8] :=
    ALLOCATED_TASK_IDS[hi(TASK_ID) shr 8] - (lo(TASK_ID));
end;
(* FREE_TASK_ID *)

begin
  (* DESTROY *)
  asm
    pushf
    cll
  end;
  (* asm *)
  (* The null task cannot be destroyed *)
  if TASK = NULL_TASK_ID then
    begin
      RUNNING_TASK_PTR^.ERROR_HANDLER[ILLEGAL_OPERATION](TASK);
      asm popf end;
      exit;
    end;
  (* if...then *)
  if TASK = RUNNING_TASK_PTR^.TASK_ID
  then
    begin
      DESTROYED_TASK_PTR := RUNNING_TASK_PTR;
      DESTROYED_TASK_WAS_RUNNING := true;
    end;
  else
    begin
      DESTROYED_TASK_PTR := FIND_AND_REMOVE
        (
          TASK,
          (From the) READY_QUEUE,
          UNCONDITIONALLY,
          THERE_IS_NO_BLOCK_TO_MATCH
        );
      if DESTROYED_TASK_PTR = nil then
        DESTROYED_TASK_PTR := FIND_AND_REMOVE
          (
            TASK,
            (From the) BLOCKED_QUEUE,
            UNCONDITIONALLY,
            THERE_IS_NO_BLOCK_TO_MATCH
          );
      if DESTROYED_TASK_PTR = nil then
        DESTROYED_TASK_PTR := FIND_AND_REMOVE
          (
            TASK,
            (From the) DELAYED_QUEUE,
            UNCONDITIONALLY,
            THERE_IS_NO_BLOCK_TO_MATCH
          );
    end;
  (* if...then *)
  DESTROYED_TASK_WAS_RUNNING := false;
end;
(* if...then..else *)

if DESTROYED_TASK_PTR = nil
then
  begin
    (* Can't find the requested task ID *)
    RUNNING_TASK_PTR^.ERROR_HANDLER[ILLEGAL_TASK_ID](TASK);
    asm popf end;
    exit;
  end;
  (* if...then *)
  begin
    (* Check to see if special control structures need to be removed. *)
    if DESTROYED_TASK_PTR^.WAITING_FOR <> nil then
      begin
        if (USER_KEYBOARD_SEM_PTR <> nil) and
          (DESTROYED_TASK_PTR^.WAITING_FOR^.SEM_PTR =
            USER_KEYBOARD_SEM_PTR)
        then
          begin
            dispose(USER_KEYBOARD_SEM_PTR);
            USER_KEYBOARD_SEM_PTR := nil;
          end;
        (* if...then *)
        else
          if (USER_MOUSE_EVENT_PTR <> nil) and
            (DESTROYED_TASK_PTR^.WAITING_FOR^.SEM_PTR =
              SEMAPHORE_PTR(USER_MOUSE_EVENT_PTR)) then
            begin
              dispose(USER_MOUSE_EVENT_PTR);
              USER_MOUSE_EVENT_PTR := nil;
            end;
            (* if...then *)
          (* Get rid of whatever the task was waiting for *)
          dispose(DESTROYED_TASK_PTR^.WAITING_FOR);
          end;
        (* if...then *)
      end;
    (* De-allocate the task's mailbox (if applicable) *)
    REMOVE_MBOX_OF(DESTROYED_TASK_PTR^.TASK_ID);
    (* Switch stacks if user stack is going to be destroyed *)
    if DESTROYED_TASK_WAS_RUNNING then
      begin
        RUNNING_TASK_PTR := REMOVE_FROM_HEAD(READY_QUEUE);
        VALUE := RUNNING_TASK_PTR^.CONTEXT.SS;
        asm
          push VALUE
          pop ax
        end;
        VALUE := RUNNING_TASK_PTR^.CONTEXT.SP;
        asm mov sp, VALUE end;
        (* if...then *)
      end;
    end;
  end;
end;
(* if...then *)

```

```

(* Remove task structures *)
FREE_TASK_ID(DESTROYED_TASK_PTR^.TASK_ID);
with DESTROYED_TASK_PTR do freemem(STACK_PTR, STACK_SIZE);
dispose(DESTROYED_TASK_PTR);

if DESTROYED_TASK_WAS_RUNNING then
  (* Cannot come back here because stack is gone' *)
  asm tmp DISPATCH_TASK end;
end; (* if...then...else *)
end; (* DESTROY *)

procedure SUSPEND
(*-----*)
(TASK : TASK_IDS);
(* This routine causes a previously running or ready task to be blocked *)
(* unconditionally. *)
(*-----*)
begin (* SUSPEND *)
  asm
    pushf
    cli
  end; (* asm *)
  REQUEUE_AND_RESCHEDULE(TASK, BLOCKED);
  asm popf end;
end; (* SUSPEND *)

procedure RESUME
(*-----*)
(TASK : TASK_IDS);
(* This routine causes a previously suspended task to be made ready. The *)
(* task is not necessarily the next to run, that is based on the priorities *)
(* of the other ready tasks. *)
(*-----*)
var
  (*-----*)
  (* Pointer to the Task Control Block of the task that is being made *)
  (* READY. *)
  (*-----*)
  RESUMED_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

begin (* RESUME *)
  asm
    pushf
    cli
  end; (* asm *)
  RESUMED_TASK_PTR := FIND_AND_REMOVE
  (
    TASK,
    (From the) BLOCKED_QUEUE,
    UNCONDITIONALLY,
    THERE_IS_NO_BLOCK_TO_MATCH
  );
  if RESUMED_TASK_PTR <> nil
  then
    begin
      INSERT INTO(READY_QUEUE, RESUMED_TASK_PTR);
      REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
    end (* if...then *)
  else
    if RUNNING_TASK_PTR^.TASK_ID <> NULL_TASK_ID then
      with RUNNING_TASK_PTR do
        ERROR_HANDLER[TASK_IS_NOT_ACTIVE](TASK_ID);
      end;
    asm popf end;
  end; (* RESUME *)

function GET_MILLISECOND_TICKS : longint;
(*-----*)
(* This routine returns the number of milliseconds that have elapsed since *)
(* the program started executing. *)
(*-----*)
begin (* GET_MILLISECOND_TICKS *)
  asm
    pushf
    cli
  end; (* asm *)
  GET_MILLISECOND_TICKS := MILLISECOND_TICKS;
  asm popf end;
end; (* GET_MILLISECOND_TICKS *)

procedure WAIT_FOR_DELAY
(*-----*)
(DURATION : TIME);
(* This system call causes the currently running task to become blocked *)
(* for the specified length of time. There is no guarantee that the task *)
(* will begin executing when the time expires, the only guarantee is that *)
(* the task will become ready to execute at that time. *)
(* No attempt is made to detect 'Global Ticks' rollover, as a result of *)
(* this after the program has been running for a long time ((2^30-1)* msec *)
(* i.e. 12.4 days) there is a possibility of a delay request resulting in *)
(* an immediate rescheduling. *)
(*-----*)
var
  (*-----*)
  (* Temporary variable used to hold the MILLISECOND_TICKS value when *)
  (* the task will be 'awakened'. *)
  (*-----*)
  ASSOLUTE_TIME : longint;

begin (* WAIT_FOR_DELAY *)
  asm
    pushf
    cli
  end; (* asm *)
  with DURATION do
    ASSOLUTE_TIME := MILLISECOND_TICKS +
      trunc(longint(DAYS) * MILLISECONDS_PER_DAY +
        longint(HOURS) * MILLISECONDS_PER_HOUR +
        longint(MINUTES) * MILLISECONDS_PER_MINUTE +
        longint(SECONDS) * MILLISECONDS_PER_SECOND +
        longint(MILLISECONDS) * MILLISECONDS_PER_SECOND +
        1 (* Ensure that we always 'round' up *)
      );
  CHECK_AND_SETUP_BLOCK(ASSOLUTE_TIME, nil);
  asm popf end;
end; (* WAIT_FOR_DELAY *)

procedure PREEMPTASLE_DELAY
(*-----*)
(DURATION : TIME);
(* This system call causes the currently running task to be delayed by the *)
(* specified number of milliseconds. *)
(*-----*)
var
  (*-----*)
  (* Temporary variable used to hold the MILLISECOND_TICKS value when *)
  (* the task will be allowed to continue executing. *)
  (*-----*)
  DELAY_COMPLETE_TIME : longint;

begin (* PREEMPTASLE_DELAY *)
  asm
    pushf
    cli
  end; (* asm *)
  with DURATION do
    DELAY_COMPLETE_TIME := MILLISECOND_TICKS +
      trunc
      (
        longint(DAYS) * MILLISECONDS_PER_DAY +
        longint(HOURS) * MILLISECONDS_PER_HOUR +
        longint(MINUTES) * MILLISECONDS_PER_MINUTE +
        longint(SECONDS) * MILLISECONDS_PER_SECOND +
        longint(MILLISECONDS) * MILLISECONDS_PER_SECOND +
        1 (* Ensure that we always 'round' up *)
      );
  repeat (* Nothing! *) until MILLISECOND_TICKS >= DELAY_COMPLETE_TIME;
  end; (* PREEMPTASLE_DELAY *)

procedure CHANGE_PRIORITY
(*-----*)
(PRIORITY : USER_PRIORITIES);
(* This routine simply changes the priority of the currently running task *)
(* and preempts the running task (causing a dispatch of the now highest *)
(* priority task) if appropriate. *)
(*-----*)
begin (* CHANGE_PRIORITY *)
  asm
    pushf
    cli
  end; (* asm *)
  RUNNING_TASK_PTR^.PRIORITY := PRIORITIES(PRIORITY);
  REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
  end; (* CHANGE_PRIORITY *)

procedure SIGNAL_SEMAPHORE
(*-----*)
(var SEM : SEMAPHORE);
(* This routine performs an Up() operation on the specified semaphore. If *)
(* there are tasks waiting on the semaphore then one is awakened. If there *)
(* are none waiting then the semaphore is incremented. *)
(*-----*)
var
  (*-----*)
  (* Pointer to Task Control Block which is waiting on the semaphore *)
  (* (may very well be nil). *)
  (*-----*)
  WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

begin (* SIGNAL_SEMAPHORE *)
  asm
    pushf
    cli
  end; (* asm *)
  if TASKING_CONFIGURATION.PRIORITY_INHERITANCE_ENABLED then
    FIND_AND_REMOVE_LOCKED_RESOURCE(addr(SEM));
    WAITING_TASK_PTR := FIND_AND_REMOVE
    (
      ANY_TASK,
      (From the) BLOCKED_QUEUE,
      BECAUSE_THE_BLOCK_WAS_SIGNALLED,
      (At) addr(SEM)
    );
  if WAITING_TASK_PTR = nil
  then
    inc(SEM)
  else
    INSERT INTO(READY_QUEUE, WAITING_TASK_PTR);
    REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
  end; (* SIGNAL_SEMAPHORE *)

procedure WAIT_ON_SEMAPHORE
(*-----*)
(var SEM : SEMAPHORE);
(* This routine performs a Down() operation on the specified semaphore. *)
(* If the semaphore is zero then the running task is suspended. If it is *)
(* non-zero then the semaphore is decremented. *)
(*-----*)
begin (* WAIT_ON_SEMAPHORE *)
  asm
    pushf
    cli
  end; (* asm *)
  if SEM > 0
  then
    begin
      dec(SEM);
      if TASKING_CONFIGURATION.PRIORITY_INHERITANCE_ENABLED and (SEM = 0)
      then
        CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST(addr(SEM));
        REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
      end (* if...then *)
    else
      begin
        if TASKING_CONFIGURATION.PRIORITY_INHERITANCE_ENABLED then
          CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST(addr(SEM));
          CHECK_AND_SETUP_BLOCKING_TIME_DELAY, addr(SEM);
        end; (* if...then...else *)
      end;
    end; (* WAIT_ON_SEMAPHORE *)

procedure SIGNAL_EVENT
(*-----*)
(var THE_EVENT : EVENT);
(* This routine signals the event specified, if there are tasks waiting on *)
(* the event then one task is made ready. If there are no tasks waiting on *)
(* the event then the signal is saved, only one signal is maintained no *)
(* matter how many times the event is signaled. *)
(*-----*)
var
  (*-----*)
  (* Pointer to Task Control Block which is waiting on the event (may *)
  (* very well be nil). *)
  (*-----*)
  WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

begin (* SIGNAL_EVENT *)
  asm
    pushf
    cli
  end; (* asm *)
  if TASKING_CONFIGURATION.PRIORITY_INHERITANCE_ENABLED then
    FIND_AND_REMOVE_LOCKED_RESOURCE(addr(THE_EVENT));
    WAITING_TASK_PTR := FIND_AND_REMOVE
    (
      ANY_TASK,
      (From the) BLOCKED_QUEUE,
      BECAUSE_THE_BLOCK_WAS_SIGNALLED,
      (At) addr(THE_EVENT)
    );
  if WAITING_TASK_PTR = nil
  then
    THE_EVENT := SIGNALLED
  else
    INSERT INTO(READY_QUEUE, WAITING_TASK_PTR);
  end;

```

```

    REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
asm popfd end;
end; (* SIGNAL_EVENT *)

procedure BROADCAST_EVENT
(*****);
(var THE_EVENT : EVENT);
(-----)
(* This routine signals the event specified, if there are tasks waiting on *)
(* the event then they are all made ready. If there are no tasks waiting on *)
(* the event then the signal is saved, only one signal is maintained no *)
(* matter how many times the event is signalled. *)
(-----)
var
(*****);
(* Pointer to Task Control Block which is waiting on the event (may *)
(* very well be nil). *)
(-----)
WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

begin (* BROADCAST_EVENT *)
asm
    pushf
    clli
end; (* asm *)
if TASKING_CONFIGURATION.PRIDRITY_INHERITANCE_ENABLED then
    FIND_AND_REMOVE_LOCKED_RESOURCE(addr(THE_EVENT));
repeat
    WAITING_TASK_PTR := FIND_AND_REMOVE
    (
        ANY_TASK,
        (FROM THE) BLOCKED_QUEUE,
        BECAUSE_THE_BLOCK_WAS_SIGNALLED,
        (AT) addr(THE_EVENT)
    );
    if WAITING_TASK_PTR <> nil then
        INSERT_INTO(READY_QUEUE, WAITING_TASK_PTR);
until WAITING_TASK_PTR = nil;
REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
asm popfd end;
end; (* BROADCAST_EVENT *)

procedure START_PERIODIC_EVENT
(*****);
(var THE_EVENT : EVENT; INTERVAL : TIME);
(-----)
(* This routine starts the periodic signalling (by TASKING) of the event *)
(* specified at the specified interval. This routine cannot cause a task to *)
(* be suspended but a re-schedule may occur. *)
(-----)
var
(*****);
(* Pointer to the list of periodic events (possibly nil). *)
(-----)
WALKING_PTR : PERIODIC_EVENTS_PTR;

(*****);
(* Amount of time between successive signals on the event. *)
(-----)
PERIOD_INTERVAL : longint;

begin (* START_PERIODIC_EVENT *)
asm
    pushf
    clli
end; (* asm *)

(* Compute period (in milliseconds) *)
with INTERVAL do
    PERIOD_INTERVAL := trunc(longint(DAYS) * MILLISECONDS_PER_DAY +
        longint(HOURS) * MILLISECONDS_PER_HOUR +
        longint(MINUTES) * MILLISECONDS_PER_MINUTE +
        longint(SECONDS) * MILLISECONDS_PER_SECOND +
        longint(MILLISECONDS));

(* Find the place in the list to insert the event *)
if PERIODIC_EVENTS_LIST_PTR = nil
then
    begin
        (* Insert element (i.e. creates list) *)
        new(WALKING_PTR);
        with WALKING_PTR do
            begin
                EVENT_PTR := addr(THE_EVENT);
                PERIOD := PERIOD_INTERVAL;
                NEXT := nil;
            end; (* with...do *)
            PERIODIC_EVENTS_LIST_PTR := WALKING_PTR;
        end (* if...then *)
    else
        begin
            WALKING_PTR := PERIODIC_EVENTS_LIST_PTR;
            while (WALKING_PTR^.EVENT_PTR <> addr(THE_EVENT)) and
                (WALKING_PTR^.NEXT <> nil) do
                WALKING_PTR := WALKING_PTR^.NEXT;

            (* There can only be one period for each event *)
            if (WALKING_PTR^.EVENT_PTR = addr(THE_EVENT))
            then
                with RUNNING_TASK_PTR do
                    ERROR_HANDLER(ILLEGAL_OPERATION)(TASK_ID)
                else
                    begin
                        (* Insert element (always at the end of the list!) *)
                        new(WALKING_PTR^.NEXT);
                        WALKING_PTR := WALKING_PTR^.NEXT;
                        with WALKING_PTR do
                            begin
                                EVENT_PTR := addr(THE_EVENT);
                                PERIOD := PERIOD_INTERVAL;
                                NEXT := nil;
                            end; (* with...do *)
                            end; (* if...then..else *)
                        end; (* if...then..else *)

                    REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
                    asm popfd end;
                    end; (* START_PERIODIC_EVENT *)

procedure STOP_PERIODIC_EVENT
(*****);
(var THE_EVENT : EVENT);
(-----)
(* This routine stops the periodic signalling (by TASKING) of the event *)
(* specified. This routine cannot cause a task to be suspended but a *)
(* re-schedule may occur. *)
(-----)
var
(*****);
(* Pointer to the list of periodic events (possibly nil) *)
(-----)
WALKING_PTR, PREVIOUS_PTR : PERIODIC_EVENTS_PTR;

begin (* STOP_PERIODIC_EVENT *)
asm
    pushf
    clli
end; (* asm *)

(* Find the place in the list to delete the event *)
PREVIOUS_PTR := nil;
WALKING_PTR := PERIODIC_EVENTS_LIST_PTR;
while (WALKING_PTR <> nil) and
    (WALKING_PTR^.EVENT_PTR = addr(THE_EVENT)) do
    begin
        PREVIOUS_PTR := WALKING_PTR;
        WALKING_PTR := WALKING_PTR^.NEXT;
        end; (* while...do *)

    (* There must be one period for the event *)
    if (PREVIOUS_PTR = nil) or (PREVIOUS_PTR^.EVENT_PTR <> addr(THE_EVENT))
    then
        with RUNNING_TASK_PTR do
            ERROR_HANDLER(ILLEGAL_OPERATION)(TASK_ID)
        else
            begin
                if PREVIOUS_PTR <> nil
                then
                    PREVIOUS_PTR := WALKING_PTR^.NEXT
                else
                    PERIODIC_EVENTS_LIST_PTR := nil;
                    dispose(WALKING_PTR);
                    end; (* if...then..else *)

                REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
                asm popfd end;
                end; (* STOP_PERIODIC_EVENT *)

procedure WAIT_ON_EVENT
(*****);
(var THE_EVENT : EVENT);
(-----)
(* This routine causes the calling task to be suspended waiting for the *)
(* specified event to occur. If the event has already occurred then the *)
(* calling task will immediately become ready, although a re-schedule may *)
(* occur. *)
(-----)
begin (* WAIT_ON_EVENT *)
asm
    pushf
    clli
end; (* asm *)
if THE_EVENT = SINGALED
then
    begin
        THE_EVENT := UNSINGALED;
        if TASKING_CONFIGURATION.PRIORITY_INHERITANCE_ENABLED then
            CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST(addr(THE_EVENT));
        REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
        end (* if...then *)
    else
        begin
            if TASKING_CONFIGURATION.PRIORITY_INHERITANCE_ENABLED then
                CHECK_OR_ADD_RESOURCE_TO_LOCKED_LIST(addr(THE_EVENT));
            CHECK_AND_SETUP_BLOCK(INO_TIME_DELAY, addr(THE_EVENT));
            end; (* if...then..else *)
        asm popfd end;
        end; (* WAIT_ON_EVENT *)

procedure SIGNAL_BINARY_SEMAPHORE
(*****);
(var SEM : BINARY_SEMAPHORE);
(-----)
(* This routine performs an Up() operation on the specified semaphore. If *)
(* there are tasks waiting on the semaphore then one is awakened. If there *)
(* are none waiting then the semaphore is set. *)
(-----)
begin (* SIGNAL_BINARY_SEMAPHORE *)
asm
    pushf
    clli
end; (* asm *)
if SEM = 1
then
    begin
        with RUNNING_TASK_PTR do
            ERROR_HANDLER(ILLEGAL_OPERATION)(TASK_ID);
        REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
        end (* if...then *)
    else
        SIGNAL_EVENT(EVENT(SEM));
    asm popfd end;
    end; (* SIGNAL_BINARY_SEMAPHORE *)

procedure WAIT_ON_BINARY_SEMAPHORE
(*****);
(var SEM : BINARY_SEMAPHORE);
(-----)
(* This routine performs a Down() operation on the specified semaphore. *)
(* If the semaphore is zero then the running task is suspended. If it is *)
(* non-zero then the semaphore is cleared. *)
(-----)
begin (* WAIT_ON_BINARY_SEMAPHORE *)
    WAIT_ON_EVENT(EVENT(SEM));
end; (* WAIT_ON_BINARY_SEMAPHORE *)

procedure SIGNAL_CONDITION_VARIABLE
(*****);
(var C_VAR : CONDITION_VARIABLE);
(-----)
(* This routine signals the condition variable specified, if there is a *)
(* task waiting on the condition variable then it is made ready (or one of *)
(* the multiple waiting tasks is made ready. If there are no tasks waiting *)
(* on the condition variable then the signal is lost. *)
(-----)
var
(*****);
(* Pointer to Task Control Block which is waiting on the condition *)
(* variable (may very well be nil). *)
(-----)
WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;

begin (* SIGNAL_CONDITION_VARIABLE *)
asm
    pushf
    clli
end; (* asm *)
WAITING_TASK_PTR := FIND_AND_REMOVE
(
    ANY_TASK,
    (FROM THE) BLOCKED_QUEUE,
    BECAUSE_THE_BLOCK_WAS_SIGNALLED,
    (AT) addr(C_VAR)
);
    if WAITING_TASK_PTR <> nil then
        INSERT_INTO(READY_QUEUE, WAITING_TASK_PTR);
        REQUEUE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
        asm popfd end;
        end; (* SIGNAL_CONDITION_VARIABLE *)

procedure WAIT_ON_CONDITION_VARIABLE
(*****);
(var C_VAR : CONDITION_VARIABLE);
(-----)

```

```

(* This routine causes the calling task to become blocked. Because *)
(* condition variables have no memory, the calling task is guaranteed to *)
(* become blocked. As soon as another task signals the condition variable *)
(* the waiting task will be made ready (or one of the multiple waiting tasks *)
(* will be made ready). *)
begin (* WAIT_ON_CONDITION_VARIABLE *)
asm
pushf
cli
end; (* asm *)
CHECK_AND_SETUP_BLOCK(NO_TIME_DELAY, addr(C_VAR));
asm popf end;
end; (* WAIT_ON_CONDITION_VARIABLE *)

procedure SEND_MESSAGE TO
(* RECIPIENT_TASK : TASK_IDS; XMIT_MESSAGE_PTR : pointer; *)
(* This routine sends the message (actually the pointer) to the task by *)
(* placing it into the task's mailbox (which is automatically created if *)
(* necessary). Although the calling task cannot block while sending a *)
(* message it is possible for a reschedule to occur. *)
var
(* Pointer to Task Control Block which is waiting on the message *)
(* (may very well be nil). *)
WAITING_TASK_PTR : TASK_CONTROL_BLOCK_PTR;
(* Pointer to mailbox structure which is to receive the message. *)
RECIPIENT_MBOX_PTR : MAILBOX_PTR;
(* Message list elements, used to 'walk' through the *)
(* message list until the end of the list is found (messages must be *)
(* received in FIFO manner). *)
WALKING_PTR : MESSAGE_PTR;
(* Temporary message pointer used to insert the message into the *)
(* message list. *)
TEMP_PTR : MESSAGE_PTR;
begin (* SEND_MESSAGE_TO *)
asm
pushf
cli
end; (* asm *)
RECIPIENT_MBOX_PTR := FIND_OR_CREATE_MBOX_FOR(RECIPIENT_TASK);
WALKING_PTR := RECIPIENT_MBOX_PTR^.CONTENTS;
(* Find the end of the list, allocate storage and create link *)
if WALKING_PTR = nil
then
begin
new(TEMP_PTR);
RECIPIENT_MBOX_PTR^.CONTENTS := TEMP_PTR;
end (* if...then *)
else
begin
while WALKING_PTR^.NEXT <= nil do WALKING_PTR := WALKING_PTR^.NEXT;
new(WALKING_PTR^.NEXT);
TEMP_PTR := WALKING_PTR^.NEXT;
end; (* if...then...else *)
(* Copy the mail into the recipient's mailbox *)
TEMP_PTR^.INFO_ADDRESS := XMIT_MESSAGE_PTR;
(* Tell the recipient that it is there, which enables interrupts *)
SIGNAL_SEMAPHORE(RECIPIENT_MBOX_PTR^.SEM);
asm popf end;
end; (* SEND_MESSAGE_TO *)

procedure RECEIVE_MESSAGE
(* RCV_MESSAGE_PTR : pointer; *)
(* This routine retrieves a message (actually a pointer to a message) from *)
(* the caller's mailbox (which is created if necessary), if there are no *)
(* messages in the mailbox then nil is returned. The task cannot become *)
(* blocked by calling this routine but a reschedule is possible. *)
var
(* Pointer to mailbox structure which has received the message (may *)
(* very well point to empty mailbox). *)
RECIPIENT_MBOX_PTR : MAILBOX_PTR;
(* Temporary pointer used to extract the mail from the message list. *)
TEMP_PTR : MESSAGE_PTR;
begin (* RECEIVE_MESSAGE *)
asm
pushf
cli
end; (* asm *)
(* Setup return value in case the mailbox is empty *)
RCV_MESSAGE_PTR := nil;
RECIPIENT_MBOX_PTR := FIND_OR_CREATE_MBOX_FOR(RUNNING_TASK_PTR^.TASK_ID);
(* Check to see if there is mail to receive *)
if RECIPIENT_MBOX_PTR^.CONTENTS <= nil then
begin
(* Copy the mail into the recipients buffer *)
TEMP_PTR := RECIPIENT_MBOX_PTR^.CONTENTS;
RCV_MESSAGE_PTR := TEMP_PTR^.INFO_ADDRESS;
(* Deallocate the mail storage *)
RECIPIENT_MBOX_PTR^.CONTENTS := TEMP_PTR^.NEXT;
dispose(TEMP_PTR);
end; (* if...then *)
REQUERE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
asm popf end;
end; (* RECEIVE_MESSAGE *)

procedure WAIT_AND_RECEIVE_MESSAGE
(* RCV_MESSAGE_PTR : pointer; *)
(* This routine retrieves a message (actually a pointer to a message) from *)
(* the caller's mailbox (which is created if necessary), if there are no *)
(* messages in the mailbox then nil is returned. The task cannot become *)
(* blocked by calling this routine but a reschedule is possible. *)
var
(* Pointer to mailbox structure which has received the message (may *)
(* very well point to empty mailbox). *)
RECIPIENT_MBOX_PTR : MAILBOX_PTR;
(* Temporary pointer used to extract the mail from the message list. *)
TEMP_PTR : MESSAGE_PTR;
begin (* WAIT_AND_RECEIVE_MESSAGE *)
asm
pushf
cli
end; (* asm *)
(* Setup return value in case the mailbox is empty *)
RCV_MESSAGE_PTR := nil;
RECIPIENT_MBOX_PTR := FIND_OR_CREATE_MBOX_FOR(RUNNING_TASK_PTR^.TASK_ID);
(* Check to see if there is mail to receive *)
if RECIPIENT_MBOX_PTR^.CONTENTS <= nil then
begin
(* Copy the mail into the recipients buffer *)
TEMP_PTR := RECIPIENT_MBOX_PTR^.CONTENTS;
RCV_MESSAGE_PTR := TEMP_PTR^.INFO_ADDRESS;
(* Deallocate the mail storage *)
RECIPIENT_MBOX_PTR^.CONTENTS := TEMP_PTR^.NEXT;
dispose(TEMP_PTR);
end; (* if...then *)
REQUERE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
asm popf end;
end; (* WAIT_AND_RECEIVE_MESSAGE *)

(* ***** *)
TEMP_PTR := MESSAGE_PTR;
begin (* WAIT_AND_RECEIVE_MESSAGE *)
asm
pushf
cli
end; (* asm *)
RECIPIENT_MBOX_PTR := FIND_OR_CREATE_MBOX_FOR(RUNNING_TASK_PTR^.TASK_ID);
(* Check to see if there is mail to receive, which enables interrupts *)
WAIT_ON_SEMAPHORE(RECIPIENT_MBOX_PTR^.SEM);
(* Copy the mail into the recipients buffer *)
TEMP_PTR := RECIPIENT_MBOX_PTR^.CONTENTS;
RCV_MESSAGE_PTR := RECIPIENT_MBOX_PTR^.INFO_ADDRESS;
(* Deallocate the mail storage *)
RECIPIENT_MBOX_PTR^.CONTENTS := TEMP_PTR^.NEXT;
dispose(TEMP_PTR);
asm popf end;
end; (* WAIT_AND_RECEIVE_MESSAGE *)

procedure ENABLE_MOUSE_ACTIONS
(* (MOUSE_ACTIONS : word); *)
(* This routine enables the specified MS-Mouse events to be received by *)
(* the task which is (or will be) waiting to receive mouse actions. *)
begin (* ENABLE_MOUSE_ACTIONS *)
asm
pushf
cli
end; (* asm *)
if not MOUSE_INSTALLED
then
with RUNNING_TASK_PTR do
ERROR_HANDLER(INSUFFICIENT_RESOURCES)(TASK_ID)
else
begin
if USER_MOUSE_EVENT_PTR = nil then
begin
new(USER_MOUSE_EVENT_PTR);
USER_MOUSE_EVENT_PTR := UNSIGNALLED;
end; (* if...then *)
PUSH_MOUSE_EVENT_HANDLER(MOUSE_ACTIONS, MOUSE_EVENT_HANDLER);
ENABLE_MOUSE_DRIVER;
end; (* if...then...else *)
asm popf end;
end; (* ENABLE_MOUSE_ACTIONS *)

procedure WAIT_AND_RECEIVE_MOUSE_ACTIONS
(* (var MOUSE_INFO : MOUSE_PARAMETERS; *)
(* This routine receives MS-Mouse action parameters from the MS-Mouse *)
(* driver. If none of the enabled mouse actions has occurred then the task *)
(* is blocked waiting for one (or more) action to occur. *)
begin (* WAIT_AND_RECEIVE_MOUSE_MESSAGE *)
asm
pushf
cli
end; (* asm *)
if USER_MOUSE_EVENT_PTR = nil
then
begin
with RUNNING_TASK_PTR do
ERROR_HANDLER(INSUFFICIENT_RESOURCES)(TASK_ID);
REQUERE_AND_RESCHEDULE(RUNNING_TASK_PTR^.TASK_ID, READY);
end (* if...then *)
else
begin
WAIT_ON_EVENT(USER_MOUSE_EVENT_PTR);
MOUSE_INFO := MOUSE_INFO_FROM_HANDLER;
end; (* if...then...else *)
asm popf end;
end; (* WAIT_AND_RECEIVE_MOUSE_MESSAGE *)

function WAIT_ON_READKEY : char;
(* This routine provides a means for an application task to perform I/O in *)
(* a 'blocking' manner. This means that while there is no input from the PC *)
(* keyboard the calling task is suspended, as soon as input arrives from the *)
(* keyboard the caller is awakened and given the ASCII code of the key that *)
(* was pressed by the user (i.e., a blocking equivalent of 'readkey'). *)
var
(* Intel 80x86 register variables used to pass register values to *)
(* software interrupt handler. *)
REGS : registers;
begin (* WAIT_ON_READKEY *)
asm
pushf
cli
end; (* asm *)
if USER_KEYBOARD_SEM_PTR = nil then
begin
new(USER_KEYBOARD_SEM_PTR);
USER_KEYBOARD_SEM_PTR := 0;
(* Flush keyboard buffer, throw away all characters in the buffer. *)
REGS.AX := $0C;
REGS.AL := $06;
REGS.DX := $EFFF;
intr($16, REGS);
asm cli end;
end; (* if...then *)
WAIT_ON_SEMAPHORE(USER_KEYBOARD_SEM_PTR);
(* Using this service instead of 'readkey' allows F11 & F12 detection. *)
REGS.AX := $0800;
intr($21, REGS);
asm cli end;
(* ***** *)
(* To ensure that the application will get extended keys the null part *)
(* of the extended key is returned and the controlling semaphore is *)
(* signaled to allow the application to come back and get the extended *)
(* part of the key without waiting. *)
WAIT_ON_READKEY := chr(REGS.AL);
if REGS.AL = ord(NULL_CHR) then
SIGNAL_SEMAPHORE(USER_KEYBOARD_SEM_PTR);
asm popf end;
end; (* WAIT_ON_READKEY *)

(* ***** *)
(* Scheduler Handling Operations *)
(* ***** *)

procedure LOCK_SCHEDULER;
(* This routine prevents all task rescheduling from occurring. Care must *)
(* be taken to ensure that a call to UNLOCK_SCHEDULER is made for each call *)
(* to LOCK_SCHEDULER. In order to maximize system performance, execution *)
(* time within 'locked' regions should be an absolute minimum. *)

```

```

begin (* LOCK_SCHEDULER *)
asm
  pushf
  cll
end; (* asm *)
inc(IN_PREEMPTABLE_REGION);
asm popf end;
end; (* LOCK_SCHEDULER *)

procedure UNLOCK_SCHEDULER;
(* ***** *)
(* This routine allows task rescheduling to occur. Care must be taken to *)
(* ensure that this routine is only called after LOGK_SCHEDULER has already *)
(* been called. *)
(* ***** *)
begin (* UNLOCK_SCHEDULER *)
asm
  pushf
  cll
end; (* asm *)
if IN_PREEMPTABLE_REGION > 0
then
  dec(IN_PREEMPTABLE_REGION)
else
  with RUNNING_TASK_PTR do ERROR_HANDLER[ILLEGAL_OPERRTION](TASK_ID);
asm popf end;
end; (* UNLOCK_SCHEDULER *)

(*****
*) End of user functions
*)
(*****
*) Default Error Handlers
*)
(*$F+ Handlers must be 'far' because they are used as procedural variables *)
(*****

type (*****
*) String used to create time stamps for default error handlers.
*)
STAMP_STRING = string(27);

procedure OPEN_ERROR_LOG;
(*****
*) Creates (or appends) the TASKING error log file in the same directory as *)
(* the program executable file. *)
(*****
var (*****
*) Used to determine the directory of the application executable, *)
(* that is where the error log file will be placed. *)
(*****
DIR : dirstr; NAME namestr; EXT : extstr;

begin (* OPEN_ERROR_LOG *)
(* Open Error Log file (with append if possible) *)
fsplit(paramstr(0), DIR, NAME, EXT);
assign(ERROR_LOG, DIR + 'TASKING.ERR');
(*$I-*) reset(ERROR_LOG); (*$I+*)
if ioresult = 0
then
  append(ERROR_LOG)
else
  rewrite(ERROR_LOG);
end; (* OPEN_ERROR_LOG *)

function TIME_STAMP : STAMP_STRING; near;
(* Returns a time stamp in the form: 'DD/M/YYMM:MM:SSAM'. *)
(*****
var YEAR, MONTH, DAY, DAY_OF_WEEK : word;
    HOUR, MIN, SEC, SECI00TH : word;
    STAMP : STAMP_STRING;
    TEMP : string(6);

begin (* TIME_STAMP *)
(* Date stamp *)
getdate(YEAR, MONTH, DAY, DAY_OF_WEEK);
STAMP := ' ';
STAMP := chr(ord(MONTH div 10) + ord('0'));
STAMP := STAMP + chr(ord(MONTH mod 10) + ord('0')) + '/';
STAMP := STAMP + chr(ord(DAY div 10) + ord('0'));
STAMP := STAMP + chr(ord(DAY mod 10) + ord('0')) + '/';
STAMP := STAMP + chr(ord(YEAR mod 100) div 10) + ord('0');
STAMP := STAMP + chr(ord(YEAR mod 10) + ord('0')) + 'M';

(* Time stamp *)
gettime(HOUR, MIN, SEC, SECI00TH);
str(HOUR mod 12, TEMP);
if TEMP = '0'
then
  TEMP := '12'
else if HOUR < 10 then
  TEMP := TEMP + '0';
STAMP := STAMP + TEMP + ':';
if MIN < 10 then STAMP := STAMP + '0';
str(MIN, TEMP);
STAMP := STAMP + TEMP + ':';
if SEC < 10 then STAMP := STAMP + '0';
str(SEC, TEMP);
STAMP := STAMP + TEMP;
if (HOUR div 12) = 0
then
  TIME_STAMP := STAMP + 'am'
else
  TIME_STAMP := STAMP + 'pm';
end; (* TIME_STAMP *)

procedure TASK_ALREADY_ACTIVE_ERROR_HANDLER(TASK_ID : TASK_IDS);
begin (* TASK_ALREADY_ACTIVE_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Warning! Task is already active (' , TASK_ID, '), request ignored. ');
close(ERROR_LOG);
end; (* TASK_ALREADY_ACTIVE_ERROR_HANDLER *)

procedure INSUFFICIENT_RESOURCES_ERROR_HANDLER(TASK_ID : TASK_IDS);
begin (* INSUFFICIENT_RESOURCES_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Warning! Insufficient resources (' , TASK_ID, '), request ignored. ');
close(ERROR_LOG);
halt($FF);
end; (* INSUFFICIENT_RESOURCES_ERROR_HANDLER *)

procedure TASK_Is_NOT_ACTIVE_ERROR_HANDLER(TASK_ID : TASK_IDS);
begin (* TASK_Is_NOT_ACTIVE_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Warning! Task is not active (' , TASK_ID, '), request ignored. ');
close(ERROR_LOG);
end; (* TASK_Is_NOT_ACTIVE_ERROR_HANDLER *)

procedure TASK_ALREADY_SUSPENDED_ERROR_HANDLER(TASK_ID : TASK_IDS);
begin (* TASK_ALREADY_SUSPENDED_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Warning! Task already suspended (' , TASK_ID, '), request ignored. ');
close(ERROR_LOG);
end; (* TASK_ALREADY_SUSPENDED_ERROR_HANDLER *)

procedure ILLEGAL_TASK_ID_ERROR_HANDLER(TASK_ID : TASK_IDS);
begin (* ILLEGAL_TASK_ID_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Warning! Illegal task identifier (' , TASK_ID, '), request ignored. ');
close(ERROR_LOG);
end; (* ILLEGAL_TASK_ID_ERROR_HANDLER *)

procedure ILLEGAL_OPERRTION_ERROR_HANDLER(TASK_ID : TASK_IDS);
begin (* ILLEGAL_OPERRTION_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Error! Illegal operation on task (' , TASK_ID, '), destroying task (' ,
  RUNNING_TASK_PTR^ .TASK_ID, ')';
DESTROY(RUNNING_TASK_PTR^.TASK_ID);
close(ERROR_LOG);
end; (* ILLEGAL_OPERRTION_ERROR_HANDLER *)

procedure DEADLOCK_DETECTED_ERROR_HANDLER;
begin (* DEADLOCK_DETECTED_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Error! Deadlock detected! Terminating application!');
writeln(TIME_STAMP, ' ',
  'Error! Deadlock detected! Terminating application!');
close(ERROR_LOG);
halt($FF);
end; (* DEADLOCK_DETECTED_ERROR_HANDLER *)

procedure PROGRAM_CANNOT_TERMINATE_ERROR_HANDLER;
begin (* PROGRAM_CANNOT_TERMINATE_ERROR_HANDLER *)
OPEN_ERROR_LOG;
writeln(ERROR_LOG, TIME_STAMP, ' ',
  'Error! Termination delayed, (' , NUMBER_OF_CHILD_PROGRAMS,
  ') child program(s) active. ');
close(ERROR_LOG);
end; (* PROGRAM_CANNOT_TERMINATE_ERROR_HANDLER *)

function TASKING_HEAP_ERROR_HANDLER(SIZE : word)
(*****
*) The run-time systems expects the heap error handler to return status *)
(* Information according to the following rules: *)
(* 0 indicates failure, and causes a run-time error to occur immediately. *)
(* 1 indicates failure, and causes new() or getmem() to return a nil *)
(* pointer. *)
(* 2 indicates success, and causes a retry (which could also cause another *)
(* call to the heap error function. *)
(*****
begin (* TASKING_HEAP_ERROR_HANDLER *)
if SIZE < 0
then
  begin
    OPEN_ERROR_LOG;
    writeln(ERROR_LOG, TIME_STAMP, ' ',
      'Error! Unable to allocate Heap storage for ' , SIZE, ' bytes. ');
    writeln(TIME_STAMP, ' ',
      'Error! Unable to allocate Heap storage for ' , SIZE, ' bytes. ');
    (* Cause run-time error *)
    TASKING_HEAP_ERROR_HANDLER := 0;
    close(ERROR_LOG);
    end (* if...then *)
  else
    (* False alarm *)
    TASKING_HEAP_ERROR_HANDLER := 2;
  end; (* TASKING_HEAP_ERROR_HANDLER *)

(*****
*) End of Default Error Handlers
*)
(*$F-*)
(*****

procedure REPORT_TASKING_STATISTICS; far;
(*****
*) This routine is linked into the exit procedure chain if TASKING *)
(* statistics are to be reported. *)
(*****

const (*****
*) Allows quick and easy conversion between hexadecimal digits and *)
(* the corresponding ASCII character. *)
(*****
HEX_DIGIT : array [0..15] of char = ('0','1','2','3','4','5','6','7',
  '8','9','A','B','C','D','E','F');

(*****
*) IRQ.RPT header for hardware IRQs at application termination. *)
(*****
IRQ_HEADER =
CR + LF +
  ' Vector      IRQ      Interrupts      Period (msec)  +
  ' Frequency (Hz)';

(*****
*) SERVICES.RPT header for software service calls at application *)
(* termination. *)
(*****
SERVICE_HEADER =
CR + LF +
  ' Vector(s)      Service Calls      Period (msec)  +
  ' Frequency (Hz)';

(*****
*) TASKING.RPT header for tasks at application termination. *)
(*****
TASKING_HEADER_1 = ' Tasks at Termination: ';
TASKING_HEADER_2 = ' Task State Task +
  ' Stack Size Stack Used Stack Used CPU Used';
TASKING_HEADER_3 = ' (Delayed Until) ID Priority +
  ' (Words) (Words) (%) (%)';
TASKING_HEADER_4 = ' _____ +
  ' _____ +
  ' _____ +
  ' _____';

type (*****
*) String used to convert long integers (double words) into strings *)
(* with commas embedded, i.e., worst case is -2,147,483,647. *)
(*****
LONGINT_STRING = string(14);

var (*****
*) Loop control variable of task IDs, allows all task IDs to be *)
(* checked for activity so none are missed in task report. *)
(*****
ACTIVE_TASK : TASK_IDS;

(*****
*) Set used to search/access all TASK_IDS. *)
(*****

```

```

(*****
TASK_ID_SET := TASK_ID_SETS;
(*****
(* Pointer to Task Control Blocks which are found in any of the *)
(* TASKING queues. *)
(*****
ACTIVE_TASK_PTR := TASK_CONTROL_BLOCK_PTR;
(*****
(* Total application execution time (as reported by TASKING). *)
(*****
ELAPSED_EXECUTION_TIME := real;
(*****
(* Used in computing ELAPSED TIME. *)
(*****
LCV_YEAR, MONTH, DAY, DAY_OF_WEEK, HOUR, MIN, SEC, SECL00 := word;
(*****
(* Absolute stop time (as reported by TASKING). *)
(*****
STOP_TIME := real;
(*****
(* Total application execution time (as reported by MS-DOS). *)
(*****
ELAPSED_TIME := real;
(*****
(* Amount of the CPU that the application used. *)
(*****
CPU_PERCENT, TOTAL_CPU_PERCENT := real;
(*****
(* Loop control variables used to report on interrupt handlers *)
(* protected by TASKING (both hardware and software). *)
(*****
IRQ, INTERRUPT_NUMBER := integer;
(*****
(* Temporary variables used to detect ranges of interrupt with the *)
(* same report parameters (makes a less cluttered report). *)
(*****
PREVIOUS_INTERRUPT_NUMBER, START, FINISH := integer;
PREVIOUS_SERVICE_CALLS := longint;
(*****
(* Counts the number of unprotected interrupt handlers so that the *)
(* output will not be too near output line. *)
(*****
NUM_UNPROTECTED := integer;
(*****
(* Index into each tasks stack area, used to determine how much of *)
(* the stack was actually used by the task. *)
(*****
MAX_STACK_USED := longint;
(*****
(* Percentage of task stack space, used to determine how much *)
(* stack headroom exists. *)
(*****
STACK_PERCENT, MAX_STACK_PERCENT := real;
(*****
(* Total amount of memory used by the application tasks for stacks *)
(* (only counts tasks which are active at termination). *)
(*****
APPLICATION_STACK_SPACE := longint;
(*****
(* Text filename used for all reports generated by TASKING. *)
(*****
REPORT := text;
(*****
function ADD_COMMAS
(*****
(NUMBER := longint) := LONGINT_STRING;
(*****
(* This procedure simply adds commas into the NUMBER and returns *)
(* the corresponding character string. *)
(*****
var TEMP_STR := LONGINT_STRING;
    CH := integer;
begin (* ADD_COMMAS *)
    str(NUMBER, TEMP_STR);
    CH := length(TEMP_STR) - 3;
    while CH > 0 do
        begin
            insert(',', TEMP_STR, CH + 1);
            CH := CH - 3;
        end; (* while..do *)
    ADD_COMMAS := TEMP_STR;
end; (* ADD_COMMAS *)
begin (* REPORT TASKING STATISTICS *)
(* Restore exit procedure *)
exitproc := SAVE_EXIT;
(* Determine program termination time *)
getdate(YEAR, MONTH, DAY, DAY_OF_WEEK);
gettime(HOUR, MIN, SEC, SECL00);
STOP_TIME := 0.0;
for LCV := 1 to MONTH-1 do
    STOP_TIME := STOP_TIME + DAYS_IN(LCV);
STOP_TIME := ((STOP_TIME + DAY) * 24 + HOUR) * 60 + MIN) * 60 +
    SEC + SECL00 / 100;
ELAPSED_TIME := STOP_TIME - TASKING_STATISTICS.START_TIME;
if ELAPSED_TIME = 0 then ELAPSED_TIME := -1;
(* Determine elapsed execution time *)
ELAPSED_EXECUTION_TIME := MILLISECOND_TICKS / 1000 (Sec);
if ACTUAL_TASKING_CONFIGURATION.STATISTICS in
(HARDWARE_INTERRUPT_STATISTICS, ALL_STATISTICS) then
begin
(* Create Hardware Interrupt Report *)
assign(REPORT, 'IRQ.RPT');
(*SI*) rewrite(REPORT); (* SI*)
writeln(REPORT, 'IRQ HEADER');
for IRQ := 0 to 3F do
    with TASKING_STATISTICS do
        begin
            if IRQ < 8
            then
                write(REPORT, HEX_DIGIT[(IRQ + 308) shr 4] : 5,
                    HEX_DIGIT[(IRQ + 308) and 30F]);
            else
                write(REPORT, HEX_DIGIT[(IRQ + 368) shr 4] : 5,
                    HEX_DIGIT[(IRQ + 368) and 30F]);
            case IRQ of
                30 : write(REPORT, ' Timer 0 ');
                31 : write(REPORT, ' Keyboard ');
                32 : write(REPORT, ' Slave 8259 ');
                33 : write(REPORT, ' COM 1 ');
                34 : write(REPORT, ' COM 2 ');
                36 : write(REPORT, ' Floppy Disk ');
                37 : write(REPORT, ' LPT 1 ');
                38 : write(REPORT, ' Real-Time Clock');
                3D : write(REPORT, ' Co-processor ');
                3E : write(REPORT, ' Fixed Disk ');
                else write(REPORT, ' ', IRQ, ' : 8 - (IRQ div 10));
            end; (* case..of *)
            write(REPORT, 'HARDWARE INTERRUPTS[IRQ] 9);
            if HARDWARE_INTERRUPTS[IRQ] < 0
            then
                writeln(REPORT, (ELAPSED_TIME * 1000 (mSec) /
                    HARDWARE_INTERRUPTS[IRQ]) : 16 : 2,
                    (HARDWARE_INTERRUPTS[IRQ] /
                    ELAPSED_TIME) : 20 : 2);
            else
                writeln(REPORT, 'N/A' 16, 'N/A' : 20);
            end; (* with...do *)
        end; (* if...then *)
    close(REPORT);
end; (* if...then *)
if ACTUAL_TASKING_CONFIGURATION.STATISTICS in
(MSDOS_SERVICES_STATISTICS, ALL_STATISTICS) then
begin
(* Create Software Interrupt Service Report *)
assign(REPORT, 'SERVICES.RPT');
(*SI*) rewrite(REPORT); (* SI*)
NUM_UNPROTECTED := 0;
FINISH := 0;
START := -1;
PREVIOUS_INTERRUPT_NUMBER := -1;
write(REPORT, CR, LF, 'Unprotected Interrupts:', CR, LF, ' ');
for INTERRUPT_NUMBER := 0 to 3FF do
    if not (INTERRUPT_NUMBER in PROTECTED_INTERRUPTS) then
        begin
            if START = -1
            then
                begin
                    START := INTERRUPT_NUMBER,
                    FINISH := INTERRUPT_NUMBER;
                end; (* if...then *)
            else
                if (INTERRUPT_NUMBER = PREVIOUS_INTERRUPT_NUMBER + 1)
                then
                    FINISH := INTERRUPT_NUMBER
                else
                    begin
                        if START = FINISH
                        then
                            begin
                                write(REPORT, ' ', HEX_DIGIT[START shr 4],
                                    HEX_DIGIT[START and 30F]);
                                inc(NUM_UNPROTECTED);
                                if (NUM_UNPROTECTED mod 25) = 0 then
                                    write(REPORT, CR, LF, ' ');
                                end; (* if...then *)
                            end;
                        else
                            begin
                                write(REPORT, ' ', HEX_DIGIT[START shr 4],
                                    HEX_DIGIT[START and 30F],
                                    '...', HEX_DIGIT[FINISH shr 4],
                                    HEX_DIGIT[FINISH and 30F]);
                                inc(NUM_UNPROTECTED, 3);
                                if (NUM_UNPROTECTED mod 25) < 3 then
                                    write(REPORT, CR, LF, ' ');
                                end; (* if...then..else *)
                                START := INTERRUPT_NUMBER;
                                FINISH := INTERRUPT_NUMBER;
                            end; (* if...then..else *)
                        PREVIOUS_INTERRUPT_NUMBER := INTERRUPT_NUMBER;
                    end; (* if...then *)
                    writeln(REPORT, 'SERVICE HEADER');
                    FINISH := 0;
                    START := -1;
                    PREVIOUS_SERVICE_CALLS := -1;
                    PREVIOUS_INTERRUPT_NUMBER := -1;
                    for INTERRUPT_NUMBER := 0 to 3FF do
                        if INTERRUPT_NUMBER in PROTECTED_INTERRUPTS then
                            begin
                                if START = -1 then START := INTERRUPT_NUMBER;
                                if ((TASKING_STATISTICS.SERVICE_CALLS[INTERRUPT_NUMBER] =
                                    PREVIOUS_SERVICE_CALLS and
                                    (INTERRUPT_NUMBER = (PREVIOUS_INTERRUPT_NUMBER + 1))) or
                                    (PREVIOUS_SERVICE_CALLS = -1))
                                then
                                    FINISH := INTERRUPT_NUMBER
                                else
                                    begin
                                        with TASKING_STATISTICS do
                                            begin
                                                if START = FINISH
                                                then
                                                    write(REPORT, HEX_DIGIT[START shr 4] : 6,
                                                        HEX_DIGIT[START and 30F],
                                                        SERVICE_CALLS[INTERRUPT_NUMBER] . 18);
                                                else
                                                    write(REPORT, HEX_DIGIT[START shr 4] : 4,
                                                        HEX_DIGIT[START and 30F], '...',
                                                        HEX_DIGIT[FINISH shr 4],
                                                        HEX_DIGIT[FINISH and 30F],
                                                        SERVICE_CALLS[INTERRUPT_NUMBER] : 16);
                                                    if SERVICE_CALLS[INTERRUPT_NUMBER] < 0
                                                    then
                                                        writeln(REPORT, (ELAPSED_TIME * 1000 (mSec) /
                                                            SERVICE_CALLS[INTERRUPT_NUMBER])
                                                            : 18 : 2,
                                                            (SERVICE_CALLS[INTERRUPT_NUMBER] /
                                                            ELAPSED_TIME) : 20 : 2);
                                                    else
                                                        writeln(REPORT, 'N/A' : 18, 'N/A' : 20);
                                                    end; (* with...do *)
                                                START := INTERRUPT_NUMBER;
                                                FINISH := INTERRUPT_NUMBER;
                                                end; (* if...then..else *)
                                            end;
                                        PREVIOUS_SERVICE_CALLS :=
                                            TASKING_STATISTICS.SERVICE_CALLS[INTERRUPT_NUMBER];
                                        PREVIOUS_INTERRUPT_NUMBER := INTERRUPT_NUMBER;
                                    end; (* if...then *)
                                if START < FINISH
                                then
                                    with TASKING_STATISTICS do
                                        begin
                                            if START = FINISH
                                            then
                                                write(REPORT, HEX_DIGIT[START shr 4] : 4,
                                                    HEX_DIGIT[START and 30F], '...',
                                                    HEX_DIGIT[FINISH shr 4],
                                                    HEX_DIGIT[FINISH and 30F],
                                                    SERVICE_CALLS[INTERRUPT_NUMBER] : 16);
                                                if SERVICE_CALLS[INTERRUPT_NUMBER] < 0
                                                then
                                                    writeln(REPORT, (ELAPSED_TIME * 1000 (mSec) /
                                                        SERVICE_CALLS[INTERRUPT_NUMBER]) : 18
                                                            ,
                                                            (SERVICE_CALLS[INTERRUPT_NUMBER] /
                                                            ELAPSED_TIME) : 20 : 2);
                                                    else
                                                        writeln(REPORT, 'N/A' : 18, 'N/A' : 20);
                                                    end; (* with...do *)
                                                close(REPORT);
                                                end; (* if...then *)
                                            end;
                                        if ACTUAL_TASKING_CONFIGURATION.STATISTICS in

```



```

(TASK_STATISTICS, ALL_STATISTICS) then
with TASKING_STATISTICS do
begin
  (* Create Task Activity Report *)
  APPLICATION_STACK_SPACE := 0;
  MAX_STACK_PERCENT := 0;
  TOTAL_CPU_PERCENT := 0;
  assign(REPORT, 'TASKING.RPT');
  (*SI-*) rewrite(REPORT); (* $!+*)
  if ACTUAL_TASKING_CONFIGURATION.TASKING_MODEL = COOPERATIVE
  then
    write(REPORT, 'Cooperative Multi-tasking')
  else
    write(REPORT, 'Preemptive Multi-tasking');
  writeln(REPORT, 'Statistical Information');
  write(REPORT, ' -Priority Inheritance');
  if ACTUAL_TASKING_CONFIGURATION.PRIORITY_INHERITANCE_ENABLED
  then
    writeln(REPORT, 'Enabled')
  else
    writeln(REPORT, 'Disabled');
  case ACTUAL_TASKING_CONFIGURATION.PRIORITY_SCHEDULING_POLICY of
  STATIC_PRIORITIES : writeln(REPORT, ' -Static Priorities');
  ROTATING_PRIORITIES : writeln(REPORT, ' -Rotating Priorities');
  end; (* case .of *)
  writeln(REPORT, 'Task Activity:');
  writeln(REPORT, 'Context Switches = ', CONTEXT_SWITCHES : 12,
  (-, round(CONTEXT_SWITCHES / ELAPSED_TIME),
  per second));
  writeln(REPORT, 'Cooperative = ',
  COOPERATIVE_CONTEXT_SWITCHES : 12,
  (-, round(COOPERATIVE_CONTEXT_SWITCHES /
  ELAPSED_TIME), ' per second');
  writeln(REPORT, 'Preemptive = ',
  CONTEXT_SWITCHES-COOPERATIVE_CONTEXT_SWITCHES:12,
  (-, round((CONTEXT_SWITCHES -
  COOPERATIVE_CONTEXT_SWITCHES) /
  ELAPSED_TIME), ' per second'));
  if ACTUAL_TASKING_CONFIGURATION.TASKING_MODEL = PREEMPTIVE then
  writeln(REPORT, 'Target Time Slice = ',
  ACTUAL_TASKING_CONFIGURATION.TARGET_TIMESLICE,
  ' usec');
  if CONTEXT_SWITCHES <> 0 then
  writeln(REPORT, 'Achieved Time Slice = -',
  round(ELAPSED_TIME / CONTEXT_SWITCHES * 1e6),
  ' usec');
  if HARDWARE_INTERRUPTS(0) <> 0 then
  writeln(REPORT, 'Available CPU Bandwidth = ',
  (100.0 * NULL_CLOCK_TICKS /
  HARDWARE_INTERRUPTS(0)) : 6 : 3, ' %');
  if NUMBER_OF_PERIODIC_EVENTS <> 0 then
  writeln(REPORT, 'Periodic Event Faults = ',
  (100.0 * NUMBER_OF_PERIODIC_EVENTS_MISSED /
  NUMBER_OF_PERIODIC_EVENTS) : 6 : 3, ' %');
  writeln(REPORT, TASKING_HEADER 1);
  writeln(REPORT, TASKING_HEADER 2);
  writeln(REPORT, TASKING_HEADER 3);
  writeln(REPORT, TASKING_HEADER 4);
  for TASK_ID_SET := low(TASK_ID_SETS) to high(TASK_ID_SETS) do
  for ACTIVE_TASK := (ord(TASK_ID_SET shl 8) + low(byte) to
  (ord(TASK_ID_SET shl 8) + high(byte)) do
  if lo(ACTIVE_TASK) in ALLOCATED_TASK_IDS(TASK_ID_SET) then
  begin
    if RUNNING_TASK_PTR^.TASK_ID = ACTIVE_TASK
    then
      begin
        ACTIVE_TASK_PTR := RUNNING_TASK_PTR;
        write(REPORT, 'Running', ' ' : 8);
        end (* if...then *)
      else
        begin
          ACTIVE_TASK_PTR := FIND_AND_REMOVE
          (
            ACTIVE_TASK,
            (From the) READY_QUEUE,
            UNCONDITIONALLY,
            THERE_IS_NO_BLOCK_TO_MATCH
          );
          if ACTIVE_TASK_PTR = nil
          then
            begin
              ACTIVE_TASK_PTR := FIND_AND_REMOVE
              (
                ACTIVE_TASK,
                (From the) BLOCKED_QUEUE,
                UNCONDITIONALLY,
                THERE_IS_NO_BLOCK_TO_MATCH
              );
              if ACTIVE_TASK_PTR = nil
              then
                begin
                  ACTIVE_TASK_PTR := FIND_AND_REMOVE
                  (
                    ACTIVE_TASK,
                    (From the) DELAYED_QUEUE,
                    UNCONDITIONALLY,
                    THERE_IS_NO_BLOCK_TO_MATCH
                  );
                  write(REPORT, 'Delayed = ',
                  ACTIVE_TASK_PTR^.
                  WAITING_FOR^.ABSOLUTE_TIME : 8);
                  end (* if...then *)
                else
                  write(REPORT, 'Blocked', ' ' : 8);
                  end (* if...then *)
                else
                  write(REPORT, 'Ready', ' ' : 8);
                  end (* if...then...else *)
              with ACTIVE_TASK_PTR^ do
                begin
                  MAX_STACK_USED := 0;
                  while mem[seg(STACK_PTR) : ofs(STACK_PTR) +
                  MAX_STACK_USED] = DEFAULT_STACK_VALUE do
                    inc(MAX_STACK_USED, 2);
                    dec(MAX_STACK_USED, 2);
                  STACK_SIZE := STACK_SIZE div 2;
                  MAX_STACK_USED := STACK_SIZE - (MAX_STACK_USED div 2);
                  if PRIORITY > high(USER_PRIORITIES) then
                    PRIORITY := high(USER_PRIORITIES);
                  if STACK_SIZE = 0
                  then
                    STACK_PERCENT := -1
                  else
                    STACK_PERCENT := MAX_STACK_USED / STACK_SIZE * 100;
                  if STACK_PERCENT > MAX_STACK_PERCENT
                  then MAX_STACK_PERCENT := STACK_PERCENT;
                  CPU_PERCENT := 100 * CLOCK_TICKS /
                  HARDWARE_INTERRUPTS(0);
                  TOTAL_CPU_PERCENT := TOTAL_CPU_PERCENT + CPU_PERCENT;
                  writeln(REPORT, TASK_ID : 4, PRIORITY : 8,
                  STACK_SIZE : 11,
                  MAX_STACK_USED : 11,
                  STACK_PERCENT : 11 : 1,
                  CPU_PERCENT : 10 : 1);
                  inc(APPLICATION_STACK_SPACE, STACK_SIZE * 2);
                  end; (* with...do *)
                end; (* if...then *)
              writeln(REPORT);

```

```

writeln(REPORT, 'Maximum percent of stack used: ',
MAX_STACK_PERCENT : 3 : 1, '%');
writeln(REPORT);
writeln(REPORT, 'Heap Information (bytes):');
writeln(REPORT, 'Total Available: ',
ADD_COMMAS(TOTAL_AVAILABLE_HEAP) : 7);
writeln(REPORT, 'Used for Application Stacks: ',
ADD_COMMAS(APPLICATION_STACK_SPACE) : 7);
writeln(REPORT, 'Available to Application: ',
ADD_COMMAS(TOTAL_AVAILABLE_HEAP -
APPLICATION_STACK_SPACE) : 7);
writeln(REPORT);
writeln(REPORT, 'Execution Time:');
writeln(REPORT, 'Effective = ', ELAPSED_EXECUTION_TIME : 8 : 1,
seconds);
writeln(REPORT, 'Absolute = ', ELAPSED_TIME : 7 : 2,
seconds);
close(REPORT);
end; (* with...do *)
end; (* REPORT_TASKING_STATISTICS *)

procedure TASKING_EXIT; far;
(*-----*)
(* This routine is linked into the exit procedure chain to ensure that all *)
(* interrupt vectors that this unit hooks are unhooked before terminating. *)
(* The exit procedure chain is executed no matter what has happened to cause *)
(* the program to terminate. *)
(* The first time this routine is executed it is because the 'main' of the *)
(* user application terminated, at that time all tasks become alive. Any *)
(* subsequent execution of this routine actually terminates the program. *)
(* If the application 'main' (or any task) terminates abnormally then the *)
(* program will terminate (whether tasks are alive or not). *)
(*-----*)
var
  (*-----*)
  (* Loop control variables used to restore interrupt handlers *)
  (* protected by TASKING (both hardware and software). *)
  (*-----*)
  IRQ, INTERRUPT_NUMBER : integer;

begin (* TASKING_EXIT *)
  (* Application is trying to terminate, wait until it can if necessary... *)
  if NUMBER_OF_CHILD_PROGRAMS <> 0 then
  begin
    PROGRAM_CANNOT_TERMINATE_ERROR_HANDLER;
    while NUMBER_OF_CHILD_PROGRAMS <> 0 do
      begin
        ACTUAL_TASKING_CONFIGURATION.TASKING_MODEL := PREEMPTIVE;
        asm sti end;
        end; (* if...then *)
      end; (* if...then *)

    (* Make sure that no more context switches are performed *)
    LOCK_SCHEDULER;

    (* Unprogram the timer chip *)
    asm cli end;
    port(TIMER_CONTROL_PORT) := SYSTEM_TIMER_CONTROL_WORD;
    port(TIMER_0_DATA_PORT) := lo(MSDOS_TIMER_0_VALUE);
    port(TIMER_0_DATA_PORT) := hi(MSDOS_TIMER_0_VALUE);

    (* Restore revector'd interrupt handler *)
    with OLD_REVECTORED_INTERRUPT do
      begin
        mem[$0000 : REVECTORED_INTERRUPT_NUMBER shl 2] := OFFSET;
        mem[$0000 : (REVECTORED_INTERRUPT_NUMBER shl 2) + 2] := SEGMENT;
        end; (* with...do *)

    (* Restore H/W Interrupt Handler(s) *)
    for IRQ := 0 to 7 do
      with HARDWARE_HANDLER(IRQ) do
        begin
          mem[$0000 : (($08 + IRQ) shl 2)] := OFFSET;
          mem[$0000 : (($08 + IRQ) shl 2) + 2] := SEGMENT;
          end; (* with...do *)
        for IRQ := 8 to 15 do
          with HARDWARE_HANDLER(IRQ) do
            begin
              mem[$0000 : (($08 + IRQ) shl 2)] := OFFSET;
              mem[$0000 : (($08 + IRQ) shl 2) + 2] := SEGMENT;
              end; (* with...do *)
            (* Restore S/W Interrupt Handlers *)
            for INTERRUPT_NUMBER := 0 to $FF do
              if INTERRUPT_NUMBER in PROTECTED_INTERRUPTS then
                with INTERRUPT_VECTOR_HANDLER(INTERRUPT_NUMBER) do
                  begin
                    mem[$0000 : (INTERRUPT_NUMBER shl 2)] := OFFSET;
                    mem[$0000 : (INTERRUPT_NUMBER shl 2) + 2] := SEGMENT;
                    end; (* with...do *)
                  (* Restore exit procedure? *)
                  if ACTUAL_TASKING_CONFIGURATION.STATISTICS <> NO_STATISTICS
                  then
                    exitproc := @REPORT_TASKING_STATISTICS
                  else
                    exitproc := SAVE_EXIT;
                  asm sti end;
                  end; (* TASKING_EXIT *)

procedure TASKING_START; far;
(*-----*)
(* This routine is linked into the exit procedure chain to ensure that all *)
(* interrupt vectors that this unit needs are hooked before initiating multi- *)
(* tasking. The exit procedure chain is executed no matter what has *)
(* happened to cause the program to terminate, but only normal termination *)
(* initiates TASKING. *)
(*-----*)
const
  (*-----*)
  (* Creates a convenient way to install the TASKING hardware *)
  (* interrupt handlers. *)
  (*-----*)
  HARDWARE_INTERCEPTORS : array[2..15] of pointer =
  (
    @IRQ_2_INTERCEPTOR, @IRQ_3_INTERCEPTOR, @IRQ_4_INTERCEPTOR,
    @IRQ_5_INTERCEPTOR, @IRQ_6_INTERCEPTOR, @IRQ_7_INTERCEPTOR,
    @IRQ_8_INTERCEPTOR, @IRQ_9_INTERCEPTOR, @IRQ_A_INTERCEPTOR,
    @IRQ_B_INTERCEPTOR, @IRQ_C_INTERCEPTOR, @IRQ_D_INTERCEPTOR,
    @IRQ_E_INTERCEPTOR, @IRQ_F_INTERCEPTOR
  );

var
  (*-----*)
  (* Loop control variables used to save interrupt handlers protected *)
  (* by TASKING (both hardware and software). *)
  (*-----*)
  IRQ, INTERRUPT_NUMBER : integer;

  (*-----*)
  (* $0x86 registers user to make MS-DOS system calls. *)
  (*-----*)
  REGS : registers;

  (*-----*)
  (* Used to compute new value for the system clock timer. *)
  (*-----*)
  TEMP_TIMER_0_VALUE : longint;
  (*-----*)

```

```

(* Used in computing ELAPSED_TIME. *)
LCV, YEAR, MONTH, DAY, QAY_OF_WEEK, HOUR, MIN, SEC, SECL00 : word;
end; (* if...then *)
end; (* POLICY_VALUE *)

procedure STATISTICS_VALUE
(* (CONTROL : string; var PARAMETER : STATISTIC_OPTIONS); *)
(* This procedure examines the global LINE searching for CONTROL, if *)
(* it is found then it's string value is examined to see if it matches *)
(* one of the valid statistic gathering modes. *)
begin (* STATISTICS VALUE *)
  if UP_STRING(copy(LINE, 1, POSITION - 1)) = CONTROL then
    begin
      delete(LINE, 1, POSITION);
      while LINE[length(LINE)] = ' ' do delete(LINE, length(LINE), 1);
      POSITION := 1;
      if UP_STRING(LINE) = NO_STATISTICS_FLAG
      then
        PARAMETER := NO_STATISTICS
      else if UP_STRING(LINE) = TASK_STATISTICS_FLAG
      then
        PARAMETER := TASK_STATISTICS
      else if UP_STRING(LINE) = MSDOS_SERVICES_STATISTICS_FLAG
      then
        PARAMETER := MSDOS_SERVICES_STATISTICS
      else if UP_STRING(LINE) = HARDWARE_INTERRUPT_STATISTICS_FLAG
      then
        PARAMETER := HARDWARE_INTERRUPT_STATISTICS
      else if UP_STRING(LINE) = ALL_STATISTICS_FLAG
      then
        PARAMETER := ALL_STATISTICS;
      end; (* if...then *)
    end; (* STATISTICS_VALUE *)

procedure MODEL_VALUE(CONTROL : string; var PARAMETER : TASKING_MODELS;
(* This procedure examines the global LINE searching for CONTROL, if *)
(* it is found then it's string value is examined to see if it matches *)
(* one of the valid tasking models. *)
begin (* MODEL_VALUE *)
  if UP_STRING(copy(LINE, 1, POSITION - 1)) = CONTROL then
    begin
      delete(LINE, 1, POSITION);
      while LINE[length(LINE)] = ' ' do delete(LINE, length(LINE), 1);
      POSITION := 1;
      if UP_STRING(LINE) = COOPERATIVE_FLAG
      then
        PARAMETER := COOPERATIVE
      else if UP_STRING(LINE) = PREEMPTIVE_FLAG
      then
        PARAMETER := PREEMPTIVE;
      end; (* if...then *)
    end; (* MODEL_VALUE *)

begin (* PROCESS INI FILE *)
  PROCESS_INI_FILE := false;
  INI_FILENAME := paramstr(0);
  delete(INI_FILENAME, length(INI_FILENAME) - 2, 3);
  INI_FILENAME := INI_FILENAME + 'INI';
  assign(INI_FILE, INI_FILENAME);
  (*SI*) reset(INI_FILE); (*SI*)
  if ioread = 0 then
    begin
      PROCESS_INI_FILE := true;
      while not eof(INI_FILE) do
        begin
          readln(INI_FILE, LINE);
          POSITION := pos(COMMENT_DELIMITER, LINE);
          if POSITION < 0 then
            delete(LINE, POSITION, length(LINE));
          if UP_STRING(copy(LINE, 1, length(TASKING_LABEL))) =
            TASKING_LABEL then
            begin
              PARSING := TASKING_SECTION;
              readln(INI_FILE, LINE);
              POSITION := pos(COMMENT_DELIMITER, LINE);
              if POSITION < 0 then
                delete(LINE, POSITION, length(LINE));
              end; (* if...then *)
              POSITION := pos('=', LINE);
              if (POSITION < 0) and (PARSING = TASKING_SECTION) then
                with CONFIG do
                  begin
                    POLICY_VALUE
                    (
                      PRIORITY_SCHEDULING_POLICY_CONTROL,
                      PRIORITY_SCHEDULING_POLICY
                    );
                    BOOLEAN_VALUE
                    (
                      PRIORITY_INHERITANCE_CONTROL,
                      PRIORITY_INHERITANCE_ENABLED
                    );
                    STATISTICS_VALUE
                    (
                      STATISTICS_CONTROL,
                      STATISTICS
                    );
                    MODEL_VALUE
                    (
                      TASKING_MODEL_CONTROL,
                      TASKING_MODEL
                    );
                    INTEGER_VALUE
                    (
                      TARGET_TIMESLICE_CONTROL,
                      TARGET_TIMESLICE
                    );
                end; (* Tasking Section *)
              end; (* while...do *)
            close(INI_FILE);
          end; (* if...then *)
        end; (* PROCESS_INI_FILE *)
    end; (* TASKING_START *)
    if (exitcode = 0) and (erroraddr = nil)
    then
      begin
        INI_FILE_CONFIGURATION := TASKING_CONFIGURATION;
        if PROCESS_INI_FILE(INI_FILE_CONFIGURATION)
        then
          (* 'Take' INI file configuration parameters *)
          ACTUAL_TASKING_CONFIGURATION := INI_FILE_CONFIGURATION
        else
          (* 'Take' user configuration parameters *)
          ACTUAL_TASKING_CONFIGURATION := TASKING_CONFIGURATION;
        TASKING_CONFIGURATION := ACTUAL_TASKING_CONFIGURATION;
        (* Pass control to 'next' exit procedure *)
        exitproc := @TASKING_EXIT;
        gettime(HOUR, MIN, SEC, SECL00);
        getdate(YEAR, MONTH, DAY, DAY_OF_WEEK);
        with TASKING_STATISTICS do

```

```

begin
    START_TIME := 0.0;
    for LCV := 1 to MONTH-1 do
        START_TIME := START_TIME + DAYS_IN(LCV);
        START_TIME := (((START_TIME + DAY) * 24 +
            HOUR) * 60 + MIN) * 60 + SEC + SECL00 / 100;
    end; (* with...do *)

(* Save revector interrupt handler *)
getintvec(REVECTORED_INTERRUPT_NUMBER,
    OLD_REVECTORED_INTERRUPT_VECTOR);

(* Chain into S/W Interrupt Handlers *)
for INTERRUPT_NUMBER := 0 to 0FF do
    if INTERRUPT_NUMBER in PROTECTED_INTERRUPTS then
        begin
            getintvec(INTERRUPT_NUMBER,
                INTERRUPT_VECTOR_HANDLER{INTERRUPT_NUMBER}.VECTOR);
            setintvec(INTERRUPT_NUMBER, @SW_INTERRUPT_INTERCEPTOR);
        end; (* if...then *)

(* Chain into Keyboard handler *)
getintvec(KEYBOARD_INTERRUPT_NUMBER,
    HARDWARE_HANDLER{1}.VECTOR);
setintvec(KEYBOARD_INTERRUPT_NUMBER, @KEYBOARD_INTERCEPTOR);

(* Chain into IRQ 2 through IRQ 7 handlers *)
for IRQ := 2 to 7 do
    begin
        getintvec($00 + IRQ, HARDWARE_HANDLER{IRQ}.VECTOR);
        setintvec($00 + IRQ, HARDWARE_INTERCEPTORS{IRQ});
    end;
(* Chain into IRQ 8 through IRQ 15 handlers *)
for IRQ := 8 to 15 do
    begin
        getintvec($68 + IRQ, HARDWARE_HANDLER{IRQ}.VECTOR);
        setintvec($68 + IRQ, HARDWARE_INTERCEPTORS{IRQ});
    end;

(* Enable re-scheduling when creating tasks *)
ALLOW_RESCHEDULE_IN_CREATE := true;

(* Chain into Clock handler *)
asm cll end;
with HARDWARE_HANDLER{0} do
    begin
        OFFSET := memm($0000: SYSTEM_TIMER_INTERRUPT_NUMBER shl 2);
        SEGMENT := memm($0000: {SYSTEM_TIMER_INTERRUPT_NUMBER shl 2}+2);
    end; (* with...do *)
    CLOCK_VECTOR := HARDWARE_HANDLER{0}.VECTOR;
    memm($0000 : SYSTEM_TIMER_INTERRUPT_NUMBER shl 2) :=
        ofs(CLOCK_INTERCEPTOR);
    memm($0000 : {SYSTEM_TIMER_INTERRUPT_NUMBER shl 2} + 2) :=
        seg(CLOCK_INTERCEPTOR);

(* Enforce target timeslice range restrictions (if necessary) *)
with ACTUAL_TASKING_CONFIGURATION do
    if TARGET_TIMESLICE < MINIMUM_TIMESLICE
    then
        TARGET_TIMESLICE := MINIMUM_TIMESLICE
    else
        if TARGET_TIMESLICE > MAXIMUM_TIMESLICE then
            TARGET_TIMESLICE := MAXIMUM_TIMESLICE;

(* Compute new timer counter value *)
TEMP_TIMER_0_VALUE := round(TIMER_CLOCK_FREQUENCY / 1e6 *
    ACTUAL_TASKING_CONFIGURATION.TARGET_TIMESLICE);

(* Verify that the timer will be at least as fast as TASKING needs *)
if TEMP_TIMER_0_VALUE > ONE_MSEC_TIMER_0_VALUE
then
    NEW_TIMER_0_VALUE := ONE_MSEC_TIMER_0_VALUE
else
    NEW_TIMER_0_VALUE := TEMP_TIMER_0_VALUE;

(* Establish system clock scaling factors *)
TIMER_PERIOD := NEW_TIMER_0_VALUE /
    TIMER_CLOCK_FREQUENCY {Sec} * 1e6;

(* Reprogram the timer chip *)
port{TIMER_CONTROL_PORT} := SYSTEM_TIMER_CONTROL_WORD;
port{TIMER_0_DATA_PORT} := lo(NEW_TIMER_0_VALUE);
port{TIMER_0_DATA_PORT} := hi(NEW_TIMER_0_VALUE);

(* Talk about a 'goto'... *)
asm jmp DISPATCH_TASK end;
end (* if...then *)
else
    exitproc := SAVE_EXIT;
end; (* TASKING_START *)

{*****}
(* This is the TASKING unit initialization, there is not much to do except *)
(* setup the system stack and install the exit procedure and interrupt *)
(* handlers. *****)
{*****}
begin (* TASKING_INITIALIZATION *)
    TASKING_STATISTICS.TOTAL_AVAILABLE_HEAP := memavail;

(* Clear the null task stack *)
fillchar(NULL_STACK_PTR, NULL_STACK_SIZE, byte(DEFAULT_STACK_VALUE));

(* Initialize TASK_ID array *)
for TASK_ID_SET := low{TASK_ID_SETS} to high{TASK_ID_SETS} do
    ALLOCATED_TASK_IDS{TASK_ID_SET} := {};

(* Install exit procedure *)
SAVE_EXIT := exitproc;
exitproc := @TASKING_START;

heaperror := @TASKING_HEAP_ERROR_HANDLER;
end. (* TASKING_INITIALIZATION *)

```

# 9. Appendix B Genetic Algorithm Source Code

## 9.1 RTOS-APP.PAS

```

program REAL_TIME_APPLICATION;
.....
[* This program interprets an INI file and a CSV (Comma separated text)
[* file to configure itself to implement a multitasking design. After the
[* design has been executed a report is generated which indicates how well
[* the design time restraints were met.
.....
Compiler Options (Ver. 7.0)
[*SA+ Word Alignment
[*SB- Debug Code Generation ON (Sort of)
[*SD- Requires /V option to TPC to activate
[*SL+ Local Debug symbols ON (Sort of)
[*SF- Requires /V option to TPC to activate
[*SG- Far calls only as needed always
[*SI- Generic 80x86 code only
[*SM $8000,50000,100000 I/O Checking OFF
[*SN- Memory (Stack, Minheap, Maxheap)
[*SO- Software Emulation of 80x87
[*SP- No 80x87 run-time emulation
[*SQ- Overlays NOT allowed
[*SR- Overflow Checking OFF
[*SS- Range Checking OFF
[*ST+ Stack Checking OFF
[*SV- Force Typed '?' references
[*SX+ Var-string Checking OFF
[*SX+ Enable Extended syntax
.....
uses TASKING;

const
.....
[* Limit on the number of tasks allowed, basically the internal data
[* structure must fit within 64K, this limitation restricts the number
[* of tasks.
.....
MAX_TASKS = 50;

.....
[* The application will execute for a length of time equal to the
[* longest task period times this factor. This will ensure that the
[* results gathered are statistically significant.
.....
EXECUTION_DWELL_TIME_FACTOR = 100;

.....
[* The absolute execution of the application is limited so that the
[* results can be gathered in a reasonable amount of time.
.....
MAXIMUM_DWELL_TIME = 5 * 60; [* Seconds

type
.....
[* Intertask communication is accomplished via TASKING resources (e.g.
[* events, semaphores, messages, condition variables), the resources
[* can either be received or signaled by the task.
.....
RESOURCE_ACTIONS = (SIGNAL_RESOURCE, RECEIVE_RESOURCE);

.....
[* The TASKING types are used for all intertask communications, this
[* typemark defines those communication constructs.
.....
RESOURCE_TYPES =
MESSAGE_RESOURCE, SEMAPHORE_RESOURCE,
BINARY_SEMAPHORE_RESOURCE, EVENT_RESOURCE,
CONDITION_VARIABLE_RESOURCE, GENERIC_RESOURCE
;

.....
[* The tasks are allowed to communicate with each other, this data
[* structure defines the manner in which that communication can take
[* place.
.....
COMMUNICATION_SPECIFICATION = record
.....
[* Not all tasks use all for any/ of the intertask communication
[* resources, this element is used to determine which resources are
[* active for each task.
.....
DEFINED : boolean;
(* The user defines the resource by a text string name, this is used
[* to determine when the same resource is being used by more than a
[* single task.
.....
RESOURCE_NAME : string(25);
.....
[* If a message is being used as the communication mechanism then
[* the task ID of the task that is going to receive the message is
[* needed, this element holds this information (if applicable).
.....
RESOURCE_TASK_ID : TASK_IDS;
.....
[* The most important specification parameter for the resource is
[* the absolute time from the beginning of the period that the
[* intertask communication is to take place.
.....
RESOURCE_TIME : word;
.....
[* The actual intertask communication resource is dynamically
[* allocated by the task itself.
.....
case RESOURCE_TYPE : RESOURCE_TYPES of
MESSAGE_RESOURCE : (MSG_PTR : ^longint);
SEMAPHORE_RESOURCE : (SEM_PTR : ^SEMAPHORE);
BINARY_SEMAPHORE_RESOURCE : (B_SEM_PTR : ^BINARY_SEMAPHORE);
EVENT_RESOURCE : (EVENT_PTR : ^EVENT);
CONDITION_VARIABLE_RESOURCE : (C_VAR_PTR : ^CONDITION_VARIABLE);
GENERIC_RESOURCE : (RESOURCE_PTR : pointer);
end; [* COMMUNICATION_SPECIFICATION

end;

type
.....
[* Depending on the timing requirements of the task, the task is
[* classified into one of the following categories.
[* Note: All of the following timing descriptions assume that the
[* workload time is less than (or equal to) the task period.
[* The diagrams are 'not' to scale, only relative timing
[* information is intended to be conveyed.
.....
TASK_CATEGORIES = (
.....
[* Case #1 - Proactive task - deadline occurs during the task
[* execution, i.e., the task has a critical point in its

```

```

.....
execution that it must get to before the 'deadline'
amount of time has elapsed.
.....
Deadline
Workload
Period
Time
.....
PROACTIVE;
.....
[* Case #2 - Reactive task - deadline occurs after the task has
[* finished its execution, i.e. the critical event is
[* outside the realm of this task. It does not actually
[* perform the critical operation but the timing of that
[* operation is critical to the execution (actually next
[* execution) of this task.
.....
Deadline
Workload
Period
Time
.....
REACTIVE;
.....
[* Case #3 - Passive task - no deadline. The task merely performs
[* some work but does not have a critical deadline to
[* meet (other than the implicit deadline of periodicity).
.....
Deadline
Workload
Period
Time
.....
PASSIVE;
.....
[* Case #4 - Background task - no deadline, no period. The task
[* performs continuous work, since it never completes a
[* cycle it 'cannot' be late. A task such as this will
[* consume resources and force the application to execute
[* at 100% CPU utilization.
.....
Deadline
Workload
Period
Time
.....
BACKGROUND;

.....
[* All of the information about the task is maintained in this data
[* structure. The task design and intertask constructs are maintained
[* here as well as all performance monitoring parameters.
.....
TASK_SPECIFICATION = record
.....
[* Used to allow messages to be sent between the tasks.
.....
TASK_ID : TASK_IDS;
.....
[* Used to allow the task priorities to be reported at termination.
.....
TASK_PRIORITY : USER_PRIORITIES;
.....
[* Used to determine which task ID the user wants the message sent
[* to (if messages are being used). Also used when the report file
[* is generated to make it more user friendly (i.e. the report is
[* made in terms of the user's task names).
.....
TASK_NAME : string(25);
.....
[* Based on the characteristics of the task it is categorized so
[* that the correct statistics are computed for the task. Not all
[* statistics are valid for all task characteristics.
.....
TASK_CATEGORY : TASK_CATEGORIES;
.....
[* Amount of time between successive executions of the task.
.....
PERIOD : word;
.....
[* Amount of time from the beginning of the tasks execution to its
[* 'critical time'.
.....
DEADLINE : word;
.....
[* How 'important' the deadline is: 10 = Hard Deadline (meeting the
[* the deadline is critical), 1 = Soft (meeting the deadline would
[* be 'nice'), linear scale for all other values.
.....
DEADLINE_HARDNESS : 1..10;
.....
[* Amount of time that the task is busy working (i.e., the amount
[* of time that the task executes).
.....
WORKLOAD : word;
.....
[* Each task is allowed to have up to two intertask communication
[* resources (for obvious reasons, each must include both types of
[* communication actions).
.....
INTERTASK_COMMUNICATION :
array (RESOURCE_ACTIONS, 1..2) of COMMUNICATION_SPECIFICATION;
.....
[* As the various timing requirements of the tasks are monitored,
[* the application records how well the requirements are met using
[* the following data structure.
.....
PERIODIC_TIME_MONITOR, DEADLINE_TIME_MONITOR : record
.....
[* Number of times that the event occurred at the correct time.
.....
SUCCESSFUL : longint;
.....
[* Number of times that the event occurred.
.....
TOTAL : longint;
end; [* TIME_MONITOR

end; [* TASK_SPECIFICATION

var
.....
[* Used for creating the supervisor task.
.....
SUPERVISOR : TASK_IDS;
TASK_ATTR : TASK_ATTRIBUTES;

.....
[* The application report file is generated if the program command
[* line is blank. If anything is on the command line then 'no' report
[* report file is generated. Whether the file is generated or not the
[* result is returned via the MS-DOS errorlevel.
.....
GENERATE_REPORT_FILE : boolean;

```

```

(* The main database for the application, all tasks are completely *)
(* specified by the information in this array. *)
TASK : array[1..MAX_TASKS] of TASK_SPECIFICATION;

(* The number of tasks that the user has specified in the .CSV file. *)
NUMBER_OF_TASKS : integer;

(* The largest task period is used to determine the execution time of *)
(* the application, this variable hold that value. *)
MAXIMUM_MSEC_PERIOD : integer;

(* Name of this program .EXE file, used to determine the filenames of *)
(* the .RPT, .IMI and .CSV files. *)
PROGRAM_NAME : string;

procedure SUPERVISOR_TASK(TASK_ID : TASK_IDS; PRIORITY : USER_PRIORITIES); far;
(* This task is "not" specified by the user (i.e. .CSV file). Instead this *)
(* task is needed to bring the application to an orderly shutdown and to *)
(* produce the .RPT file. In addition, this task sets the MS-DOS errorlevel *)
(* to the total application timeliness. *)
var
(* The total time that the application is going to be allowed to *)
(* execute. *)
EXECUTION_TIME : TIME;
(* The number of seconds that the application is going to execute. *)
DWEELL_SECONDS : integer;
(* The file handle for the application report. *)
REPORT : text;
(* Loop Control Variable for accessing all tasks in the task *)
(* specification database. *)
TASK_LCV : integer;
(* The 'timeliness' of the application is computed for each of the *)
(* following areas of execution. *)
PERIODICITY_RESULT, DEADLINE_RESULT, APPLICATION_RESULT : record
TIMELINESS : real;
NUMBER_OF_TIMES : integer;
end; (* TIMELINESS *)

begin (* SUPERVISOR TASK *)
(* Calculate application dwell time *)
DWEELL_SECONDS := round(MAXIMUM_MSEC_PERIOD * EXECUTION_DWEELL_TIME_FACTOR /
1000);
if DWEELL_SECONDS > MAXIMUM_DWEELL_TIME then halt(1);

(* Calculate a time delay to run the application for the dwell time *)
with EXECUTION_TIME do
begin
DAYS := 0;
HOURS := 0;
MINUTES := DWEELL_SECONDS div 60;
DWEELL_SECONDS := DWEELL_SECONDS mod 60;
SECONDS := DWEELL_SECONDS;
MILLISECONDS := 0;
end; (* with...do *)

(* Wait for dwell time to elapse *)
WAIT_FOR_DELAY EXECUTION_TIME;

(* Destroy all application tasks (makes results more accurate)... *)
for TASK_LCV := 1 to NUMBER_OF_TASKS do
DESTROY(TASK_LCV, TASK_ID);

with PERIODICITY_RESULT do
begin
TIMELINESS := 0.0;
NUMBER_OF_TIMES := 0;
end; (* with...do *)
with DEADLINE_RESULT do
begin
TIMELINESS := 0.0;
NUMBER_OF_TIMES := 0;
end; (* with...do *)
with APPLICATION_RESULT do
begin
TIMELINESS := 0.0;
NUMBER_OF_TIMES := 0;
end; (* with...do *)
for TASK_LCV := 1 to NUMBER_OF_TASKS do
with TASK(TASK_LCV) do
begin
with PERIODIC_TIME_MONITOR do
begin
if TOTAL <> 0 then
PERIODICITY_RESULT.TIMELINESS := PERIODICITY_RESULT.TIMELINESS
+ SUCCESSFUL / TOTAL * 100;
inc(PERIODICITY_RESULT.NUMBER_OF_TIMES);
end; (* with...do *)
case TASK_CATEGORY of
PROACTIVE,
REACTIVE :
begin
with DEADLINE_TIME_MONITOR do
begin
if TOTAL <> 0 then
DEADLINE_RESULT.TIMELINESS := DEADLINE_RESULT.TIMELINESS
+ SUCCESSFUL / TOTAL * 100 *
(1 - DEADLINE_HARDNESS / 10);
inc(DEADLINE_RESULT.NUMBER_OF_TIMES);
end; (* with...do *)
end; (* Proactive ) Reactive Tasks *)
PASSIVE : begin (* Passive Task *) end;
BACKGROUND : begin (* Background Task *) end;
end; (* case...of *)
end; (* with...do *)
if PERIODICITY_RESULT.NUMBER_OF_TIMES <> 0 then
with APPLICATION_RESULT do
begin
TIMELINESS := PERIODICITY_RESULT.TIMELINESS + TIMELINESS;
NUMBER_OF_TIMES := PERIODICITY_RESULT.NUMBER_OF_TIMES +
NUMBER_OF_TIMES;
end; (* with...do *)
if DEADLINE_RESULT.NUMBER_OF_TIMES <> 0 then

```

```

with APPLICATION_RESULT do
begin
TIMELINESS := DEADLINE_RESULT.TIMELINESS + TIMELINESS;
NUMBER_OF_TIMES := DEADLINE_RESULT.NUMBER_OF_TIMES +
NUMBER_OF_TIMES;
end; (* with...do *)
with APPLICATION_RESULT do
if NUMBER_OF_TIMES <> 0 then
TIMELINESS := TIMELINESS / NUMBER_OF_TIMES;

if GENERATE_REPORT_FILE then
begin
(* Create Report File *)
assign(REPORT, PROGRAM_NAME + '.RPT');
(*SI*) rewrite(REPORT); (*SI*)
if ioresult <> 0 then halt(1);

writeln(REPORT, ' Individual Task Timeliness');
for TASK_LCV := 1 to NUMBER_OF_TASKS do
with TASK(TASK_LCV) do
begin
writeln(REPORT, ' Task Name: ', TASK_NAME,
', Priority = ', TASK_PRIORITY);
with PERIODIC_TIME_MONITOR do
begin
write(REPORT, ' Periodicity: (', PERIOD, ' msec)',
', SUCCESSFUL / TOTAL);
if TOTAL <> 0 then
writeln(REPORT, ' (', SUCCESSFUL / TOTAL *
100 : 2 : 2, '%)');
end; (* with...do *)
case TASK_CATEGORY of
PROACTIVE,
REACTIVE :
with DEADLINE_TIME_MONITOR do
begin
write(REPORT, ' Deadline: (', DEADLINE, ' msec)',
', SUCCESSFUL / TOTAL);
if TOTAL <> 0 then
writeln(REPORT, ' (', SUCCESSFUL / TOTAL *
100 : 2 : 2, '%)');
end; (* with...do *)
PASSIVE : begin (* Passive Task *) end;
BACKGROUND : begin (* Background Task *) end;
end; (* case...of *)
writeln(REPORT);
end; (* with...do *)
writeln(REPORT);
writeln(REPORT, ' Average Task Timeliness');
with PERIODICITY_RESULT do
if NUMBER_OF_TIMES <> 0 then
writeln(REPORT, ' Periodicity: ', TIMELINESS /
NUMBER_OF_TIMES : 2 : 2,
'%');
with DEADLINE_RESULT do
if NUMBER_OF_TIMES <> 0 then
writeln(REPORT, ' Deadline: ', TIMELINESS /
NUMBER_OF_TIMES : 2 : 2, '%');
writeln(REPORT);
writeln(REPORT, ' Application Timeliness = ',
APPLICATION_RESULT.TIMELINESS : 2 : 2, '%');
close(REPORT);
end; (* if...then *)

(* shut down the application in order to provide exit code... *)
halt(round(APPLICATION_RESULT.TIMELINESS));
end; (* SUPERVISOR TASK *)

procedure GENERAL_TASK(TASK_ID : TASK_IDS; PRIORITY : USER_PRIORITIES); far;
(* This task is used to emulate all user application tasks. It does this *)
(* by assuming the characteristics of the user task specification. *)
var
(* These are the absolute delays, they are used in evaluating how *)
(* well the task performed. *)
PERIODIC_DELAY, DEADLINE_DELAY : TIME;
(* The periodic execution of the task is handled by TASKING, an *)
(* event is used to accomplish this. *)
PERIODIC_EVENT : EVENT;
(* Stop watch/ variables for evaluating performance. *)
START_TIME, STOP_TIME : longint;
(* Database of the tasks critical events (for "each" periodic *)
(* execution. *)
CRITICAL : array[1..6] of record
(* The relative delay for the specific event (from the previous *)
(* event). *)
DELAY : TIME;
(* The actual critical event classification. *)
MILESTONE : (
NOT_APPLICABLE,
DEADLINE_EXPIRES,
WORKLOAD_EXPIRES,
WAIT_FOR_RESOURCE_1,
SIGNAL_RESOURCE_1,
WAIT_FOR_RESOURCE_2,
SIGNAL_RESOURCE_2);
end; (* CRITICAL *)
(* Loop Control Variables used to create the critical event list. *)
TMC_LCV, INSERT : integer;

procedure FIND_OR_CREATE_RESOURCE(ACTION : RESOURCE_ACTIONS;
NUMBER : integer);
(* This routine searches the tasks' resources for the resource specified. *)
(* If it is found, the tasks are made to point to the same resource, if it *)
(* is not found, a new resource is created. *)
var
(* Indicates the existence of this critical event within the events *)
(* of the application (recall that resources can be shared by more *)
(* than one task). *)
FOUND : boolean;
(* Loop Control Variable to search all existing task specifications. *)
TASK_LCV : integer;

```

```

begin (* FIND_OR_CREATE_RESOURCE *)
  FOUND := false;
  (* Check for existing resource *)
  for TASK_LCV := 1 to MAX_TASKS do
    begin
      if not ((ACTION = RECEIVE_RESOURCE) and (NUMBER = 1)) then
        with TASK[TASK_LCV].INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 1) do
          if (RESOURCE_PTR <> nil)
            and
              (RESOURCE_NAME = TASK[TASK_ID].
                INTERTASK_COMMUNICATION(ACTION, NUMBER).
                RESOURCE_NAME) then
            begin
              TASK[TASK_ID].
                INTERTASK_COMMUNICATION(ACTION, NUMBER).RESOURCE_PTR :=
                RESOURCE_PTR;
              FOUND := true;
            end; (* if...then *)
          if not ((ACTION = SIGNAL_RESOURCE) and (NUMBER = 1)) then
            with TASK[TASK_LCV].INTERTASK_COMMUNICATION[SIGNAL_RESOURCE, 1] do
              if (RESOURCE_PTR <> nil)
                and
                  (RESOURCE_NAME = TASK[TASK_ID].
                    INTERTASK_COMMUNICATION(ACTION, NUMBER).
                    RESOURCE_NAME) then
                begin
                  TASK[TASK_ID].
                    INTERTASK_COMMUNICATION[ACTION, NUMBER].RESOURCE_PTR :=
                    RESOURCE_PTR;
                  FOUND := true;
                end; (* if...then *)
              if not ((ACTION = RECEIVE_RESOURCE) and (NUMBER = 2)) then
                with TASK[TASK_LCV].INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 2) do
                  if (RESOURCE_PTR <> nil)
                    and
                      (RESOURCE_NAME = TASK[TASK_ID].
                        INTERTASK_COMMUNICATION[ACTION, NUMBER].
                        RESOURCE_NAME) then
                    begin
                      TASK[TASK_ID].
                        INTERTASK_COMMUNICATION[ACTION, NUMBER].RESOURCE_PTR :=
                        RESOURCE_PTR;
                      FOUND := true;
                    end; (* if...then *)
                  if not ((ACTION = SIGNAL_RESOURCE) and (NUMBER = 2)) then
                    with TASK[TASK_LCV].INTERTASK_COMMUNICATION[SIGNAL_RESOURCE, 2] do
                      if (RESOURCE_PTR <> nil)
                        and
                          (RESOURCE_NAME = TASK[TASK_ID].
                            INTERTASK_COMMUNICATION[ACTION, NUMBER].
                            RESOURCE_NAME) then
                            begin
                              TASK[TASK_ID].
                                INTERTASK_COMMUNICATION(ACTION, NUMBER).RESOURCE_PTR :=
                                RESOURCE_PTR;
                              FOUND := true;
                            end; (* if...then *)
                        end; (* for...to...do *)
                    end;
                if not FOUND then
                  with TASK[TASK_ID].INTERTASK_COMMUNICATION[ACTION, NUMBER] do
                    case RESOURCE_TYPE of
                      MESSAGE_RESOURCE : getmem(pointer(MSG_PTR), sizeof(longint));
                      SEMAPHORE_RESOURCE :
                        begin
                          getmem(pointer(SEM_PTR), sizeof(SEMAPHORE));
                          SEM_PTR := 1;
                        end; (* Semaphore *)
                      BINARY_SEMAPHORE_RESOURCE
                        begin
                          getmem(pointer(B_SEM_PTR), sizeof(BINARY_SEMAPHORE));
                          B_SEM_PTR := 1;
                        end; (* Binary Semaphore *)
                      EVENT_RESOURCE :
                        begin
                          getmem(pointer(EVENT_PTR), sizeof(EVENT));
                          EVENT_PTR := SIGNALLED;
                        end; (* Event *)
                      CONDITION_VARIABLE_RESOURCE
                        getmem(pointer(C_VAR_PTR), sizeof(CONDITION_VARIABLE));
                    end; (* case...of *)
                  end; (* FIND_OR_CREATE_RESOURCE *)
                end;
            begin (* GENERAL_TASK *)
              with TASK[TASK_ID] do
                begin
                  (* Initialize time sequences *)
                  for TIME_LCV := 1 to 6 do
                    with CRITICAL[TIME_LCV] do
                      begin
                        with DELAY do
                          begin
                            DAYS := 0;
                            HOURS := 0;
                            MINUTES := 0;
                            SECONDS := 0;
                            MILLISECONDS := 0;
                            end; (* with...do *)
                            MILESTONE := NOT_APPLICABLE;
                            end; (* with...do *)
                        (* Determine critical time sequence (absolute times only) *)
                        if (DEADLINE <= WORKLOAD) and (WORKLOAD <= PERIOD) then
                          begin
                            TASK_CATEGORY := PROACTIVE;
                            with CRITICAL[1] do
                              begin
                                DELAY.MILLISECONDS := DEADLINE;
                                MILESTONE := DEADLINE_EXPIRES;
                              end; (* with...do *)
                            with CRITICAL[2] do
                              begin
                                DELAY.MILLISECONDS := WORKLOAD;
                                MILESTONE := WORKLOAD_EXPIRES;
                              end; (* with...do *)
                            end; (* if...then *)
                          else if (WORKLOAD < DEADLINE) and (DEADLINE <= PERIOD) then
                            begin
                              TASK_CATEGORY := REACTIVE;
                              with CRITICAL[1] do
                                begin
                                  DELAY.MILLISECONDS := WORKLOAD;
                                  MILESTONE := WORKLOAD_EXPIRES;
                                end; (* with...do *)
                                with CRITICAL[2] do
                                  begin
                                    DELAY.MILLISECONDS := DEADLINE;
                                    MILESTONE := DEADLINE_EXPIRES;
                                  end; (* with...do *)
                                end; (* if...then *)
                              end;
                            else if (WORKLOAD < PERIOD) and (PERIOD <= DEADLINE) then
                              begin
                                TASK_CATEGORY := PASSIVE;
                                with CRITICAL[1] do
                                  begin
                                    DELAY.MILLISECONDS := WORKLOAD;
                                    MILESTONE := WORKLOAD_EXPIRES;
                                  end; (* with...do *)
                                end;
                              end;
                            end;
                          end;
                        (* Setup absolute deadline timer *)
                        with DEADLINE_DELAY do
                          begin
                            DAYS := 0;
                            HOURS := 0;
                            MINUTES := 0;
                            SECONDS := 0;
                            MILLISECONDS := DEADLINE;
                            end; (* with...do *)
                        (* Setup periodic event *)
                        PERIODIC_EVENT := UNSIGNALLED;
                        with PERIODIC_DELAY do
                          begin
                            DAYS := 0;
                            HOURS := 0;
                            MINUTES := 0;
                            SECONDS := 0;
                            MILLISECONDS := PERIOD;
                            end; (* with...do *)
                        (* Setup relative delays *)
                        for TIME_LCV := 6 downto 2 do
                          CRITICAL[TIME_LCV].DELAY.MILLISECONDS :=
                            CRITICAL[TIME_LCV].DELAY.MILLISECONDS -
                            CRITICAL[TIME_LCV-1].DELAY.MILLISECONDS;
                        (* Start emulating the task behavior *)
                        START_PERIODIC_EVENT(PERIODIC_EVENT, PERIODIC_DELAY);
                        repeat
                          (* Perform and evaluate delays (as appropriate) *)
                          START_TIME := GET_MILLISECOND_TICKS;
                          for TIME_LCV := 1 to 6 do
                            with CRITICAL[TIME_LCV] do
                              begin
                                case MILESTONE of
                                  NOT_APPLICABLE : (* Do nothing *);
                                  WORKLOAD_EXPIRES : PREEMPTABLE_DELAY(DELAY);
                                  DEADLINE_EXPIRES :
                                    begin
                                      PREEMPTABLE_DELAY(DELAY);
                                      STOP_TIME := GET_MILLISECOND_TICKS;
                                      if (STOP_TIME - START_TIME) <=
                                        DEADLINE_DELAY.MILLISECONDS then
                                        inc(DEADLINE_TIME_MONITOR.SUCCESSFUL);
                                        inc(DEADLINE_TIME_MONITOR.TOTAL);
                                      end; (* if...then *)
                                    WAIT_FOR_RESOURCE_1 :
                                      with TASK[TASK_ID].
                                        INTERTASK_COMMUNICATION[RECEIVE_RESOURCE, 1] do
                                          case RESOURCE_TYPE of
                                            MESSAGE_RESOURCE :
                                              WAIT_AND_RECEIVE_MESSAGE(pointer(MSG_PTR));
                                            SEMAPHORE_RESOURCE :
                                              WAIT_ON_SEMAPHORE(SEM_PTR);
                                            BINARY_SEMAPHORE_RESOURCE :

```

```

        WAIT ON BINARY_SEMAPHORE(B_SEM_PTR);
EVENT RESOURCE :
    WAIT ON EVENT(EVENT_PTR);
CONDITION VARIABLE RESOURCE :
    WAIT ON CONDITION_VARIABLE(C_VAR_PTR);
end; (* case...of *)
SIGNAL_RESOURCE_1 :
with TASK(TASK_ID)
    INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1) do
case RESOURCE_TYPE of
MESSAGE_RESOURCE :
    SEND MESSAGE TO(RESOURCE_TASK_ID, pointer(MSG_PTR));
SEMAPHORE_RESOURCE :
    SIGNAL_SEMAPHORE(SEM_PTR);
BINARY_SEMAPHORE_RESOURCE :
    SIGNAL_BINARY_SEMAPHORE(B_SEM_PTR);
EVENT_RESOURCE :
    SIGNAL_EVENT(EVENT_PTR);
CONDITION_VARIABLE_RESOURCE :
    SIGNAL_CONDITION_VARIABLE(C_VAR_PTR);
end; (* case...of *)
WAIT FOR RESOURCE_2 :
with TASK(TASK_ID)
    INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 2) do
case RESOURCE_TYPE of
MESSAGE_RESOURCE :
    WAIT AND RECEIVE MESSAGE(pointer(MSG_PTR));
SEMAPHORE_RESOURCE :
    WAIT ON SEMAPHORE(SEM_PTR);
BINARY_SEMAPHORE_RESOURCE :
    WAIT ON BINARY_SEMAPHORE(B_SEM_PTR);
EVENT_RESOURCE :
    WAIT ON EVENT(EVENT_PTR);
CONDITION_VARIABLE_RESOURCE :
    WAIT ON CONDITION_VARIABLE(C_VAR_PTR);
end; (* case...of *)
SIGNAL_RESOURCE_2 :
with TASK(TASK_ID)
    INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 2) do
case RESOURCE_TYPE of
MESSAGE_RESOURCE :
    SEND MESSAGE TO(RESOURCE_TASK_ID, pointer(MSG_PTR));
SEMAPHORE_RESOURCE :
    SIGNAL_SEMAPHORE(SEM_PTR);
BINARY_SEMAPHORE_RESOURCE :
    SIGNAL_BINARY_SEMAPHORE(B_SEM_PTR);
EVENT_RESOURCE :
    SIGNAL_EVENT(EVENT_PTR);
CONDITION_VARIABLE_RESOURCE :
    SIGNAL_CONDITION_VARIABLE(C_VAR_PTR);
end; (* case...of *)
else (* Do Nothing *)
end; (* case...of *)
end; (* for...to...do *)

WAIT ON EVENT(PERIODIC_EVENT);
STOP_TIME := GET_MILLISECOND_TICKS;
if (STOP_TIME - START_TIME) <= PERIODIC_DELAY.MILLISECONDS then
inc(PERIODIC_TIME_MONITOR.SUCCESSFUL);
inc(PERIODIC_TIME_MONITOR.TOTAL);
until false;
end; (* with...do *)
end; (* GENERAL_TASK *)

procedure PROCESS_APPLICATION_TASK_SPECIFICATION_FILE;
(* This routine parses the .CSV file and creates the task specification *)
(* database and creates the application tasks. *)
type
(* Definition of the initial priority assignment gene (other gene *)
(* definitions come directly from the TASKING unit). *)
PRIORITY_ASSIGNMENT_ALGORITHMME
(* All tasks are initially assigned the same priority. *)
UNIFORM_ASSIGNMENT;
(* All tasks are initially assigned random priorities. *)
RANDOM_ASSIGNMENT;
(* Task priorities are assigned based on execution rates: higher *)
(* execution rate implies higher priority. *)
RATE_MONOTONIC_ASSIGNMENT;
(* Task priorities are assigned based on execution deadlines: *)
(* earlier execution deadline implies higher priority. *)
DEADLINE_MONOTONIC_ASSIGNMENT;
(* Task priorities are assigned based on workload level: higher *)
(* workload implies higher priority. *)
WORKLOAD_MONOTONIC_ASSIGNMENT;

var
(* Used to create the application tasks. *)
TASK_ID : TASK_IDS;
TASK_ATTR : TASK_ATTRIBUTES;

(* Used when determining how the task priorities should be assigned. *)
INITIAL_PRIORITY_ASSIGNMENT : PRIORITY_ASSIGNMENT_ALGORITHMMS;

(* Used to get the random number generator seed so that the results *)
(* are reproducible. *)
RANDOM_NUMBER_SEED : longint;

(* Used to assign initial priorities (i.e. Workload and Deadline *)
(* Monotonic. *)
MAXIMUM_MSEC_DEADLINE : integer;
MAXIMUM_MSEC_WORKLOAD : integer;

(* Used to hold the actual priorities of the application tasks. *)
PRIORITIES : array [1..MAX_TASKS] of USER_PRIORITIES;

LCV : integer;

function UP_STRING(S : string) : string;
(* This routine returns a string corresponding to the upper case value *)
(* of the string passed to it. *)

```

```

type (
  (* The .CSV file organization is very strict. The tasks "must" *)
  (* be specified with the following characteristics and they "must" *)
  (* be in the order listed below. *)
  .....)
FIELDS =
  (
    Field Name           Type      Range      Size (Chars)  *)
TASK_NAME_FIELD,      (* Text      0-1000    25             *)
PERIOD_FIELD,        (* mSec      0-1000    5              *)
DEADLINE_FIELD,     (* mSec      0-1000    5              *)
HARDNESS_FIELD,     (* mSec      0-1000    5              *)
WORKLOAD_FIELD,     (* mSec      0-1000    5              *)
RECEIVE_RESOURCE_1_FIELD, (* mSec    0-1000    5              *)
RECEIVE_RESOURCE_TYPE_1_FIELD, (* Text    0-1000    10             *)
RECEIVE_RESOURCE_TIME_1_FIELD, (* mSec    0-1000    5              *)
SIGNAL_RESOURCE_1_FIELD, (* Text      0-1000    10             *)
SIGNAL_RESOURCE_TYPE_1_FIELD, (* Text      0-1000    10             *)
RECEIVE_RESOURCE_2_FIELD, (* Text      0-1000    5              *)
RECEIVE_RESOURCE_TYPE_2_FIELD, (* Text      0-1000    10             *)
RECEIVE_RESOURCE_TIME_2_FIELD, (* mSec    0-1000    5              *)
SIGNAL_RESOURCE_2_FIELD, (* Text      0-1000    10             *)
SIGNAL_RESOURCE_TYPE_2_FIELD, (* Text      0-1000    10             *)
SIGNAL_RESOURCE_TIME_2_FIELD, (* mSec    0-1000    5              *)
);

const (
  (* CSV file format is 'comma separated text'. *)
  .....)
COMMA = ',';

var (
  (* CSV file parsing variables. *)
  .....)
LINE       : string;
CSV_FILE   : text;
POSITION   : integer;

(* These are used as the task specification is parsed (i.e. as *)
(* the fields are stripped from the line of text). *)
PARAMETER : string;
PARAM_LCV : FIELDS;

function INTERPRET_FIELD(THE_FIELD : FIELDS; var FIELD_STR : string;
                        TASK_NUMBER : integer) : boolean;
(* This routine examines the contents of the field and determines if it *)
(* is valid. Note that very little validation is actually performed. *)

var (
  (* Temporary variable to hold function return value. *)
  .....)
VALID_RESULT : boolean;

(* Used to determine the validity of numeric fields. *)
VALUE, CODE : integer;

begin (* INTERPRET_FIELD *)
  VALID_RESULT := true;

  with TASK(TASK_NUMBER) do
  case THE_FIELD of
    TASK_NAME_FIELD :
      TASK_NAME := FIELD_STR;
      PERIOD_FIELD :
        begin
          val(FIELD_STR, VALUE, CODE);
          if (CODE = 0) and (VALUE > 0)
          then
            PERIOD := VALUE
          else
            VALID_RESULT := false;
          end; (* Period *)
        DEADLINE_FIELD :
          begin
            val(FIELD_STR, VALUE, CODE);
            if (CODE = 0) and (VALUE >= 0)
            then
              DEADLINE := VALUE
            else
              VALID_RESULT := false;
            end; (* Deadline *)
          HARDNESS_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (1 <= VALUE) and (VALUE <= 10)
              then
                DEADLINE_HARDNESS := VALUE
              else
                VALID_RESULT := false;
              end; (* Hardness *)
            WORKLOAD_FIELD :
              begin
                val(FIELD_STR, VALUE, CODE);
                if (CODE = 0) and (VALUE > 0)
                then
                  WORKLOAD := VALUE
                else
                  VALID_RESULT := false;
                end; (* Workload *)
          RECEIVE_RESOURCE_1_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 1) do
            RESOURCE_NAME := FIELD_STR;
          RECEIVE_RESOURCE_TYPE_1_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 1) do
            begin
              if FIELD_STR = 'EVENT' then
                RESOURCE_TYPE := EVENT_RESOURCE
              else if FIELD_STR = 'SEMAPHORE' then
                RESOURCE_TYPE := SEMAPHORE_RESOURCE
              else if FIELD_STR = 'MESSAGE' then
                RESOURCE_TYPE := MESSAGE_RESOURCE
              else if FIELD_STR = 'BSEMAPHORE' then
                RESOURCE_TYPE := BINARY_SEMAPHORE_RESOURCE
              else if FIELD_STR = 'COND'VAR' then
                RESOURCE_TYPE := CONDITION_VARIABLE_RESOURCE
              else if FIELD_STR <> 'NONE' then
                VALID_RESULT := false;
              DEFINED := VALID_RESULT and (FIELD_STR <> 'NONE');
            end; (* Resource Type *)
          RECEIVE_RESOURCE_TIME_1_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (VALUE >= 0)
              then
                INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 1),
                RESOURCE_TIME := VALUE
              else
                VALID_RESULT := false;
            end; (* Communication Time *)
          SIGNAL_RESOURCE_1_FIELD :
            with INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1) do
            RESOURCE_NAME := FIELD_STR;
          SIGNAL_RESOURCE_TYPE_1_FIELD :
            with INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1) do
            begin
              if FIELD_STR = 'EVENT' then
                RESOURCE_TYPE := EVENT_RESOURCE
              else if FIELD_STR = 'SEMAPHORE' then
                RESOURCE_TYPE := SEMAPHORE_RESOURCE
              else if FIELD_STR = 'MESSAGE' then
                RESOURCE_TYPE := MESSAGE_RESOURCE
              else if FIELD_STR = 'BSEMAPHORE' then
                RESOURCE_TYPE := BINARY_SEMAPHORE_RESOURCE
              else if FIELD_STR = 'COND'VAR' then
                RESOURCE_TYPE := CONDITION_VARIABLE_RESOURCE
              else if FIELD_STR <> 'NONE' then
                VALID_RESULT := false;
              DEFINED := VALID_RESULT and (FIELD_STR <> 'NONE');
            end; (* Resource Type *)
          SIGNAL_RESOURCE_TIME_1_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (VALUE >= 0)
              then
                INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1),
                RESOURCE_TIME := VALUE
              else
                VALID_RESULT := false;
            end; (* Communication Time *)
          SIGNAL_RESOURCE_2_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 2) do
            RESOURCE_NAME := FIELD_STR;
          RECEIVE_RESOURCE_TYPE_2_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 2) do
            begin
              if FIELD_STR = 'EVENT' then
                RESOURCE_TYPE := EVENT_RESOURCE
              else if FIELD_STR = 'SEMAPHORE' then
                RESOURCE_TYPE := SEMAPHORE_RESOURCE
              else if FIELD_STR = 'MESSAGE' then
                RESOURCE_TYPE := MESSAGE_RESOURCE
              else if FIELD_STR = 'BSEMAPHORE' then
                RESOURCE_TYPE := BINARY_SEMAPHORE_RESOURCE
              else if FIELD_STR = 'COND'VAR' then
                RESOURCE_TYPE := CONDITION_VARIABLE_RESOURCE
              else if FIELD_STR <> 'NONE' then
                VALID_RESULT := false;
              DEFINED := VALID_RESULT and (FIELD_STR <> 'NONE');
            end; (* Resource Type *)
          SIGNAL_RESOURCE_TIME_2_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (VALUE >= 0)
              then
                INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 2),
                RESOURCE_TIME := VALUE
              else
                VALID_RESULT := false;
            end; (* Communication Time *)
          else VALID_RESULT := true;
        end; (* case...of *)
      INTERPRET_FIELD := VALID_RESULT;
    end; (* INTERPRET_FIELD *)

begin (* PROCESS_CSV_FILE *)
  (* Remove report file from previous run (if it exists) *)
  assign(CSV_FILE, PROGRAM_NAME + '.RPT');
  (*SI-*) reset(CSV_FILE); (*SI+*)
  if IORESULT = 0 then
    begin
      close(CSV_FILE);
      erase(CSV_FILE);
    end; (* if...then *)
  assign(CSV_FILE, PROGRAM_NAME + '.CSV');
  (*SI-*) reset(CSV_FILE); (*SI+*)
  if IORESULT = 0 then
    begin
      NUMBER_OF_TASKS := 0;
      while not eof(CSV_FILE) do
        begin
          readln(CSV_FILE, LINE);
          (* Make sure there is a comma at the end of the line *)
          LINE := LINE + COMMA;
          while LINE(1) = ' ' do delete(LINE, 1, 1);
          if UP_STRING(copy(LINE, 1, pos(COMMA, LINE)-1)) = 'TASK' then
            begin
              delete(LINE, 1, pos(COMMA, LINE));
              (* Parse valid task identification line *)
              inc(NUMBER_OF_TASKS);
              PARAM_LCV := low(FIELDS);
              repeat
                PARAMETER := UP_STRING(copy(LINE, 1, pos(COMMA, LINE)-1));
                if not INTERPRET_FIELD(PARAM_LCV, PARAMETER,
                                      NUMBER_OF_TASKS) then
                  halt(1);
                PARAM_LCV := succ(PARAM_LCV);
                delete(LINE, 1, pos(COMMA, LINE));
              until (length(LINE) = 0) or (PARAM_LCV > high(FIELDS));
              end; (* if...then *)
            end; (* while...do *)
          close(CSV_FILE);
        end; (* if...then *)
      end; (* PROCESS_CSV_FILE *)

begin (* PROCESS_APPLICATION_TASK_SPECIFICATION_FILE *)
  PROCESS_INI_FILE;

  (* Setup default task parameters *)
  with TASK_ATTR do
  begin
with INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1) do
begin
  (* The .CSV file organization is very strict. The tasks "must" *)
  (* be specified with the following characteristics and they "must" *)
  (* be in the order listed below. *)
  .....)
FIELDS =
  (
    Field Name           Type      Range      Size (Chars)  *)
TASK_NAME_FIELD,      (* Text      0-1000    25             *)
PERIOD_FIELD,        (* mSec      0-1000    5              *)
DEADLINE_FIELD,     (* mSec      0-1000    5              *)
HARDNESS_FIELD,     (* mSec      0-1000    5              *)
WORKLOAD_FIELD,     (* mSec      0-1000    5              *)
RECEIVE_RESOURCE_1_FIELD, (* mSec    0-1000    5              *)
RECEIVE_RESOURCE_TYPE_1_FIELD, (* Text    0-1000    10             *)
RECEIVE_RESOURCE_TIME_1_FIELD, (* mSec    0-1000    5              *)
SIGNAL_RESOURCE_1_FIELD, (* Text      0-1000    10             *)
SIGNAL_RESOURCE_TYPE_1_FIELD, (* Text      0-1000    10             *)
RECEIVE_RESOURCE_2_FIELD, (* Text      0-1000    5              *)
RECEIVE_RESOURCE_TYPE_2_FIELD, (* Text      0-1000    10             *)
RECEIVE_RESOURCE_TIME_2_FIELD, (* mSec    0-1000    5              *)
SIGNAL_RESOURCE_2_FIELD, (* Text      0-1000    10             *)
SIGNAL_RESOURCE_TYPE_2_FIELD, (* Text      0-1000    10             *)
SIGNAL_RESOURCE_TIME_2_FIELD, (* mSec    0-1000    5              *)
);

const (
  (* CSV file format is 'comma separated text'. *)
  .....)
COMMA = ',';

var (
  (* CSV file parsing variables. *)
  .....)
LINE       : string;
CSV_FILE   : text;
POSITION   : integer;

(* These are used as the task specification is parsed (i.e. as *)
(* the fields are stripped from the line of text). *)
PARAMETER : string;
PARAM_LCV : FIELDS;

function INTERPRET_FIELD(THE_FIELD : FIELDS; var FIELD_STR : string;
                        TASK_NUMBER : integer) : boolean;
(* This routine examines the contents of the field and determines if it *)
(* is valid. Note that very little validation is actually performed. *)

var (
  (* Temporary variable to hold function return value. *)
  .....)
VALID_RESULT : boolean;

(* Used to determine the validity of numeric fields. *)
VALUE, CODE : integer;

begin (* INTERPRET_FIELD *)
  VALID_RESULT := true;

  with TASK(TASK_NUMBER) do
  case THE_FIELD of
    TASK_NAME_FIELD :
      TASK_NAME := FIELD_STR;
      PERIOD_FIELD :
        begin
          val(FIELD_STR, VALUE, CODE);
          if (CODE = 0) and (VALUE > 0)
          then
            PERIOD := VALUE
          else
            VALID_RESULT := false;
          end; (* Period *)
        DEADLINE_FIELD :
          begin
            val(FIELD_STR, VALUE, CODE);
            if (CODE = 0) and (VALUE >= 0)
            then
              DEADLINE := VALUE
            else
              VALID_RESULT := false;
            end; (* Deadline *)
          HARDNESS_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (1 <= VALUE) and (VALUE <= 10)
              then
                DEADLINE_HARDNESS := VALUE
              else
                VALID_RESULT := false;
              end; (* Hardness *)
            WORKLOAD_FIELD :
              begin
                val(FIELD_STR, VALUE, CODE);
                if (CODE = 0) and (VALUE > 0)
                then
                  WORKLOAD := VALUE
                else
                  VALID_RESULT := false;
                end; (* Workload *)
          RECEIVE_RESOURCE_1_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 1) do
            RESOURCE_NAME := FIELD_STR;
          RECEIVE_RESOURCE_TYPE_1_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 1) do
            begin
              if FIELD_STR = 'EVENT' then
                RESOURCE_TYPE := EVENT_RESOURCE
              else if FIELD_STR = 'SEMAPHORE' then
                RESOURCE_TYPE := SEMAPHORE_RESOURCE
              else if FIELD_STR = 'MESSAGE' then
                RESOURCE_TYPE := MESSAGE_RESOURCE
              else if FIELD_STR = 'BSEMAPHORE' then
                RESOURCE_TYPE := BINARY_SEMAPHORE_RESOURCE
              else if FIELD_STR = 'COND'VAR' then
                RESOURCE_TYPE := CONDITION_VARIABLE_RESOURCE
              else if FIELD_STR <> 'NONE' then
                VALID_RESULT := false;
              DEFINED := VALID_RESULT and (FIELD_STR <> 'NONE');
            end; (* Resource Type *)
          RECEIVE_RESOURCE_TIME_1_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (VALUE >= 0)
              then
                INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 1),
                RESOURCE_TIME := VALUE
              else
                VALID_RESULT := false;
            end; (* Communication Time *)
          SIGNAL_RESOURCE_1_FIELD :
            with INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1) do
            RESOURCE_NAME := FIELD_STR;
          SIGNAL_RESOURCE_TYPE_1_FIELD :
            with INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1) do
            begin
              if FIELD_STR = 'EVENT' then
                RESOURCE_TYPE := EVENT_RESOURCE
              else if FIELD_STR = 'SEMAPHORE' then
                RESOURCE_TYPE := SEMAPHORE_RESOURCE
              else if FIELD_STR = 'MESSAGE' then
                RESOURCE_TYPE := MESSAGE_RESOURCE
              else if FIELD_STR = 'BSEMAPHORE' then
                RESOURCE_TYPE := BINARY_SEMAPHORE_RESOURCE
              else if FIELD_STR = 'COND'VAR' then
                RESOURCE_TYPE := CONDITION_VARIABLE_RESOURCE
              else if FIELD_STR <> 'NONE' then
                VALID_RESULT := false;
              DEFINED := VALID_RESULT and (FIELD_STR <> 'NONE');
            end; (* Resource Type *)
          SIGNAL_RESOURCE_TIME_1_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (VALUE >= 0)
              then
                INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 1),
                RESOURCE_TIME := VALUE
              else
                VALID_RESULT := false;
            end; (* Communication Time *)
          SIGNAL_RESOURCE_2_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 2) do
            RESOURCE_NAME := FIELD_STR;
          RECEIVE_RESOURCE_TYPE_2_FIELD :
            with INTERTASK_COMMUNICATION(RECEIVE_RESOURCE, 2) do
            begin
              if FIELD_STR = 'EVENT' then
                RESOURCE_TYPE := EVENT_RESOURCE
              else if FIELD_STR = 'SEMAPHORE' then
                RESOURCE_TYPE := SEMAPHORE_RESOURCE
              else if FIELD_STR = 'MESSAGE' then
                RESOURCE_TYPE := MESSAGE_RESOURCE
              else if FIELD_STR = 'BSEMAPHORE' then
                RESOURCE_TYPE := BINARY_SEMAPHORE_RESOURCE
              else if FIELD_STR = 'COND'VAR' then
                RESOURCE_TYPE := CONDITION_VARIABLE_RESOURCE
              else if FIELD_STR <> 'NONE' then
                VALID_RESULT := false;
              DEFINED := VALID_RESULT and (FIELD_STR <> 'NONE');
            end; (* Resource Type *)
          SIGNAL_RESOURCE_TIME_2_FIELD :
            begin
              val(FIELD_STR, VALUE, CODE);
              if (CODE = 0) and (VALUE >= 0)
              then
                INTERTASK_COMMUNICATION(SIGNAL_RESOURCE, 2),
                RESOURCE_TIME := VALUE
              else
                VALID_RESULT := false;
            end; (* Communication Time *)
          else VALID_RESULT := true;
        end; (* case...of *)
      INTERPRET_FIELD := VALID_RESULT;
    end; (* INTERPRET_FIELD *)

begin (* PROCESS_CSV_FILE *)
  (* Remove report file from previous run (if it exists) *)
  assign(CSV_FILE, PROGRAM_NAME + '.RPT');
  (*SI-*) reset(CSV_FILE); (*SI+*)
  if IORESULT = 0 then
    begin
      close(CSV_FILE);
      erase(CSV_FILE);
    end; (* if...then *)
  assign(CSV_FILE, PROGRAM_NAME + '.CSV');
  (*SI-*) reset(CSV_FILE); (*SI+*)
  if IORESULT = 0 then
    begin
      NUMBER_OF_TASKS := 0;
      while not eof(CSV_FILE) do
        begin
          readln(CSV_FILE, LINE);
          (* Make sure there is a comma at the end of the line *)
          LINE := LINE + COMMA;
          while LINE(1) = ' ' do delete(LINE, 1, 1);
          if UP_STRING(copy(LINE, 1, pos(COMMA, LINE)-1)) = 'TASK' then
            begin
              delete(LINE, 1, pos(COMMA, LINE));
              (* Parse valid task identification line *)
              inc(NUMBER_OF_TASKS);
              PARAM_LCV := low(FIELDS);
              repeat
                PARAMETER := UP_STRING(copy(LINE, 1, pos(COMMA, LINE)-1));
                if not INTERPRET_FIELD(PARAM_LCV, PARAMETER,
                                      NUMBER_OF_TASKS) then
                  halt(1);
                PARAM_LCV := succ(PARAM_LCV);
                delete(LINE, 1, pos(COMMA, LINE));
              until (length(LINE) = 0) or (PARAM_LCV > high(FIELDS));
              end; (* if...then *)
            end; (* while...do *)
          close(CSV_FILE);
        end; (* if...then *)
      end; (* PROCESS_CSV_FILE *)

begin (* PROCESS_APPLICATION_TASK_SPECIFICATION_FILE *)
  PROCESS_INI_FILE;

  (* Setup default task parameters *)
  with TASK_ATTR do
  begin

```



```

STACK_WORDS_NEEDED := 4000;
(* Use default error handlers *)
ERROR_HANDLERS[TASK_ALREADY_ACTIVE] := nil;
ERROR_HANDLERS[INSUFFICIENT_RESOURCES] := nil;
ERROR_HANDLERS[TASK_IS_NOT_ACTIVE] := nil;
ERROR_HANDLERS[TASK_ALREADY_SUSPENDED] := nil;
ERROR_HANDLERS[ILLEGAL_TASK_ID] := nil;
ERROR_HANDLERS[ILLEGAL_OPERATION] := nil;
end; (* with...do *)

PROCESS_CS_V_FILE;

(* Determine maximum execution period, deadline and workload *)
MAXIMUM_MSEC_PERIOD := 1;
MAXIMUM_MSEC_DEADLINE := 1;
MAXIMUM_MSEC_WORKLOAD := 1;
for LCV := 1 to NUMBER_OF_TASKS do
begin
if TASK[LCV].PERIOD > MAXIMUM_MSEC_PERIOD then
MAXIMUM_MSEC_PERIOD := TASK[LCV].PERIOD;
if TASK[LCV].DEADLINE > MAXIMUM_MSEC_DEADLINE then
MAXIMUM_MSEC_DEADLINE := TASK[LCV].DEADLINE;
if TASK[LCV].WORKLOAD > MAXIMUM_MSEC_WORKLOAD then
MAXIMUM_MSEC_WORKLOAD := TASK[LCV].WORKLOAD;
end; (* for...to...do *)

(* Assign initial task priorities *)
case INITIAL_PRIORITY_ASSIGNMENT of
UNIFORM_ASSIGNMENT :
for LCV := 1 to NUMBER_OF_TASKS do
PRIORITIES[LCV] := (high(USER_PRIORITIES) + low(USER_PRIORITIES))
div 2;
RANDOM_ASSIGNMENT :
begin
randseed := RANDOM_NUMBER_SEED;
for LCV := 1 to NUMBER_OF_TASKS do
PRIORITIES[LCV] := random(high(USER_PRIORITIES) -
low(USER_PRIORITIES)) +
low(USER_PRIORITIES);
end; (* Random Assignment *)
RATE_MONOTONIC_ASSIGNMENT :
for LCV := 1 to NUMBER_OF_TASKS do
PRIORITIES[LCV] := low(USER_PRIORITIES) + high(USER_PRIORITIES) -
round(high(USER_PRIORITIES) /
MAXIMUM_MSEC_PERIOD * TASK[LCV].PERIOD);
DEADLINE_MONOTONIC_ASSIGNMENT :
for LCV := 1 to NUMBER_OF_TASKS do
PRIORITIES[LCV] := low(USER_PRIORITIES) + high(USER_PRIORITIES) -
round(high(USER_PRIORITIES) /
MAXIMUM_MSEC_DEADLINE * TASK[LCV].DEADLINE);
WORKLOAD_MONOTONIC_ASSIGNMENT :
for LCV := 1 to NUMBER_OF_TASKS do
PRIORITIES[LCV] := round((high(USER_PRIORITIES) - 1) /
MAXIMUM_MSEC_WORKLOAD * TASK[LCV].WORKLOAD);
end; (* case...of *)

for LCV := 1 to NUMBER_OF_TASKS do
begin
TASK_ATTR.PRIORITY := PRIORITIES[LCV];
CREATE(TASK_ID, TASK_ATTR, GENERAL_TASK);
TASK[LCV].TASK_ID := TASK_ID;
TASK[LCV].TASK_PRIORITY := TASK_ATTR.PRIORITY;
end; (* for...to...do *)
end; (* PROCESS_APPLICATION_TASK_SPECIFICATION_FILE *)

begin (* REAL_TIME_APPLICATION *)
(* Determine whether report file is required or not *)
GENERATE_REPORT_FILE := (paramcount = 0);

(* Determine base filename for input/output files *)
PROGNAME := paramstr(0);
delete(PROGNAME, length(PROGNAME) - 3, 4);

(* Establish task database and create application tasks *)
PROCESS_APPLICATION_TASK_SPECIFICATION_FILE;

(* Setup the supervisor task *)
with TASK_ATTR do
begin
PRIORITY := high(USER_PRIORITIES);
STACK_WORDS_NEEDED := 2000;
(* Use default error handlers *)
ERROR_HANDLERS[TASK_ALREADY_ACTIVE] := nil;
ERROR_HANDLERS[INSUFFICIENT_RESOURCES] := nil;
ERROR_HANDLERS[TASK_IS_NOT_ACTIVE] := nil;
ERROR_HANDLERS[TASK_ALREADY_SUSPENDED] := nil;
ERROR_HANDLERS[ILLEGAL_TASK_ID] := nil;
ERROR_HANDLERS[ILLEGAL_OPERATION] := nil;
end; (* with...do *)
CREATE(SUPERVISOR, TASK_ATTR, SUPERVISOR_TASK);
end. (* REAL_TIME_APPLICATION *)

```

## 9.2 GRAPH-GA.PAS

```

program GENETIC_ALGORITHM_THESIS(input, output);
(* This program uses a Genetic Algorithm (GA) to search the problem space *
 * for the best solution to the problem specified. *)
(* Compiler Options (Ver. 7.0) *)
(*$A+ Word Alignment *)
(*$B+ Short Circuit Boolean Evaluation *)
(*$D+ Debug Code Generation ON (Sort of) *)
(*$L+ Requires /V option to BPC to activate *)
(*$M+ Local Debug symbols ON (Sort of) *)
(*$E- Requires /V option to BPC to activate *)
(*$I- Far calls only as needed *)
(*$M $FFFO,$FFFO,$FFFO I/O Checking OFF *)
(*$N+ Software Emulation of 80x87 *)
(*$O- Overlays NOT allowed *)
(*$Q+ Standard 'string' parameters *)
(*$R- Overflow Checking OFF *)
(*$S- Range Checking OFF *)
(*$U- Stack Checking OFF *)
(*$T+ Var-string Checking OFF *)
(*$X+ Force Typed 'g' references *)
(*$+ Enable Extended syntax *)
uses dos, crt, graph, GRAPHICS, BMP_UTIL,
(*$IFDEF MATH_GA *)
MATH_GA;
(*$ELSE *)
RTOS_GA;
(*$ENDIF *)

const
(* Used to allow a single routine to be used as 'paint screen' or as *
 * 'update as needed' drawing routine. *)
FORCE_UPDATE = true;
DO_NOT_FORCE_UPDATE = false;

(* Command line controllable parameters. *)
(* Control flag used to specify that the reverse video colors *
 * should be used. *)
REVERSE_VIDEO = boolean = false;
(* Control flag used to specify that the display should be 640x480, *
 * 800x600 or 1024x768, 1280x1024. *)
VIDEO_MODES = NO_DISPLAY_MODE, VGA, SVGA, VESA, SVESA);
type
DISPLAY_MODE = VIDEO_MODES;
const
(* When in graphics mode the follow colors are used for drawing the *
 * screen. *)
GRAPHIC_COLORS = record
BACKGROUND_COLOR : word;
AXIS_COLOR : word;
GRID_COLOR : word;
AXIS_LABEL_COLOR : word;
GRAPH_COLOR : word;
STATUS_COLOR : word;
STATUS_HIGHLIGHT_COLOR : word;
MAIN_TITLE_COLOR : word;
BORDER_COLOR : word;
end; (* GRAPHIC_COLORS *)

const
(* INI file only configurable parameters. *)
(* Directory where IMAGE_XX.BMP files will be stored. *)
IMAGE_DIRECTORY = 'disc';
(* The graphic display uses the following colors in normal video. *)
NORMAL_GRAPHIC_COLORS =
(
BACKGROUND_COLOR : black;
AXIS_COLOR : lightgray;
GRID_COLOR : darkgray;
AXIS_LABEL_COLOR : white;
GRAPH_COLOR : cyan;
STATUS_COLOR : yellow;
STATUS_HIGHLIGHT_COLOR : white;
MAIN_TITLE_COLOR : lightcyan;
BORDER_COLOR : green
);
(* The graphic display uses the following colors in reverse video. *)
REVERSE_GRAPHIC_COLORS =
(
BACKGROUND_COLOR : white;
AXIS_COLOR : darkgray;
GRID_COLOR : lightgray;
AXIS_LABEL_COLOR : black;
GRAPH_COLOR : blue;
STATUS_COLOR : red;
STATUS_HIGHLIGHT_COLOR : lightred;
MAIN_TITLE_COLOR : lightblue;
BORDER_COLOR : green
);
(* The maximum number of graphs on the screen at any one time *)
MAX_NUMBER_OF_GRAPHS = 3;

(* User controllable parameters (run-time controls). *)
(* Controls how many graphs are being displayed. *)
NUMBER_OF_GRAPHS_TO_DISPLAY : 0..MAX_NUMBER_OF_GRAPHS = 0;
(* Allows the user to pause the display. *)
USER_PAUSED : boolean = false;
(* Allows for an orderly shutdown of the application. *)
USER_QUIT : boolean = false;

(* The following constants are used to control sizes/offsets of the *)
(* graphic screen. *)
POINT_SIZE = 3; TICK_LENGTH = 4; EDGE_OFFSET = 4; CURSOR_SIZE = 5;
(* MS-DOS exit codes (i.e. 'errorlevel'). *)
INVALID_PARAMETER_ERROR_LEVEL = 20;
NON_SUPPORTED_VIDEO_MODE_ERROR_LEVEL = 30;
(* General purpose special characters. *)
BELL = chr($07); (* Bell character *)
SPACE = chr($20); (* Blank character *)
CR = chr($0D); (* Carriage Return *)
LF = chr($0A); (* Line Feed *)
BS = chr($08); (* Backspace *)
ESC = chr($1B); (* Escape *)
NULL = chr($00); (* Precedes extended keys *)
type
(* The axis cross point can be placed anywhere on the graph. *)
AXIS_TYPES = (AUTO_AXIS, FORCE_LL, FORCE_LR, FORCE_UL, FORCE_UR);
(* The graph can be positioned at various places on the screen. *)
GRAPH_POSITIONS = (FULL_SCREEN, TOP_HALF, BOTTOM_HALF, TOP_THIRD,
MIDDLE_THIRD, BOTTOM_THIRD);
(* Type to specify parameters for the two axis. *)
AXIS = (X_AXIS, Y_AXIS);
(* Type to specify how the graph should be drawn. *)
GRAPH_TYPES = (BAR_GRAPH, LINE_GRAPH);
(* Graph lines are a sequence of points defined by this type. *)
GRAPH_PTR = ^POINTS;
POINTS = array(0..POPULATION_SIZE) of record
POSITION : array[AXIS] of real;
end; (* LINES *)
(* The user controllable attributes (as well as some fixed parameters) *
 * of the graphs are kept in this record type. *)
GRAPH_SCALE_INFO = record
(* Main title to the graph. *)
TITLE : string[25];
(* Line or bar graph indicator. *)
GRAPH_TYPE : GRAPH_TYPES;
(* FIRST and LAST refer to the indices of the VALUES array. *)
FIRST, LAST : integer;
(* For each axis the maximum, minimum and number of tick marks *
 * must be specified. The range is calculated as the graph is setup. *)
HIGH_LABEL, LOW_LABEL,
RANGE, NUMBER_OF_TICK_MARKS : array[AXIS] of longint;
(* The label for the axis units. *)
LABEL_NAME : array[AXIS] of string[10];
(* Used to turn on/off peak detection for each of the axis. *)
PEAK_DETECTION : array[AXIS] of boolean;
(* Used to turn on/off average detection for the Y-Axis. *)
AVERAGE_DETECTION : boolean;
end; (* GRAPH_SCALE_INFO *)
(* Data structure that keeps all of the information about a graph so *
 * that the graph display functions can access the correct graph when *
 * more than one is being displayed at a time. *)
GRAPH_INFO = record
(* Scaling for X and Y axis (see above). *)
GRAPH_SCALES : GRAPH_SCALE_INFO;
(* X and Y values for the previous peaks/averages (allows peaks/ *
 * averages to be removed). *)
PREVIOUS_PEAK, PREVIOUS_PEAK_VALUE, PREVIOUS_AVERAGE_VALUE : real;
(* Offset of the graph into the graphic window on the screen. *)
HORIZONTAL_OFFSET, VERTICAL_OFFSET : real;
(* Sizes of the graphic window on the screen (function of the *
 * video mode selected). *)
HORIZONTAL_SIZE, VERTICAL_SIZE : real;
(* Offset of the graph axis within the graphic window. *)
X_AXIS_OFFSET, Y_AXIS_OFFSET : real;
(* The actual data that is being graphed, and the array limits. *)
VALUES : record
FIRST, LAST : integer;
DATA : array(0..POPULATION_SIZE) of real;
end; (* VALUES *)
(* Pointers to graphic points that are being graphed after the *
 * appropriate conversion from value to pixel has been done. *)
GRAPH_POINTS, PREVIOUS_GRAPH_POINTS : GRAPH_PTR;
(* The color of the line being graphed. *)
GRAPH_COLOR : word;
end; (* GRAPH_INFO *)
var
(* Maximum X and Y values for the graphics mode selected. *)
MAX_X, MAX_Y : integer;

```

```

(* The data structure for the graphs being displayed. *)
.....
PLOT : array (1..MAX_NUMBER_OF_GRAPHS) of GRAPH_INFO;
.....
(* The graphic display uses the following colors variables for both *)
(* normal and reversed color modes. *)
.....
SCREEN : GRAPHIC_COLORS;
.....
type
(* The user is allowed to change the scales being used, this data *)
(* structure holds the user choices. *)
.....
USER_GRAPH_SCALES = record
(* Contains all of the graph axis information. *)
.....
SCALES : GRAPH_SCALE_INFO;
(* Used to turn on/off peak detection for each of the axis. *)
.....
PEAK_DETECTION : array (AXIS) of boolean;
(* Used to turn on/off average detection for the Y-axis. *)
.....
AVERAGE_DETECTION : boolean;
end; (* USER_GRAPH_SCALES *)
var
USER_GRAPH : array (1..MAX_NUMBER_OF_GRAPHS) of USER_GRAPH_SCALES;
.....
(* The offset from the edge of the screen to the graph is a function *)
(* of the graphics mode selected. *)
.....
GRAPH_OFFSET : integer; MESSAGE_REGION : integer;
.....
(* Graphic Video Driver mode parameter *)
.....
GRAPHICS_MODE : integer;
.....
(* Address of previous exit procedure in the exit call chain. *)
.....
SAVE_EXIT : pointer;
.....
function INT_STR(I: longint): string;
.....
(* Convert any integer type to a string. *)
.....
var S : string(11);
begin (* INT_STR *)
str(I, S);
INT_STR := S;
end; (* INT_STR *)
function REAL_STR(R: real; F1, F2 : integer): string;
.....
(* Convert real type to a string (same format as write(R : F1 : F2)). *)
.....
var S : string(11);
begin (* REAL_STR *)
str(R : F1 : F2, S);
REAL_STR := S;
end; (* REAL_STR *)
function X_POSITION(var X : real; GRAPH_NUMBER : integer) : real;
.....
(* Determines the X position within the graphics window that corresponds *)
(* to the X value passed in. *)
.....
var
(* Temporary graphic screen position. *)
.....
POSITION : real;
begin (* X_POSITION *)
with PLOT(GRAPH_NUMBER) do
with GRAPH_SCALES do
begin
X := X - LOW_LABEL(X_AXIS) - 1;
POSITION := HORIZONTAL_OFFSET + 1 +
(X * (HORIZONTAL_SIZE) / RANGE(X_AXIS));
if (HORIZONTAL_OFFSET <= POSITION) and
(POSITION <= (HORIZONTAL_OFFSET + HORIZONTAL_SIZE))
then
X := X + LOW_LABEL(X_AXIS) - 1
else
X := LOW_LABEL(X_AXIS) + 1;
X_POSITION := POSITION;
end; (* with...do *)
end; (* X_POSITION *)
function X_VALUE(POSITION : integer; GRAPH_NUMBER : integer) : real;
.....
(* Determines the value of the function at the graphic position passed in. *)
(* this function assumes the entire screen is the viewport. *)
.....
begin (* X_VALUE *)
with PLOT(GRAPH_NUMBER) do
with GRAPH_SCALES do
begin
X_VALUE := LOW_LABEL(X_AXIS) +
(POSITION - HORIZONTAL_OFFSET) /
HORIZONTAL_SIZE * RANGE(X_AXIS);
end; (* with...do *)
end; (* X_VALUE *)
function Y_POSITION(var Y : real; GRAPH_NUMBER : integer) : real;
.....
(* Determines the Y position within the graphics window that corresponds *)
(* to the Y value passed in. *)
.....
var
(* Temporary graphic screen position. *)
.....
POSITION : real;
begin (* Y_POSITION *)
with PLOT(GRAPH_NUMBER) do
with GRAPH_SCALES do
begin
Y := Y - LOW_LABEL(Y_AXIS) - 1;
POSITION := VERTICAL_OFFSET + 1 +
(VERTICAL_SIZE - (Y * VERTICAL_SIZE) / RANGE(Y_AXIS));
if POSITION > (VERTICAL_OFFSET + VERTICAL_SIZE - 1) then
Y_POSITION := VERTICAL_OFFSET + VERTICAL_SIZE - 1
else if POSITION < (VERTICAL_OFFSET + 1) then
Y_POSITION := VERTICAL_OFFSET + 1
else
Y_POSITION := POSITION;
end; (* Y_POSITION *)
Y := Y + LOW_LABEL(Y_AXIS) + 1;
end; (* with...do *)
end; (* Y_POSITION *)
(* This function assumes the entire screen is the viewport. *)
.....
begin (* Y_VALUE *)
with PLOT(GRAPH_NUMBER) do
with GRAPH_SCALES do
begin
Y_VALUE := LOW_LABEL(Y_AXIS) +
(VERTICAL_SIZE + VERTICAL_OFFSET - POSITION) /
VERTICAL_SIZE * RANGE(Y_AXIS);
end; (* with...do *)
end; (* Y_VALUE *)
procedure UPDATE_STATUS_MESSAGES(FORCE_UPDATE : boolean);
.....
(* Updates the status messages with the present value of the status flags. *)
.....
type
(* This type allows smaller strings to be manipulated for the date *)
(* status since the date is a fixed length. *)
.....
DATE_STRINGS = string(8);
.....
(* Type used to contain strings for the present PC clock time. *)
.....
TIME_STRINGS = string(12);
.....
const
(* The previous status parameters are maintained static so that *)
(* the update is flicker-free. *)
.....
PREVIOUS_USER_PAUSED : boolean = false;
PREVIOUS_DATE : DATE_STRINGS = '';
DATE : DATE_STRINGS = '';
PREVIOUS_INDIVIDUAL : integer = POPULATION_SIZE;
PREVIOUS_GENERATION : integer = -1;
.....
(* Values of the present and previously displayed PC clock time *)
(* the update is flicker-free. *)
.....
PREVIOUS_TIME : TIME_STRINGS = '';
TIME : TIME_STRINGS = '';
function TIME_NOW : TIME_STRINGS;
.....
(* This routine returns a string corresponding to the present TIME in *)
(* the PC. The format is 'HH:MM:SS am/pm'. *)
.....
var
(* Used to determine the present time (as reported by the PC). *)
.....
HR, MIN, SEC, SECL00TH : word;
.....
(* Temporary variables to hold the intermediate string results. *)
.....
S : TIME_STRINGS; TEMP : string(6);
begin (* TIME_NOW *)
gettime(HR, MIN, SEC, SECL00TH);
str(HR mod 12, S);
if S = '0' then S := '12';
S := S + ':';
if MIN < 10 then S := S + '0';
str(MIN, TEMP);
S := S + TEMP;
S := S + ':';
if SEC < 10 then S := S + '0';
str(SEC, TEMP);
S := S + TEMP;
if (HR div 12) = 0
then
TIME_NOW := S + ' am'
else
TIME_NOW := S + ' pm';
end; (* TIME_NOW *)
function DATE_NOW : DATE_STRINGS;
.....
(* This routine returns a string corresponding to the present DATE in *)
(* the PC. The string is of the form 'MM-DD-YY'. *)
.....
var
(* Used to determine the present date (as reported by the PC). *)
.....
YEAR, MONTH, DAY, DAY_OF_WEEK : word;
.....
(* Temporary variable to hold the intermediate string results. *)
.....
S : DATE_STRINGS;
begin (* DATE_NOW *)
getdate(YEAR, MONTH, DAY, DAY_OF_WEEK);
S := '';
S := chr(ord(MONTH div 10) + ord('0'));
S := S + chr(ord(MONTH mod 10) + ord('0')) + '-';
S := S + chr(ord(DAY div 10) + ord('0'));
S := S + chr(ord(DAY mod 10) + ord('0')) + '-';
S := S + chr(ord(YEAR mod 100 div 10) + ord('0'));
S := S + chr(ord(YEAR mod 10) + ord('0'));
DATE_NOW := S;
end; (* DATE_NOW *)
begin (* UPDATE_STATUS_MESSAGES *)
(* Update Time string *)
settextstyle(defaultfont, horizdir, 1);
settextjustify(righttext, centertext);
TIME := TIME_NOW;
if FORCE_UPDATE or (TIME < PREVIOUS_TIME) then
begin
setcolor(SCREEN.BACKGROUND_COLOR);
outtextxy
[
MAX_X - 2 * EDGE_OFFSET,
MAX_Y + 2 * MESSAGE_REGION div 3,
PREVIOUS_TIME
];
PREVIOUS_TIME := TIME;
setcolor(SCREEN.STATUS_COLOR);
outtextxy
[
MAX_X - 2 * EDGE_OFFSET,
MAX_Y + 2 * MESSAGE_REGION div 3,
PREVIOUS_TIME
];
end; (* if...then *)
end; (* UPDATE_STATUS_MESSAGES *)

```

```

(* Update Date string *)
settextjustify(righttext, centertext);
DATE := DATE NOW;
if FORCE_UPDATE or (DATE <> PREVIOUS_DATE) then
begin
  setcolor(SCREEN.BACKGROUND_COLOR);
  outtextxy
  (
    MAX_X - 2 * EDGE_OFFSET,
    MAX_Y + MESSAGE_REGION div 3,
    PREVIOUS_DATE
  );
  setcolor(SCREEN.STATUS_COLOR);
  PREVIOUS_DATE := DATE;
  outtextxy
  (
    MAX_X - 2 * EDGE_OFFSET,
    MAX_Y + MESSAGE_REGION div 3,
    PREVIOUS_DATE
  );
end; (* If...then *)

(* Update Generation/Paused status *)
settextjustify(lefttext, centertext);
setcolor(SCREEN.BACKGROUND_COLOR);
if PREVIOUS_USER_PAUSED
then
  outtextxy
  (
    2 * EDGE_OFFSET,
    MAX_Y + 1 * MESSAGE_REGION div 3,
    'Paused (' + INT_STR(GENERATION) + ')'
  );
else if FORCE_UPDATE or (PREVIOUS_GENERATION <> GENERATION) then
  outtextxy
  (
    2 * EDGE_OFFSET,
    MAX_Y + 1 * MESSAGE_REGION div 3,
    'Generation ' + INT_STR(GENERATION)
  );
PREVIOUS_USER_PAUSED := USER_PAUSED;
setcolor(SCREEN.STATUS_COLOR);
if PREVIOUS_USER_PAUSED
then
  outtextxy
  (
    2 * EDGE_OFFSET,
    MAX_Y + 1 * MESSAGE_REGION div 3,
    'Paused (' + INT_STR(GENERATION) + ')'
  );
else if FORCE_UPDATE or (PREVIOUS_GENERATION <> GENERATION) then
  outtextxy
  (
    2 * EDGE_OFFSET,
    MAX_Y + 1 * MESSAGE_REGION div 3,
    'Generation ' + INT_STR(GENERATION)
  );
PREVIOUS_GENERATION := GENERATION;

if FORCE_UPDATE or (INDIVIDUAL <> PREVIOUS_INDIVIDUAL) then
begin
  settextjustify(lefttext, centertext);
  setcolor(SCREEN.BACKGROUND_COLOR);
  outtextxy
  (
    2 * EDGE_OFFSET,
    MAX_Y + 2 * MESSAGE_REGION div 3,
    'Individual ' + INT_STR(PREVIOUS_INDIVIDUAL) + '/' +
    INT_STR(POPULATION_SIZE)
  );
  PREVIOUS_INDIVIDUAL := INDIVIDUAL;
  setcolor(SCREEN.STATUS_COLOR);
  outtextxy
  (
    2 * EDGE_OFFSET,
    MAX_Y + 2 * MESSAGE_REGION div 3,
    'Individual ' + INT_STR(PREVIOUS_INDIVIDUAL) + '/' +
    INT_STR(POPULATION_SIZE)
  );
end; (* if...then *)
end; (* UPDATE_STATUS_MESSAGES *)

procedure SET_GRAPH_SCALES(var USER : USER_GRAPH_SCALES; GRAPH_NUMBER:integer);
(* Takes the user scales and moves them into the global graph scales. *)
begin (* SET_GRAPH_SCALES *)
  with USER_GRAPH(GRAPH_NUMBER).SCALES do
  begin
    case GRAPH_NUMBER of
      1 : begin
          FIRST := 1;
          LABT := POPULATION_SIZE;
        end;
      2 : begin
          FIRST := 0;
          LABT := 100;
        end;
      3 : begin
          FIRST := 0;
          LABT := 100;
        end;
    end; (* case...of *)
    if USER.SCALES.HIGH_LABEL(X_AXIS) > USER.SCALES.LOW_LABEL(X_AXIS)
    then
      begin
        HIGH_LABEL(X_AXIS) := USER.SCALES.HIGH_LABEL(X_AXIS);
        LOW_LABEL(X_AXIS) := USER.SCALES.LOW_LABEL(X_AXIS);
      end (* if...then *)
    else
      write(BELL);
      NUMBER_OF_TICK_MARKS(X_AXIS) :=
        USER.SCALES.NUMBER_OF_TICK_MARKS(X_AXIS);
      if USER.SCALES.HIGH_LABEL(Y_AXIS) > USER.SCALES.LOW_LABEL(Y_AXIS)
      then
        begin
          HIGH_LABEL(Y_AXIS) := USER.SCALES.HIGH_LABEL(Y_AXIS);
          LOW_LABEL(Y_AXIS) := USER.SCALES.LOW_LABEL(Y_AXIS);
        end (* if...then *)
      else
        write(BELL);
        NUMBER_OF_TICK_MARKS(Y_AXIS) :=
          USER.SCALES.NUMBER_OF_TICK_MARKS(Y_AXIS);
        with USER_GRAPH(GRAPH_NUMBER) do
        begin
          PEAK_DETECTION(X_AXIS) := USER.PEAK_DETECTION(X_AXIS);
          PEAK_DETECTION(Y_AXIS) := USER.PEAK_DETECTION(Y_AXIS);
        end; (* with...do *)
      end; (* if...then *)
    end; (* SET_GRAPH_SCALES *)
  end;
end;

procedure INITIALIZE;
(* This procedure parses the command-line parameters, initializes all *)
(* global data structures, reads the INI file. *)
var
  (* Loop Control Variable for parsing the command-line parameters. *)
  LCV, GRAPH_NUMBER : integer;

  (* The command-line parameter begin parsed. *)
  PARAM : string;

  (* Used to detect invalid command-line parameters so that the 'help' *)
  (* info can be displayed. *)
  INVALID_PARAM : boolean;

  function UP_STRING(S : string) : string;
  (* This routine returns a string corresponding to the upper case value *)
  (* of the string passed to it. *)
  var LCV : integer;
  begin (* UP_STRING *)
    UP_STRING[0] := S[0];
    for LCV := 1 to length(S) do
      UP_STRING(LCV) := upcase(S(LCV));
    end; (* UP_STRING *)
  end;

  procedure PROCESS_INI_FILE;
  (* This procedure finds and reads the INI file. All "valid" INI file *)
  (* lines are processed (invalid lines are silently ignored). *)
  const
    (* The following are used to parse the INI file lines of text. *)
    CONFIG_LABEL = 'CONFIG';
    NORMAL_COLORS_LABEL = 'NORMALCOLORS';
    REVERSE_COLORS_LABEL = 'REVERSECOLORS';
    EVALUATIONS_LABEL = 'EVALUATIONS';
    FITNESS_HISTOGRAM_LABEL = 'FITNESSHISTOGRAM';
    DIVERSITY_HISTOGRAM_LABEL = 'DIVERSITYHISTOGRAM';
    COMMENT_DELIMITER = ';';
    TRUE_FLAG = 'TRUE';
    FALSE_FLAG = 'FALSE';
    IMAGE_DIRECTORY_CONTROL = 'IMAGEDIRECTORY';
    VIDEO_MODE_CONTROL = 'VIDEOMODE';
    VGA_MODE_FLAG = 'VGA';
    SVGA_MODE_FLAG = 'SVGA';
    VESA_MODE_FLAG = 'VESA';
    SVESA_MODE_FLAG = 'SVESA';
    BACKGROUND_COLOR_CONTROL = 'BACKGROUNDCOLOR';
    AXIS_COLOR_CONTROL = 'AXISCOLOR';
    GRID_COLOR_CONTROL = 'GRIDCOLOR';
    AXIS_LABEL_COLOR_CONTROL = 'AXISLABELCOLOR';
    GRAPH_COLOR_CONTROL = 'GRAPHCOLOR';
    STATUS_COLOR_CONTROL = 'STATUSCOLOR';
    STATUS_HIGHLIGHT_COLOR_CONTROL = 'STATUSHIGHLIGHTCOLOR';
    MAIN_TITLE_COLOR_CONTROL = 'MAINTITLECOLOR';
    BORDER_COLOR_CONTROL = 'BORDERCOLOR';
    MAX_X_CONTROL = 'MAXX';
    MIN_X_CONTROL = 'MINX';
    X_TICK_MARKS_CONTROL = 'XTICKMARKS';
    Y_PEAK_DETECTION_CONTROL = 'YPEAKDETECTION';
    MAX_Y_CONTROL = 'MAXY';
    MIN_Y_CONTROL = 'MINY';
    Y_TICK_MARKS_CONTROL = 'YTIICKMARKS';
    Y_PEAK_DETECTION_CONTROL = 'YPEAKDETECTION';
    Y_AVERAGE_DETECTION_CONTROL = 'YAVERAGEDETECTION';

    (* State variable that defines the section of the INI file that *)
    (* is being parsed. *)
    PARSING :
      UNDEFINED,
      CONFIG_SECTION,
      COLOR_SECTION,
      REVERSE_SECTION,
      EVALUATIONS_SECTION,
      FITNESS_HISTOGRAM_SECTION,
      DIVERSITY_HISTOGRAM_SECTION
      ) = UNDEFINED;

  var
    (* INI file parsing variables. *)
    INI_FILENAME : string(sizeof(dirst)+sizeof(namestr)+sizeof(extstr));
    LINE : string;
    INI_FILE : text;
    POSITION : integer;
    CONTROL_CHAR : char;

    procedure BOOLEAN_VALUE(CONTROL : string; var PARAMETER : boolean);
    (* This procedure examines the global LINE searching for CONTROL, if *)
    (* it is found, it's boolean value is set based on the rest of the LINE. *)
    begin (* BOOLEAN_VALUE *)
      if UP_STRING(copy(LINE, 1, POSITION - 1)) = CONTROL then
        begin
          delete(LINE, 1, POSITION);
          if UP_STRING(copy(LINE, 1, length(TRUE_FLAG))) = TRUE_FLAG then
            PARAMETER := true;
          else if UP_STRING(copy(LINE, 1, length(FALSE_FLAG))) = FALSE_FLAG
          then
            PARAMETER := false;
          end; (* if...then *)
        end; (* BOOLEAN_VALUE *)
      end;

    procedure INTEGER_VALUE(CONTROL : string; var PARAMETER : longint);
    (* This procedure examines the global LINE searching for CONTROL, if *)
    (* it is found, it's integer value is set based on the rest of the LINE. *)
    var
      (* Used to convert text to integers. *)
      VALUE, CODE : integer;
    begin (* INTEGER_VALUE *)
      if UP_STRING(copy(LINE, 1, POSITION - 1)) = CONTROL then
        begin
          delete(LINE, 1, POSITION);
          POSITION := 1;
          while LINE(POSITION) in ('0'..'9') do inc(POSITION);
          val(copy(LINE, 1, POSITION - 1), VALUE, CODE);
          if CODE = 0 then PARAMETER := VALUE;
          end; (* if...then *)
        end; (* INTEGER_VALUE *)
      end;

    procedure DIRECTORY_VALUE(CONTROL : string; var PARAMETER : dirstr);
    (* This procedure examines the global LINE searching for CONTROL, if *)

```

```

(* it is found then it's string value is set based on the rest of the *)
(* LINE. *)
(***** *)
var (***** *)
(* Temporary variables used to strip characters from the pathname. *)
(***** *)
DIRECTORY dirstr; FILENAME : namestr; EXTENSION : extstr;

begin (* DIRECTORY_VALUE *)
if UP_STRING(copy(LINE, 1, POSITION - 1)) = CONTROL then
begin
PARAMETER := copy(LINE, POSITION + 1, length(LINE));
if pos(SPACE, PARAMETER) < 0 then
PARAMETER(pos(SPACE, PARAMETER)) := '\';
if pos('\', PARAMETER) < 0 then
PARAMETER(0) := chr(pos('\', PARAMETER));
fsplit(PARAMETER, DIRECTORY, FILENAME, EXTENSION);
PARAMETER := UP_STRING(DIRECTORY);
end; (* if...then *)
end; (* DIRECTORY_VALUE *)

procedure VIDEO_VALUE(CONTROL : string; var PARAMETER : VIDEO_MODES);
(***** *)
(* This procedure examines the global LINE searching for CONTROL, if *)
(* it is found then it's string value is examined to see if it matches *)
(* one of the valid video modes. *)
(***** *)
begin (* VIDEO_VALUE *)
if UP_STRING(copy(LINE, 1, POSITION - 1)) = CONTROL then
begin
delete(LINE, 1, POSITION);
while LINE[length(LINE)] = SPACE do delete(LINE, length(LINE), 1);
POSITION := 1;
if UP_STRING(LINE) = VGA_MODE_FLAG
then
PARAMETER := VGA
else if UP_STRING(LINE) = SVGA_MODE_FLAG
then
PARAMETER := SVGA
else if UP_STRING(LINE) = VESA_MODE_FLAG
then
PARAMETER := VESA
else if UP_STRING(LINE) = SVESA_MODE_FLAG
then
PARAMETER := SVESA;
end; (* if...then *)
end; (* VIDEO_VALUE *)

procedure COLOR_VALUE(CONTROL : string; var PARAMETER : word);
(***** *)
(* This procedure examines the global LINE searching for CONTROL, if *)
(* it is found, it's string value is examined to see if it matches one *)
(* of the valid colors. *)
(***** *)
var (***** *)
(* Temporary variable to hold the color value, needed because it *)
(* may not be valid. *)
(***** *)
TEMP_COLOR : byte;

function COLOR_MATCH(PARAMETER : string; var COLOR : byte) : boolean;
(***** *)
(* This routine performs an associative match on the PARAMETER to *)
(* the list of possible color choices. Any enumeration of the letters *)
(* of a color will match as long as enough letters are provided to *)
(* determine which color is desired. For example: matched to 'blue' *)
(* include blu, ble, bue & be. *)
(***** *)
const (***** *)
(* Strings to match the colors to. *)
(***** *)
COLOR_STRING : array (0..15) of string =
(
'BLACK', 'BLUE', 'GREEN', 'CYAN', 'RED', 'MAGENTA', 'BROWN',
'LIGHTGRAY', 'DARKGRAY', 'LIGHTBLUE', 'LIGHTGREEN',
'LIGHTCYAN', 'LIGHTRED', 'LIGHTMAGENTA', 'YELLOW', 'WHITE'
);
var (***** *)
(* Index into parameter characters. *)
(***** *)
PARAM_LCV : integer;

(***** *)
(* Flag to indicate that the character being searched for was *)
(* found or not. *)
(***** *)
FOUND : boolean;

(***** *)
(* Local copy of the string being searched. *)
(***** *)
LOCAL_PARAMETER : string;

begin (* COLOR_MATCH *)
COLOR := 0;
repeat
LOCAL_PARAMETER := PARAMETER;
FOUND := pos(LOCAL_PARAMETER(1), COLOR_STRING(COLOR)) = 1;
PARAM_LCV := 1;
while FOUND and (PARAM_LCV <= length(PARAMETER)) do
begin
FOUND := FOUND and (pos(LOCAL_PARAMETER(PARAM_LCV),
COLOR_STRING(COLOR)) < 0);
inc(PARAM_LCV);
end; (* while...do *)
inc(COLOR);
until FOUND or (COLOR > 15);
if not FOUND
then
COLOR := $FF
else
dec(COLOR);
COLOR_MATCH := FOUND;
end; (* COLOR_MATCH *)

begin (* COLOR_VALUE *)
if UP_STRING(copy(LINE, 1, POSITION - 1)) = CONTROL then
begin
delete(LINE, 1, POSITION);
while LINE[length(LINE)] = SPACE do delete(LINE, length(LINE), 1);
POSITION := 1;
if COLOR_MATCH(UP_STRING(LINE), TEMP_COLOR) then
PARAMETER := TEMP_COLOR;
end; (* if...then *)
end; (* COLOR_VALUE *)

begin (* PROCESS_INI_FILE *)
INI_FILENAME := paramstr(0);
delete(INI_FILENAME, length(INI_FILENAME) - 2, 3);
INI_FILENAME := INI_FILENAME & INI;
assign(INI_FILE, INI_FILENAME);
(*$I*) reset(INI_FILE); (*$I+*)
if ioresult = 0 then
begin
while not eof(INI_FILE) do
begin
readln(INI_FILE, LINE);
POSITION := pos(COMMENT_DELIMITER, LINE);
if POSITION < 0 then
delete(LINE, POSITION, length(LINE));
if UP_STRING(copy(LINE, 1, length(NORMAL_COLORS_LABEL))) =
NORMAL_COLORS_LABEL then
begin
PARSING := COLOR_SECTION;
readln(INI_FILE, LINE);
POSITION := pos(COMMENT_DELIMITER, LINE);
if POSITION < 0 then
delete(LINE, POSITION, length(LINE));
end (* if...then *)
else if UP_STRING(copy(LINE, 1, length(REVERSE_COLORS_LABEL))) =
REVERSE_COLORS_LABEL then
begin
PARSING := REVERSE_SECTION;
readln(INI_FILE, LINE);
POSITION := pos(COMMENT_DELIMITER, LINE);
if POSITION < 0 then
delete(LINE, POSITION, length(LINE));
end (* if...then *)
else if UP_STRING(copy(LINE, 1, length(CONFIG_LABEL))) =
CONFIG_LABEL then
begin
PARSING := CONFIG_SECTION;
readln(INI_FILE, LINE);
POSITION := pos(COMMENT_DELIMITER, LINE);
if POSITION < 0 then
delete(LINE, POSITION, length(LINE));
end (* if...then *)
else if UP_STRING(copy(LINE, 1, length(EVALUATIONS_LABEL))) =
EVALUATIONS_LABEL then
begin
PARSING := EVALUATIONS_SECTION;
readln(INI_FILE, LINE);
POSITION := pos(COMMENT_DELIMITER, LINE);
if POSITION < 0 then
delete(LINE, POSITION, length(LINE));
end (* if...then *)
else if UP_STRING(copy(LINE, 1,
length(FITNESS_HISTOGRAM_LABEL))) =
FITNESS_HISTOGRAM_LABEL then
begin
PARSING := FITNESS_HISTOGRAM_SECTION;
readln(INI_FILE, LINE);
POSITION := pos(COMMENT_DELIMITER, LINE);
if POSITION < 0 then
delete(LINE, POSITION, length(LINE));
end (* if...then *)
else if UP_STRING(copy(LINE, 1,
length(DIVERSITY_HISTOGRAM_LABEL))) =
DIVERSITY_HISTOGRAM_LABEL then
begin
PARSING := DIVERSITY_HISTOGRAM_SECTION;
readln(INI_FILE, LINE);
POSITION := pos(COMMENT_DELIMITER, LINE);
if POSITION < 0 then
delete(LINE, POSITION, length(LINE));
end; (* if...then *)
POSITION := pos('=', LINE);
if POSITION < 0 then
case PARSING of
CONFIG_SECTION
begin
DIRECTORY_VALUE(IMAGE_DIRECTORY_CONTROL, IMAGE_DIRECTORY);
if DISPLAY_MODE = NO_DISPLAY_MODE then
VIDEO_VALUE(VIDEO_MODE_CONTROL, DISPLAY_MODE);
end; (* Configuration Section *)
COLOR_SECTION :
begin
COLOR_VALUE(BACKGROUND_COLOR_CONTROL,
NORMAL.BACKGROUND_COLOR);
COLOR_VALUE(AXIS_COLOR_CONTROL,
NORMAL.AXIS_COLOR);
COLOR_VALUE(GRID_COLOR_CONTROL,
NORMAL.GRID_COLOR);
COLOR_VALUE(AXIS_LABEL_COLOR_CONTROL,
NORMAL.AXIS_LABEL_COLOR);
COLOR_VALUE(GRAPH_COLOR_CONTROL,
NORMAL.GRAPH_COLOR);
COLOR_VALUE(STATUS_COLOR_CONTROL,
NORMAL.STATUS_COLOR);
COLOR_VALUE(STATUS_HIGHLIGHT_COLOR_CONTROL,
NORMAL.STATUS_HIGHLIGHT_COLOR);
COLOR_VALUE(MAIN_TITLE_COLOR_CONTROL,
NORMAL.MAIN_TITLE_COLOR);
COLOR_VALUE(BORDER_COLOR_CONTROL,
NORMAL.BORDER_COLOR);
end; (* Color Section *)
REVERSE_SECTION :
begin
COLOR_VALUE(BACKGROUND_COLOR_CONTROL,
REVERSE.BACKGROUND_COLOR);
COLOR_VALUE(AXIS_COLOR_CONTROL,
REVERSE.AXIS_COLOR);
COLOR_VALUE(GRID_COLOR_CONTROL,
REVERSE.GRID_COLOR);
COLOR_VALUE(AXIS_LABEL_COLOR_CONTROL,
REVERSE.AXIS_LABEL_COLOR);
COLOR_VALUE(GRAPH_COLOR_CONTROL,
REVERSE.GRAPH_COLOR);
COLOR_VALUE(STATUS_COLOR_CONTROL,
REVERSE.STATUS_COLOR);
COLOR_VALUE(STATUS_HIGHLIGHT_COLOR_CONTROL,
REVERSE.STATUS_HIGHLIGHT_COLOR);
COLOR_VALUE(MAIN_TITLE_COLOR_CONTROL,
REVERSE.MAIN_TITLE_COLOR);
COLOR_VALUE(BORDER_COLOR_CONTROL,
REVERSE.BORDER_COLOR);
end; (* Reverse Color Section *)
EVALUATIONS_SECTION :
with USER_GRAPH(1).SCALES do
begin
INTEGER_VALUE(MAX_X_CONTROL,
HIGH_LABEL(X_AXIS));
INTEGER_VALUE(MIN_X_CONTROL,
LOW_LABEL(X_AXIS));
INTEGER_VALUE(X_TICK_MARKS_CONTROL,
NUMBER_OF_TICK_MARKS(X_AXIS));
BOOLEAN_VALUE(Y_PEAK_DETECTION_CONTROL,
PEAK_DETECTION(X_AXIS));
INTEGER_VALUE(MAX_Y_CONTROL,
HIGH_LABEL(Y_AXIS));
INTEGER_VALUE(MIN_Y_CONTROL,
LOW_LABEL(Y_AXIS));
INTEGER_VALUE(Y_TICK_MARKS_CONTROL,
NUMBER_OF_TICK_MARKS(Y_AXIS));
BOOLEAN_VALUE(Y_PEAK_DETECTION_CONTROL,
PEAK_DETECTION(Y_AXIS));
BOOLEAN_VALUE(Y_AVERAGE_DETECTION_CONTROL,
AVERAGE_DETECTION);
end; (* with...do *)
FITNESS_HISTOGRAM_SECTION :
with USER_GRAPH(2).SCALES do
begin

```

```

INTEGER_VALUE(MAX_X_CONTROL, HIGH_LABEL[X_AXIS]);
INTEGER_VALUE(MIN_X_CONTROL, LOW_LABEL[X_AXIS]);
INTEGER_VALUE(X_TICK_MARKS_CONTROL, NUMBER_OF_TICK_MARKS[X_AXIS]);
BOOLEAN_VALUE(X_PEAK_DETECTION_CONTROL, PEAK_DETECTION[X_AXIS]);
INTEGER_VALUE(MAX_Y_CONTROL, HIGH_LABEL[Y_AXIS]);
INTEGER_VALUE(MIN_Y_CONTROL, LOW_LABEL[Y_AXIS]);
INTEGER_VALUE(Y_TICK_MARKS_CONTROL, NUMBER_OF_TICK_MARKS[Y_AXIS]);
BOOLEAN_VALUE(Y_PEAK_DETECTION_CONTROL, PEAK_DETECTION[Y_AXIS]);
BOOLEAN_VALUE(Y_AVERAGE_DETECTION_CONTROL, AVERAGE_DETECTION);
end; (* with...do *)
DIVERSITY_HISTOGRAM_SECTION :
with USER_GRAPH[3].SCALES do
begin
  INTEGER_VALUE(MAX_X_CONTROL, HIGH_LABEL[X_AXIS]);
  INTEGER_VALUE(MIN_X_CONTROL, LOW_LABEL[X_AXIS]);
  INTEGER_VALUE(X_TICK_MARKS_CONTROL, NUMBER_OF_TICK_MARKS[X_AXIS]);
  BOOLEAN_VALUE(X_PEAK_DETECTION_CONTROL, PEAK_DETECTION[X_AXIS]);
  INTEGER_VALUE(MAX_Y_CONTROL, HIGH_LABEL[Y_AXIS]);
  INTEGER_VALUE(MIN_Y_CONTROL, LOW_LABEL[Y_AXIS]);
  INTEGER_VALUE(Y_TICK_MARKS_CONTROL, NUMBER_OF_TICK_MARKS[Y_AXIS]);
  BOOLEAN_VALUE(Y_PEAK_DETECTION_CONTROL, PEAK_DETECTION[Y_AXIS]);
  BOOLEAN_VALUE(Y_AVERAGE_DETECTION_CONTROL, AVERAGE_DETECTION);
end; (* with...do *)
else (* Do Nothing! *)
end; (* case...of *)
end; (* while...do *)
close(INI_FILE);
end; (* if...then *)
end; (* PROCESS_INI_FILE *)

begin (* INITIALIZE *)
(* Initialize 'constants' *)
directvideo := false;

(* Parse Command-Line parameters *)
INVALID_PARAM := false;
for LCV := 1 to paramcount do
begin
  PARAM := UP_STRING(paramstr(LCV));
  case PARAM[2] of
    'R' : if PARAM = '-REVERSE'
      then
        REVERSE_VIDEO := true
      else
        INVALID_PARAM := true;
    'V' : if PARAM = '-VGA' then
      DISPLAY_MODE := VGA
    else if PARAM = '-VESA' then
      DISPLAY_MODE := VESA
    else
      INVALID_PARAM := true;
    'S' : if PARAM = '-SVGA' then
      DISPLAY_MODE := SVGA
    else if PARAM = '-SVESA' then
      DISPLAY_MODE := SVESA
    else
      INVALID_PARAM := true;
  else INVALID_PARAM := true;
end; (* case...of *)
end; (* for...to...do *)
if INVALID_PARAM then
begin
  writeln(LF,
    ' Evolution of Solutions to Real-Time Problems', CR, LF,
    LF,
    ' Usage:', CR, LF,
    LF,
    (*$IFDEF MATH_GA *)
    ' GMATH-GA [-VGA|-SVGA|-VESA|-SVESA] [-Reverse]', CR, LF,
    (*$ELSE *)
    ' GRTOS-GA [-VGA|-SVGA|-VESA|-SVESA] [-Reverse]', CR, LF,
    (*$ENDIF *)
    LF,
    ' Where:', CR, LF,
    ' -VGA = 640x480 graphics mode', CR, LF,
    ' -VGA = 800x600 graphics mode', CR, LF,
    ' -VESA = 1024x768 graphics mode (default)', CR, LF,
    ' -SVESA = 1280x1024 graphics mode', CR, LF,
    ' -Reverse = Reverse Video Color Scheme', CR, LF
  );
halt(INVALID_PARAMETER_ERROR_LEVEL);
end; (* if...then *)
GRAPH_OFFSET := 6 * EDOE_OFFSET;
with USER_GRAPH[1].SCALES do
begin
  TITLE := 'Fitness';
  GRAPH_TYPE := LINE_GRAPH;
  LABEL_NAME[X_AXIS] := 'Individual';
  HIGH_LABEL[X_AXIS] := ('Configuration #') POPULATION_SIZE;
  LOW_LABEL[X_AXIS] := ('Configuration #') 0;
  NUMBER_OF_TICK_MARKS[X_AXIS] := 10;
  PEAK_DETECTION[X_AXIS] := false;
  LABEL_NAME[Y_AXIS] := 'Fitness';
  HIGH_LABEL[Y_AXIS] := ('Score') 100;
  LOW_LABEL[Y_AXIS] := ('Score') 0;
  NUMBER_OF_TICK_MARKS[Y_AXIS] := 10;
  PEAK_DETECTION[Y_AXIS] := false;
  AVERAGE_DETECTION := false;
end; (* with...do *)
with USER_GRAPH[2].SCALES do
begin
  TITLE := 'Fitness Histogram';
  GRAPH_TYPE := BAR_GRAPH;
  LABEL_NAME[X_AXIS] := 'Fitness';
  HIGH_LABEL[X_AXIS] := ('Score') 100;
  LOW_LABEL[X_AXIS] := ('Score') 0;
  NUMBER_OF_TICK_MARKS[X_AXIS] := 20;
  PEAK_DETECTION[X_AXIS] := false;
  LABEL_NAME[Y_AXIS] := 'Fitness';
  HIGH_LABEL[Y_AXIS] := ('Percent') 100;
  LOW_LABEL[Y_AXIS] := ('Percent') 0;
  NUMBER_OF_TICK_MARKS[Y_AXIS] := 10;
  PEAK_DETECTION[Y_AXIS] := true;
  AVERAGE_DETECTION := false;
end; (* with...do *)
with USER_GRAPH[3].SCALES do
begin
  TITLE := 'Diversity Histogram';
  GRAPH_TYPE := BAR_GRAPH;
  LABEL_NAME[X_AXIS] := 'Diversity';
  HIGH_LABEL[X_AXIS] := ('Configuration #') 100;
  LOW_LABEL[X_AXIS] := ('Configuration #') 0;
  NUMBER_OF_TICK_MARKS[X_AXIS] := 20;
  PEAK_DETECTION[X_AXIS] := false;
  LABEL_NAME[Y_AXIS] := 'Diversity';
  HIGH_LABEL[Y_AXIS] := ('Percent') 100;
  LOW_LABEL[Y_AXIS] := ('Percent') 0;
  NUMBER_OF_TICK_MARKS[Y_AXIS] := 10;
  PEAK_DETECTION[Y_AXIS] := true;
  AVERAGE_DETECTION := false;
end; (* with...do *)
PROCESS_INI_FILE;
if REVERSE_VIDEO
then
  SCREEN := REVERSE
else
  SCREEN := NORMAL;
(* Handle the case when no INI file (or parameter) was found *)
if DISPLAY_MODE = NO_DISPLAY_MODE then
  DISPLAY_MODE := VESA;
(* Allocate Plot data structures *)
for GRAPH_NUMBER := 1 to MAX_NUMBER_OF_GRAPHS do
begin
  new(PLOT[GRAPH_NUMBER].GRAPH_POINTS);
  for LCV := 0 to POPULATION_SIZE do
  with PLOT[GRAPH_NUMBER].GRAPH_POINTS[LCV] do
  begin
    POSITION[X_AXIS] := 0.0;
    POSITION[Y_AXIS] := 0.0;
  end; (* for...to...do *)
  new(PLOT[GRAPH_NUMBER].PREVIOUS_GRAPH_POINTS);
  for LCV := 0 to POPULATION_SIZE do
  with PLOT[GRAPH_NUMBER].PREVIOUS_GRAPH_POINTS[LCV] do
  begin
    POSITION[X_AXIS] := 0.0;
    POSITION[Y_AXIS] := 0.0;
  end; (* for...to...do *)
  SET_GRAPH_SCALES(USER_GRAPH[GRAPH_NUMBER], GRAPH_NUMBER);
end; (* for...to...do *)
end; (* INITIALIZE *)

procedure INIT_GRAPHICS;
(* This procedure initializes the video graphics card based on the video mode that has been selected. *)
var
  (* Graphic Video Driver parameters. *)
  DRIVER : integer;
begin (* INIT_GRAPHICS *)
if DISPLAY_MODE = NO_DISPLAY_MODE then DISPLAY_MODE := VESA;
case DISPLAY_MODE of
  VGA : begin
    registerbdriver(addr(EGAVGA_DRIVER));
    MESSAGE_REGION := 24;
  end; (* VGA *)
  SVGA : begin
    installuserdriver('VESAL6', @DETECT_VESA_L6);
    registerbdriver(addr(VESAL6_DRIVER));
    MESSAGE_REGION := 36;
    VESA_L6_MODE := 0;
  end; (* SVGA *)
  VESA : begin
    installuserdriver('VESAL6', @DETECT_VESA_L6);
    registerbdriver(addr(VESAL6_DRIVER));
    MESSAGE_REGION := 36;
    VESA_L6_MODE := 1;
  end; (* VESA *)
  SVESA : begin
    installuserdriver('VESAL6', @DETECT_VESA_L6);
    registerbdriver(addr(VESAL6_DRIVER));
    MESSAGE_REGION := 36;
    VESA_L6_MODE := 2;
  end; (* SVESA *)
end; (* case...of *)
DRIVER := detect;
initgraph(DRIVER, GRAPHICS_MODE, '');
if graphresult <> GOK then
begin
  writeln(CR, LF, 'Fatal Error: Video mode not supported');
  halt(NON_SUPPORTED_VIDEO_MODE_ERROR_LEVEL);
end; (* if...then *)
end; (* INIT_GRAPHICS *)

procedure SHUTDOWN_GRAPHICS; far;
(* Restores the video controller to standard DOS text mode. *)
begin (* SHUTDOWN_GRAPHICS *)
closegraph;
end; (* SHUTDOWN_GRAPHICS *)

procedure REDRAW_GRAPH(NUMBER_OF_GRAPHS : integer);
(* Redraws the graph axis or displays the signon logo if there are no active graphs. *)
var
  (* Graphic screen position of the next status text message. *)
  STATUS_POS : integer;
procedure HIGHLIGHTED_TEXT(STR : string);
(* Displays a text string in the status message box where the first letter of the string is highlighted. *)
var
  (* Loop Control Variable for indexing through string. *)
  LCV : integer;
begin (* HIGHLIGHTED_TEXT *)
setcolor(SCREEN.STATUS_HIGHLIGHT_COLOR);
outtextxy
(
  STATUS_POS,
  MAX_Y + MESSAGE_REGION div 3,
  STR[1]
);
STATUS_POS := STATUS_POS + textwidth(STR[1]);
setcolor(SCREEN.STATUS_COLOR);
outtextxy
(
  STATUS_POS,
  MAX_Y + MESSAGE_REGION div 3,
  copy(STR, 2, length(STR))
);
end; (* HIGHLIGHTED_TEXT *)
end; (* REDRAW_GRAPH *)

```

```

);
for LCV := 2 to length(STR) do
  STATUS_POS := STATUS_POS + textwidth(STR[LCV]);
STATUS_POS := STATUS_POS + 2 * textwidth('h');
end; (* HIGHLIGHTED_TEXT *)

procedure DRAW_RIT_LOGO(X, Y : integer);
(* Draws the 'RIT' logo on the screen at the specified coordinates. *)
(*-----*)
const
  (* Pointer to a buffer that contains the drawn logo (valid only *)
  (* after the first call). This allows the logo to be displayed *)
  (* more than once but only drawn once. *)
  LOGO : pointer := nil;
var
  (*-----*)
  (* Work area of the screen coordinates. *)
  WORK_X, WORK_Y : integer;
begin
  (* DRAW_RIT_LOGO *)
  X := X - 225;
  if LOGO = nil then
    begin
      (* Create Logo *)
      setfillstyle(solidfill, red);
      bar(X - 70, Y + 60, X + 520, Y + 60);
      setcolor(white);
      rectangle(X - 65, Y - 55, X + 515, Y + 55);
      setfillstyle(solidfill, white);
      floodfill(X, Y - 10, white);
      setcolor(red);
      settextstyle(TriplexFont, vertdir, 5);
      settextjustify(centertext, centertext);
      outtextxy(X + 40, Y - 1, 'R I T');
      setfillstyle(solidfill, red);
      circle(X - 32, Y + 17, 2); floodfill(X - 32, Y + 17, red);
      circle(X - 32, Y - 17, 2); floodfill(X - 32, Y - 17, red);
      setfillstyle(solidfill, lightgray);
      bar(X - 10, Y - 50, X + 5, Y + 50);
      settextstyle(SimplexFont, horizdir, 0);
      settextjustify(lefttext, centertext);
      setusercharsize(7, 8, 3, 4);
      outtextxy(X + 40, Y + 5, 'ROCHESTER INSTITUTE OF TECHNOLOGY');
      setlinestyle(solidln, 0, thickwidth);
      setcolor(lightgray);
      line(X + 40, Y + 30, X + 470, Y + 30);

      (* Keep a copy of the logo so it doesn't need to redrawn *)
      getmem(LOGO, imagesize(X - 70, Y - 60, X + 520, Y + 60));
      getimage(X, Y - 60, X + 520, Y + 60, LOGO);
      end (* if...then *)
    else
      putimage(X - 70, Y - 60, LOGO, normalput);
    end; (* DRAW_RIT_LOGO *)
end;

procedure DRAW_COPYRIGHT;
(*-----*)
(* Draws the 'circles c' copyright notice. *)
(*-----*)
begin
  (* DRAW_COPYRIGHT *)
  outtext('c');
  circle
  (
    getx,
    gety + round(0.3 * textheight('c')),
    round(textheight('c') / 2.5)
  );
end; (* DRAW_COPYRIGHT *)

procedure CREATE_GRAPH
(*-----*)
(
  AXIS_TYPE : AXIS_TYPES;
  GRAPH_POS : GRAPH_POSITIONS;
  GRAPH_NUMBER : integer;
  SCALE_INFO : GRAPH_SCALE_INFO
);
(*-----*)
(* Creates a graph by allocating memory for the graph data structures *)
(* and initializing them as directed by the various parameters. *)
(*-----*)
var
  (*-----*)
  (* Loop Control Variable for initializing screen position values. *)
  LCV : integer;
begin
  (* CREATE_GRAPH *)
  with PLOT[GRAPH_NUMBER] do
    begin
      GRAPH_SCALES := SCALE_INFO;
      VALUES.FIRST := GRAPH_SCALES.FIRST;
      VALUES.LAST := GRAPH_SCALES.LAST;
      with GRAPH_SCALES do
        begin
          RANGE[X_AXIS] := abs(HIGH_LABEL[X_AXIS] - LOW_LABEL[X_AXIS]);
          RANGE[Y_AXIS] := abs(HIGH_LABEL[Y_AXIS] - LOW_LABEL[Y_AXIS]);
          end; (* with...do *)
          GRAPH_COLOR := SCREEN.GRAPH_COLOR;
          PREVIOUS_PEAK := getmaxx + 50;
          PREVIOUS_PEAK_VALUE := getmaxy + 50;
          PREVIOUS_AVERAGE_VALUE := getmaxy + 50;
          HORIZONTAL_OFFSET := 2 * EDGE_OFFSET + GRAPH_OFFSET;
          HORIZONTAL_SIZE := MAX_X - 2 * (GRAPH_OFFSET + EDGE_OFFSET);
          VERTICAL_OFFSET := 2 * EDGE_OFFSET + GRAPH_OFFSET;
          VERTICAL_SIZE := MAX_Y - 2 * (GRAPH_OFFSET + EDGE_OFFSET);
        end;
        case GRAPH_POS of
          FULL_SCREEN : (* Do Nothing *);
          TOP_HALF :
            begin
              VERTICAL_SIZE := VERTICAL_SIZE / 2 - VERTICAL_OFFSET;
            end;
          BOTTOM_HALF :
            begin
              VERTICAL_SIZE := VERTICAL_SIZE / 2 - VERTICAL_OFFSET;
              VERTICAL_OFFSET := VERTICAL_SIZE + 2 * VERTICAL_OFFSET;
            end;
          TOP_THIRD :
            begin
              VERTICAL_SIZE := VERTICAL_SIZE / 3 - 1.5 * VERTICAL_OFFSET;
            end;
          MIDDLE_THIRD :
            begin
              VERTICAL_SIZE := VERTICAL_SIZE / 3 - 1.5 * VERTICAL_OFFSET;
              VERTICAL_OFFSET := VERTICAL_SIZE + 3 * VERTICAL_OFFSET;
            end;
          BOTTOM_THIRD :
            begin
              VERTICAL_SIZE := VERTICAL_SIZE / 3 - 1.5 * VERTICAL_OFFSET;
              VERTICAL_OFFSET := 2 * VERTICAL_SIZE + 5 * VERTICAL_OFFSET;
            end;
        end; (* case...of *)
      end; (* with PLOT[GRAPH_NUMBER] do *)
    end; (* CREATE_GRAPH *)
  end;
end; (* case...of *)

with GRAPH_SCALES do
  begin
    (* Determine position for X Axis (Low, Centered, High) *)
    if ((LOW_LABEL[Y_AXIS] >= 0) and (HIGH_LABEL[Y_AXIS] >= 0) and
        (AXIS_TYPE = AUTO_AXIS)) or
        (AXIS_TYPE in {FORCE_LL, FORCE_LR})
    then
      X_AXIS_OFFSET := VERTICAL_OFFSET + VERTICAL_SIZE
    else if ((LOW_LABEL[X_AXIS] <= 0) and (HIGH_LABEL[X_AXIS] <= 0) and
        (AXIS_TYPE = AUTO_AXIS)) or
        (AXIS_TYPE in {FORCE_UL, FORCE_UR})
    then
      X_AXIS_OFFSET := VERTICAL_OFFSET
    else
      X_AXIS_OFFSET := VERTICAL_OFFSET + VERTICAL_SIZE / 2;
    (* Determine position for Y Axis (Left, Centered, Right) *)
    if ((LOW_LABEL[X_AXIS] >= 0) and (HIGH_LABEL[X_AXIS] >= 0) and
        (AXIS_TYPE = AUTO_AXIS)) or
        (AXIS_TYPE in {FORCE_LL, FORCE_UL})
    then
      Y_AXIS_OFFSET := HORIZONTAL_OFFSET
    else if ((LOW_LABEL[X_AXIS] <= 0) and (HIGH_LABEL[X_AXIS] <= 0) and
        (AXIS_TYPE in {FORCE_LR, FORCE_UR})
    then
      Y_AXIS_OFFSET := HORIZONTAL_OFFSET + HORIZONTAL_SIZE
    else
      Y_AXIS_OFFSET := HORIZONTAL_OFFSET + HORIZONTAL_SIZE / 2;
    end; (* with...do *)
  end;
end; (* CREATE_GRAPH *)

for LCV := 0 to POPULATION_SIZE do
  with PLOT[GRAPH_NUMBER].PREVIOUS_GRAPH_POINTS[LCV] do
    begin
      POSITION[X_AXIS] := HORIZONTAL_OFFSET;
      POSITION[Y_AXIS] := VERTICAL_OFFSET;
      end; (* with...do *)
    end; (* with...do *)
  end; (* CREATE_GRAPH *)

procedure DRAW_GRAPH_AXIS(GRAPH_NUMBER : integer);
(*-----*)
(* Draws the title, axis, grid lines, and tick marks for the graph. *)
(*-----*)
var
  (*-----*)
  (* Loop Control Variable for creating tick marks. *)
  LCV : integer;
  (*-----*)
  (* Value of axis tick mark graphic position. *)
  TEMP : longint;
  (*-----*)
  (* Value of the tick mark label. *)
  TICK_LABEL : real;
begin
  (* DRAW_GRAPH_AXIS *)
  with PLOT[GRAPH_NUMBER] do
    begin
      (* Add Graph Title *)
      setcolor(SCREEN.MAIN_TITLE_COLOR);
      settextjustify(centertext, centertext);
      case DISPLAY_MODE of
        VGA : settextstyle(TRIPLEX_FONT, horizdir, 1);
        SVGA : settextstyle(TRIPLEX_FONT, horizdir, 2);
        VESA :
          settextstyle(TRIPLEX_FONT, horizdir, 3);
      end; (* case...of *)
      outtextxy
      (
        MAX_X div 2,
        round(VERTICAL_OFFSET - GRAPH_OFFSET),
        GRAPH_SCALES.TITLE
      );
      (* Add X Tick Marks *)
      setusercharsize(1, 3, 1, 3);
      settextstyle(simplex_font, horizdir, 0);
      settextjustify(centertext, toptext);
      with GRAPH_SCALES do
        if NUMBER_OF_TICK_MARKS[X_AXIS] < 0 then
          for LCV := 0 to NUMBER_OF_TICK_MARKS[X_AXIS] do
            begin
              TEMP := round(HORIZONTAL_OFFSET +
                HORIZONTAL_SIZE /
                NUMBER_OF_TICK_MARKS[X_AXIS] * LCV);
              setcolor(SCREEN.AXIS_COLOR);
              setlinestyle(solidln, 0, normwidth);
              line
              (
                TEMP, round(X_AXIS_OFFSET),
                TEMP, round(X_AXIS_OFFSET + TICK_LENGTH)
              );
              setcolor(SCREEN.GRID_COLOR);
              setlinestyle(dottedln, 0, normwidth);
              line
              (
                TEMP, round(VERTICAL_OFFSET),
                TEMP, round(X_AXIS_OFFSET)
              );
              TICK_LABEL := RANGE[X_AXIS] /
                NUMBER_OF_TICK_MARKS[X_AXIS] * LCV +
                LOW_LABEL[X_AXIS];
              if TICK_LABEL = round(TICK_LABEL) then
                begin
                  setcolor(SCREEN.AXIS_COLOR);
                  outtextxy
                  (
                    TEMP,
                    round(X_AXIS_OFFSET + TICK_LENGTH - 1),
                    INT_STR[round(TICK_LABEL)]
                  );
                end; (* if...then *)
              end; (* for...to...do *)
            end;
          end; (* with GRAPH_SCALES do *)
          (* Add Y Tick Marks *)
          settextjustify(righttext, centertext);
          with GRAPH_SCALES do
            if NUMBER_OF_TICK_MARKS[Y_AXIS] < 0 then
              for LCV := 0 to NUMBER_OF_TICK_MARKS[Y_AXIS] do
                begin
                  TEMP := round(VERTICAL_OFFSET +
                    VERTICAL_SIZE /
                    NUMBER_OF_TICK_MARKS[Y_AXIS] * LCV);
                  setcolor(SCREEN.AXIS_COLOR);
                  setlinestyle(solidln, 0, normwidth);
                  line
                  (
                    round(Y_AXIS_OFFSET - TICK_LENGTH), TEMP,
                    round(Y_AXIS_OFFSET), TEMP
                  );
                  setcolor(SCREEN.GRID_COLOR);
                  setlinestyle(dottedln, 0, normwidth);
                  line
                  (
                    round(Y_AXIS_OFFSET), TEMP,
                    MAX_X - EDGE_OFFSET - GRAPH_OFFSET, TEMP
                  );
                end;
              end;
            end;
          end;
        end;
      end;
    end;
  end;
end; (* DRAW_GRAPH_AXIS *)

```

```

);
TICK_LABEL := RANGE(Y_AXIS) /
              NUMBER_OF_TICK_MARKS(Y_AXIS) *
              (NUMBER_OF_TICK_MARKS(Y_AXIS) - LCV) +
              LOW_LABEL(Y_AXIS);
if TICK_LABEL = round(TICK_LABEL) then
begin
  setcolor(SCREEN.AXIS_COLOR);
  outtextxy
  (
    round(Y_AXIS_OFFSET - TICK_LENGTH),
    TEMP - 1,
    INT_STR(round(TICK_LABEL))
  );
end; (* if...then *)
end; (* with...do *)
) * Draw X Axis *)
setcolor(SCREEN.AXIS_COLOR);
setlinestyle(solidln, 0, normwidth);
line
(
  round(HORIZONTAL_OFFSET),
  round(X_AXIS_OFFSET),
  round(HORIZONTAL_OFFSET + HORIZONTAL_SIZE),
  round(X_AXIS_OFFSET)
);
) * Draw Y Axis *)
setcolor(SCREEN.AXIS_COLOR);
setlinestyle(solidln, 0, normwidth);
line
(
  round(Y_AXIS_OFFSET),
  round(VERTICAL_OFFSET),
  round(Y_AXIS_OFFSET),
  round(VERTICAL_OFFSET + VERTICAL_SIZE)
);
with GRAPH_SCALES do
begin
  ) * Label Axis *)
  setcolor(SCREEN.AXIS_LABEL_COLOR);
  setusercharsize(1, 2, 1, 2);
  settextrjustif(righttext, centertext);
  outtextxy
  (
    round(HORIZONTAL_OFFSET + HORIZONTAL_SIZE),
    round(X_AXIS_OFFSET + round(0.75 * textheight('M'))),
    LABEL_NAME(X_AXIS)
  );
  settextrjustif(lefttext, centertext);
  outtextxy
  (
    round(Y_AXIS_OFFSET - textwidth('M')),
    round(VERTICAL_OFFSET - round(0.75 * textheight('M'))),
    LABEL_NAME(Y_AXIS)
  );
end; (* with...do *)
end; (* with...do *)
setlinestyle(solidln, 0, normwidth);
end; (* DRAW_GRAPH_AXIS *)

procedure DRAW_BOX(UL_X, UL_Y, LR_X, LR_Y : integer; COLOR, LINE_WIDTH:word);
(* Draws a colored, unfilled, solid line rectangle at the specified *)
(* coordinates. *)
begin (* DRAW_BOX *)
  setlinestyle(solidln, 0, LINE_WIDTH);
  setcolor(COLOR);
  rectangle(UL_X, UL_Y, LR_X, LR_Y);
end; ) * DRAW_BOX *)

begin (* REDRAW_GRAPH *)
  cleardevice;
  DRAW_BOX(0, 0, getmaxx, getmaxy, SCREEN.BACKGROUND_COLOR, normwidth);
  setfillstyle(solidfill, SCREEN.BACKGROUND_COLOR);
  floodfill(1, 1, SCREEN.BACKGROUND_COLOR);
  MAX_X := getmaxx;
  MAX_Y := getmaxy;

  DRAW_BOX
  (
    0, 0,
    MAX_X, MAX_Y,
    SCREEN.BORDER_COLOR,
    normwidth
  );
  MAX_Y := MAX_Y - MESSAGE_REGION;
  DRAW_BOX
  (
    0, MAX_Y,
    MAX_X, MAX_Y + MESSAGE_REGION,
    SCREEN.BORDER_COLOR,
    normwidth
  );
  settextrjustif(lefttext, centertext);
  case DISPLAY_MODE of
    VGA : begin
      setusercharsize(2, 5, 2, 5);
      settextrstyle(SIMPLEX_FONT, horizdir, 0);
      STATUS_POS := MAX_X div 4;
      end; (* VGA *)
    SVGA : begin
      setusercharsize(1, 2, 1, 2);
      settextrstyle(SIMPLEX_FONT, horizdir, 0);
      STATUS_POS := MAX_X div 4;
      end; (* SVGA *)
    VESA,
    SVESA : begin
      settextrstyle(SIMPLEX_FONT, horizdir, 2);
      STATUS_POS := MAX_X div 4;
      end; (* VESA|SVESA *)
  end; (* case...of *)
  HIGHLIGHTED TEXT ('Evolve');
  HIGHLIGHTED TEXT ('Pause');
  HIGHLIGHTED TEXT ('Grab image');
  HIGHLIGHTED TEXT ('Reverse colors');
  HIGHLIGHTED TEXT ('Make report');
  HIGHLIGHTED TEXT ('Quit');

  (* Always display two graphs *)
  CREATE_GRAPH(FORCE_LL, TOP_THIRD, 1, USER_GRAPH(1).SCALES);
  CREATE_GRAPH(FORCE_LL, MIDDLE_THIRD, 2, USER_GRAPH(2).SCALES);
  CREATE_GRAPH(FORCE_LL, BOTTOM_THIRD, 3, USER_GRAPH(3).SCALES);

  case NUMBER_OF_GRAPHS of
    0 : begin
      setcolor(SCREEN.MAIN_TITLE_COLOR);
      case DISPLAY_MODE of
        VGA,
        SVGA : settextrstyle(TRIPLEX_FONT, horizdir, 4);
        VESA,
        SVESA : settextrstyle(TRIPLEX_FONT, horizdir, 5);
      end; (* case...of *)
      settextrjustif(centertext, centertext);
      move to
      (
        MAX_X div 2,
        MAX_Y div 2 + 0 * textheight('H')
      );
      outtext('Evolution of Solutions to');
      move over(0, round(1.5 * textheight('H')));
      outtext('Real-Time Problems');
      case DISPLAY_MODE of
        VGA,
        SVGA : settextrstyle(TRIPLEX_FONT, horizdir, 3);
        VESA,
        SVESA : settextrstyle(TRIPLEX_FONT, horizdir, 4);
      end; (* case...of *)
      move over(0, 4 * textheight('H'));
      outtext('Greg P. Semeraro');
      move over(0, textheight('H'));
      case DISPLAY_MODE of
        VGA,
        SVGA : settextrstyle(SIMPLEX_FONT, horizdir, 2);
        VESA,
        SVESA : settextrstyle(SIMPLEX_FONT, horizdir, 3);
      end; (* case...of *)
      move over(1 * textwidth('H'), 0);
      outtext('Copyright 1997');
      move over(-8 * textwidth('H'), 0);
      DRAW_COPYRIGHT;
      DRAW_RIT_LOGO(MAX_X div 2, (MAX_Y div 4));
      end; (* No Graphs *)
      1 : begin
        DRAW_GRAPH_AXIS(1);
        end; (* One, Full-screen graph *)
      2 : begin
        DRAW_GRAPH_AXIS(1);
        DRAW_GRAPH_AXIS(2);
        end; (* Two Graphs *)
      3 : begin
        DRAW_GRAPH_AXIS(1);
        DRAW_GRAPH_AXIS(2);
        DRAW_GRAPH_AXIS(3);
        end; (* Three graphs *)
      end; (* case...of *)
      settextrstyle(defaultfont, horizdir, 1);
      UPDATE STATUS_MESSAGES(FORCE_UPDATE);
      end; (* REDRAW_GRAPH *)

  procedure UPDATE_GRAPHS(NUMBER_OF_GRAPHS : integer);
  (* This procedure updates then graphs on the screen by getting the data *)
  (* from the Sound Blaster DMA buffer and then drawing the graph. *)
  const
  (* Flag that indicates which half of the DMA buffer contains the *)
  (* present data samples. *)
  FRONT_HALF : boolean = false;

  (* Static data to allow previous values to be erased from the screen *)
  PREVIOUS_GENERATION : longint = 0;
  PREVIOUS_STD_DEV : real = 0.0;
  PREVIOUS_DIVERSITY_STD_DEV : real = 0.0;
  PREVIOUS_BEST_INDIVIDUAL_FITNESS : real = 0.0;
  PREVIOUS_ALL_TIME_BEST_INDIVIDUAL_FITNESS : real = 0.0;

  var
  (* Index into the data values arrays. *)
  X_LCV, LAST_ELEMENT : integer;

  (* Keeps track of the number of DMA interrupts that occurred since *)
  (* the last one was serviced. *)
  NUMBER_OF_INTERRUPTS : integer;

  (* Loop Control Variable to update all graphs being displayed. *)
  PLOT_NUMBER : integer;

  procedure DRAW_LINE_GRAPH(GRAPH_NUMBER : integer);
  (* Draws an entire line graph using small line segments. If previous line *)
  (* is erased if necessary (depending on the graph type selected). *)
  var
  (* Used to index into the data value array. *)
  X : integer;

  (* Used to hold the X-Axis value (converts integer to real). *)
  ACTUAL_X_VALUE : real;

  (* Used to calculate the position of the Y-Axis peak/average. *)
  PEAK_VALUE, AVERAGE_VALUE : real; PEAK_DETECTED : boolean;

  (* Used to hold the X-Axis position of the Y-Axis peak. *)
  PEAK : real;

  (* Used to swap the present and previous lines. *)
  TEMP_PTR : pointer;

  (* Used to calculate the average value of each sweep of the graph. *)
  NUMBER_IN_AVERAGE : integer;

  (* Used to hold the string value of the 'real' peak-to-average. *)
  TEMP_STR : string[20];

  (* Markers for the first and last X values that fit on the screen *)
  (* (needed when the X-Axis limits are changed from the default). *)
  FIRST_X, LAST_X : integer;
  X_VALUE_OFF_SCREEN : boolean;

  procedure PUT_X_BLOCK(COLOR : word; X_POS : real; GRAPH_NUMBER : integer);
  (* Draws a '█' along the X-Axis corresponding to the X position. *)
  var
  (* Graphic screen position of X-Axis. *)
  Y_POS : real;

```





```

MAX_X - 2 * EDGE_OFFSET,
round(PLOT(2).VERTICAL_OFFSET) - 2 * textheight('H'),
'std. Dev.' + REAL_STR(PREVIOUS_STD_DEV, 2, 2)
);
end; (* if...then *)
if GENETIC_ALGORITHM_FITNESS_CONVERGED
then
setcolor(SCREEN.AXIS_COLOR)
else
setcolor(SCREEN.BACKGROUND_COLOR);
outtextxy
(
MAX_X - 2 * EDGE_OFFSET - 3 * textwidth('H'),
round(PLOT(2).VERTICAL_OFFSET) - 1 * textheight('H'),
'Converged'
);
);
if LONG_TERM_DIVERSITY.STANDARD_DEVIATION <>
PREVIOUS_DIVERSITY_STD_DEV then
begin
setcolor(SCREEN.BACKGROUND_COLOR);
outtextxy
(
MAX_X - 2 * EDGE_OFFSET,
round(PLOT(3).VERTICAL_OFFSET) - 2 * textheight('H'),
'std. Dev.' + REAL_STR(PREVIOUS_DIVERSITY_STD_DEV, 2, 2)
);
);
setcolor(SCREEN.AXIS_COLOR);
PREVIOUS_DIVERSITY_STD_DEV := LONG_TERM_DIVERSITY.
STANDARD_DEVIATION;
outtextxy
(
MAX_X - 2 * EDGE_OFFSET,
round(PLOT(3).VERTICAL_OFFSET) - 2 * textheight('H'),
'std. Dev.' + REAL_STR(PREVIOUS_DIVERSITY_STD_DEV, 2, 2)
);
);
end; (* if...then *)
if GENETIC_ALGORITHM_DIVERSITY_CONVERGED
then
setcolor(SCREEN.AXIS_COLOR)
else
setcolor(SCREEN.BACKGROUND_COLOR);
outtextxy
(
MAX_X - 2 * EDGE_OFFSET - 3 * textwidth('H'),
round(PLOT(3).VERTICAL_OFFSET) - 1 * textheight('H'),
'Converged'
);
);
(* Update Most Fit Individual *)
if PRESENT_GENERATION.BEST <> PREVIOUS_BEST_INDIVIDUAL_FITNESS then
begin
setcolor(SCREEN.BACKGROUND_COLOR);
outtextxy
(
MAX_X - 2 * EDGE_OFFSET,
round(PLOT(1).VERTICAL_OFFSET) - 2 * textheight('H'),
'Present Best' +
REAL_STR(PREVIOUS_BEST_INDIVIDUAL_FITNESS, 4, 2)
);
);
setcolor(SCREEN.AXIS_COLOR);
PREVIOUS_BEST_INDIVIDUAL_FITNESS := PRESENT_GENERATION.BEST;
outtextxy
(
MAX_X - 2 * EDGE_OFFSET,
round(PLOT(1).VERTICAL_OFFSET) - 2 * textheight('H'),
'Present Best' +
REAL_STR(PREVIOUS_BEST_INDIVIDUAL_FITNESS, 4, 2)
);
);
end; (* if...then *)
(* Update Most Fit Individual *)
if LONG_TERM.BEST <> PREVIOUS_ALL_TIME_BEST_INDIVIDUAL_FITNESS then
begin
setcolor(SCREEN.BACKGROUND_COLOR);
outtextxy
(
MAX_X - 2 * EDGE_OFFSET,
round(PLOT(1).VERTICAL_OFFSET) 1 * textheight('H'),
'All-Time Best' +
REAL_STR(PREVIOUS_ALL_TIME_BEST_INDIVIDUAL_FITNESS, 2, 2)
);
);
setcolor(SCREEN.AXIS_COLOR);
PREVIOUS_ALL_TIME_BEST_INDIVIDUAL_FITNESS := LONG_TERM.BEST;
outtextxy
(
MAX_X - 2 * EDGE_OFFSET,
round(PLOT(1).VERTICAL_OFFSET) - 1 * textheight('H'),
'All-Time Best' +
REAL_STR(PREVIOUS_ALL_TIME_BEST_INDIVIDUAL_FITNESS, 4, 2)
);
);
end; (* if...then *)
end; (* if...then *)
end; (* if...then *)
if NUMBER_OF_GRAPHS <> 0 then
begin
with PLOT(1) do
for X_LCV := GRAPH_SCALES.FIRST to GRAPH_SCALES.LAST do
VALUES.DATA(X_LCV) := POPULATION(X_LCV).FITNESS;
);
with PLOT(2) do
for X_LCV := GRAPH_SCALES.FIRST to GRAPH_SCALES.LAST do
VALUES.DATA(X_LCV) := LONG_TERM.HISTOGRAM(X_LCV);
);
with PLOT(3) do
for X_LCV := GRAPH_SCALES.FIRST to GRAPH_SCALES.LAST do
VALUES.DATA(X_LCV) := LONG_TERM_DIVERSITY.HISTOGRAM(X_LCV);
);
DRAW_LINE_GRAPH(1);
DRAW_LINE_GRAPH(2);
DRAW_LINE_GRAPH(3);
end; (* if...then *)
end; (* UPDATE_GRAPHS *)
function IMAGE_FILENAME : string;
{*****}
(* This routine determines the next available IMAGE xx.BMP filename. *)
{*****}
var
{*****}
(* Used to assign a unique filename to the image file. *)
{*****}
IMAGE_NUMBER : integer;
{*****}
(* Result of I/O operation, used to find first available filename. *)
{*****}
RESULT : integer;
{*****}
(* Filename of file to get image. *)
{*****}
FILENAME : string(sizeof(namestr) + sizeof(dirstr) + 1);
{*****}
(* File used to find first available file. *)
{*****}
IMAGE : file;
begin (* IMAGE_FILENAME *)
IMAGE_NUMBER := 0;
repeat
FILENAME := 'IMAGE ';
if IMAGE_NUMBER < 10 then
FILENAME := FILENAME + '0';
FILENAME := FILENAME + INT_STR(IMAGE_NUMBER) + '.BMP';
inc(IMAGE_NUMBER);
assign(IMAGE, IMAGE_DIRECTORY + FILENAME);
(*$I+*) reset(IMAGE); (*$I+*)
RESULT := ioread;
if RESULT = 0 then
close(IMAGE);
until RESULT = 2;
IMAGE_FILENAME := FILENAME;
end; (* IMAGE_FILENAME *)
procedure RANDM_KEYBOARD_EVENT;
{*****}
(* This routine processes all user keyboard input by performing the action *)
(* requested (if possible). *)
{*****}
var
{*****}
(* Character to hold user input command. *)
{*****}
USER_INPUT : char;
begin (* HANDLE_KEYBOARD_EVENT *)
USER_INPUT := upcase(readkey);
case USER_INPUT of
'E' : begin
if NUMBER_OF_GRAPHS_TO_DISPLAY = 0
then
NUMBER_OF_GRAPHS_TO_DISPLAY := 3
else
NUMBER_OF_GRAPHS_TO_DISPLAY := 0;
REDRAW_GRAPH(NUMBER_OF_GRAPHS_TO_DISPLAY);
end; (* Run Evolution *)
'G' : begin
UPDATE_STATUS_MESSAGES(FORCE_UPDATE);
SAVE_IMAGE_AS_16_COLOR_BMP_FILE
(
IMAGE_DIRECTORY + IMAGE_FILENAME,
0, 0, getmaxx, getmaxy
);
UPDATE_STATUS_MESSAGES(FORCE_UPDATE);
end; (* Grab Image *)
'Q' : USER_QUIT := true;
'P' : begin
USER_PAUSED := not USER_PAUSED;
UPDATE_STATUS_MESSAGES(FORCE_UPDATE);
end; (* Pause *)
'M' : begin
USER_PAUSED := true;
PRODUCE_GENETIC_ALGORITHM_REPORT;
UPDATE_STATUS_MESSAGES(FORCE_UPDATE);
end; (* Pause *)
'R' : begin
REVERSE_VIDEO := not REVERSE_VIDEO;
if REVERSE_VIDEO
then
SCREEN := REVERSE
else
SCREEN := NORMAL;
REDRAW_GRAPH(NUMBER_OF_GRAPHS_TO_DISPLAY);
end; (* Reverse Video Colors *)
else write(BELL);
end; (* case...of *)
end; (* RANDM_KEYBOARD_EVENT *)
begin (* GENETIC_ALGORITHM_THESIS *)
(* Install exit procedure *)
SAVE_EXIT := exitproc;
exitproc := @SHUTDOWN_GRAPHICS;
(* Initialization *)
INITIALIZE;
INIT_GRAPHICS;
REDRAW_GRAPH(NUMBER_OF_GRAPHS_TO_DISPLAY);
(* Go until user wants to exit *)
repeat
if keypressed then HANDLE_KEYBOARD_EVENT;
if not USER_QUIT and not USER_PAUSED and not keypressed and
not (GENETIC_ALGORITHM_FITNESS_CONVERGED and
GENETIC_ALGORITHM_DIVERSITY_CONVERGED) then
begin
if NUMBER_OF_GRAPHS_TO_DISPLAY <> 0 then
repeat
UPDATE_STATUS_MESSAGES(DO_NOT_FORCE_UPDATE);
until (PROCESS_GENETIC_ALGORITHM_INDIVIDUAL = POPULATION_SIZE);
UPDATE_GRAPHS(NUMBER_OF_GRAPHS_TO_DISPLAY);
UPDATE_STATUS_MESSAGES(DO_NOT_FORCE_UPDATE);
end; (* While...do *)
until USER_QUIT or (GENERATION = maxlongint);
if GENERATION <> 0 then
PRODUCE_GENETIC_ALGORITHM_REPORT;
end; (* GENETIC_ALGORITHM_THESIS *)

```

## 9.3 RTOS\_GA.PAS

```
(* workload implies higher priority. *)
(*.....*)
WORKLOAD_MONOTONIC_ASSIGNMENT);
(*.....*)

(* Problem specific parameter. This structure defines the 'genes' of
 * the individuals that will make up the population. *)
GENOTYPE = record
(*.....*)
(* Cooperative or Preemptive Multitasking. *)
TASKING_MODEL : TASKING_MODELS;
(*.....*)
(* Preemptive multitasking timeslice (microseconds). *)
TARGET_TIMESLICE : longint;
(*.....*)
(* Enable/Disable Priority Inheritance *)
PRIORITY_INHERITANCE_ENABLED : boolean;
(*.....*)
(* Static or Rotating priorities. *)
PRIORITY_ALLOCATION : PRIORITY_SCHEDULING_POLICIES;
(*.....*)
(* Uniform, random, rate monotonic, deadline monotonic, or workload *)
(* Monotonic Priority assignments *)
INITIAL_PRIORITY_ASSIGNMENT : PRIORITY_ASSIGNMENT_ALGORITHMS;
(*.....*)
(* Genetic Algorithm bookkeeping variables used to determine which *)
(* individual will 'live' in the next generation (or to a *)
(* The individual's fitness evaluation parameters. *)
FITNESS, RELATIVE_FITNESS, CUMULATIVE_FITNESS : real;
(*.....*)
(* The individual's diversity parameter (there is a 1:1 mapping *)
(* of genotype values and diversity value). *)
DIVERSITY : longint;
(*.....*)
(* Flag to indicate whether the individual will 'live' in the *)
(* next generation. *)
SURVIVOR : boolean;
(*.....*)
(* Seed for random number generator so that results can be *)
(* reproduced. *)
RANDOM_NUMBER_SEED : longint;
end; (* GENOTYPE *)

(* This data structure is used to determine how well the GA is doing. *)
(* The only criteria used is convergence to a single solution (or to a *)
(* single 'fitness' value). The GA is assumed to be done when the *)
(* population has converged regardless of the actual fitness value *)
(* achieved. *)
STATISTICS = record
(*.....*)
(* Basic values for the population... *)
BEST, MEAN, WORST : real;
(*.....*)
(* Statistics based on the entire population... *)
VARIANCE, STANDARD_DEVIATION : real;
(*.....*)
(* Normalized histogram buckets. *)
HISTOGRAM : array [round(MINIMUM_FITNESS - 0.5) ..
round(MAXIMUM_FITNESS + 0.5)] of real;
end; (* STATISTICS *)

const
(* This is the current generation that has completed processing. *)
GENERATION : longint = 0;
(* This is the current individual that has complete processing. *)
INDIVIDUAL : integer = POPULATION_SIZE;
(* This is the maximum value for the genotype diversity, it is used *)
(* to scale the diversity display. *)
MAX_NORMALIZED_DIVERSITY : real = 0.0;
(* Indicates the status of the GA, can be used to terminate the GA. *)
GENETIC_ALGORITHM_FITNESS_CONVERGED : boolean = false;
GENETIC_ALGORITHM_DIVERSITY_CONVERGED : boolean = false;

var
(* Statistics are available for each generation that is processed, *)
(* the data is independent from generation to generation. *)
PRESENT_GENERATION : STATISTICS;

(* Cumulative statistics are available for the most recent *)
(* generations (see FITNESS_MAX_AVERAGE_GENERATIONS). *)
LONG_TERM : STATISTICS;

(* Diversity statistics for the present generation. *)
PRESENT_DIVERSITY : STATISTICS;

(* Cumulative diversity statistics for the most recent generations *)
(* (see DIVERSITY_MAX_AVERAGE_GENERATIONS). *)
LONG_TERM_DIVERSITY : STATISTICS;

(* These are all of the genes of the population. *)
POPULATION : array [0..POPULATION_SIZE] of GENOTYPE;

(* Used to verify the existence of the TASKING application program. *)
EXECUTABLE_FILE : file;

function PROCESS_GENETIC_ALGORITHM_INDIVIDUAL : integer;
(* This routine performs all processing on the individual (i.e., Selection, *)
(* Crossover, Mutation and Evaluation). The individual that was processed *)
(* returned by this function. *)

```

```

procedure PROCESS_GENETIC_ALGORITHM_GENERATION;
[*****]
[* This routine performs all processing on the population for an entire *]
[* generation [i.e. Selection/Crossover, Mutation and Evaluation]. *]
[*****]

procedure PRODUCE_GENETIC_ALGORITHM_REPORT;
[*****]
[* This routine produces a report on the 'best' individuals found. *]
[*****]

implementation

uses dos;

var
[*****]
[* List of all time best individuals that will be reported. *]
[*****]
ALL_TIME_BEST : array [0..(NUMBER_OF_BEST_TO_REPORT-1)] of GENOTYPE;

[*****]
[* The genes are allowed to range within these bounds so that each *]
[* individual created is guaranteed to have genes that are reasonable *]
[* for the problem [i.e. all individuals are viable]. *]
[*****]
UPPER_BOUND, LOWER_BOUND, RANGE : array [RTOS_GENES] of real;

const
[*****]
[* The index of the lowest in the list is maintained so that the list *]
[* can be kept as a circular buffer very efficiently. *]
[*****]
LOWEST_IN_LIST_OF_BEST_INDIVIDUALS : integer := 0;

[*****]
[* Flag that indicates that at least one fitness value was outside of *]
[* the allowable range [0.0 through 100.0], used in report. *]
[*****]
FITNESS_OUT_OF_RANGE : boolean := false;

[*****]
[* Text string of gene names [for INI file processing and report *]
[* generating] *]
[*****]
GENE_NAME : array [RTOS_GENES] of string =
[
  'TaskingModel',
  'TargetTimeslice',
  'PriorityInheritance',
  'PriorityAllocation',
  'InitialPriorityAssignment'
];

[*****]
[* Text strings for tasking model gene [for INI file processing and *]
[* report generating]. *]
[*****]
TASKING_MODEL_NAME : array [TASKING_MODELS] of string =
[
  'Cooperative',
  'Preemptive'
];

[*****]
[* Text strings for priority inheritance gene [for INI file processing *]
[* and report generating]. *]
[*****]
PRIORITY_INHERITANCE_NAME : array [boolean] of string =
[
  'Disabled',
  'Enabled'
];

[*****]
[* Text strings for priority scheduling policy gene [for INI file *]
[* processing and report generating]. *]
[*****]
PRIORITY_ALLOCATION_NAME :
array [PRIORITY_SCHEDULING_POLICIES] of string =
[
  'Static',
  'Rotating'
];

[*****]
[* Text strings for priority assignment algorithm gene [for INI file *]
[* processing and report generating]. *]
[*****]
INITIAL_PRIORITY_ASSIGNMENT_NAME :
array [PRIORITY_ASSIGNMENT_ALGORITHMS] of string =
[
  'UniformAssignment',
  'RandomAssignment',
  'RateMonotonicAssignment',
  'DeadlineMonotonicAssignment',
  'WorkloadMonotonicAssignment'
];

function RANDOM_VALUE [LOWER, UPPER : real] : real;
[*****]
[* Produces a random floating point number between LOWER and UPPER. *]
[*****]

begin
  [* RANDOM_VALUE *]
  RANDOM_VALUE := [random]maxint / maxint * [UPPER - LOWER] + LOWER;
end;

procedure WRITE_TO_INI_FILE [var INI_FILE : text; var GENES : GENOTYPE];
[*****]
[* This routine takes the individual's genotype and creates a TASKING INI *]
[* file that corresponds to it. The TASKING program will then use these *]
[* parameters to evaluate the effectiveness of the RTOS for the application. *]
[*****]

begin
  [* WRITE TO INI FILE *]
  writeln (INI_FILE, 'Tasking');
  writeln (INI_FILE, 'Statistics=taskStatistics');
  with GENES do
  begin
    RANDOM_NUMBER_SEED := [longint]random[SFFFF] shl 16 or
    writeln (INI_FILE, GENE_NAME[TASKING_MODEL_GENE], '=',
      TASKING_MODEL_NAME[TASKING_MOEOL]);
    writeln (INI_FILE, GENE_NAME[TARGET_TIMESLICE_GENE], '=',
      TARGET_TIMESLICE);
    writeln (INI_FILE,
      GENE_NAME[PRIORITY_INHERITANCE_GENE], '=',
      PRIORITY_INHERITANCE_NAME[PRIORITY_INHERITANCE_ENABLED]);
    writeln (INI_FILE,
      GENE_NAME[PRIORITY_ALLOCATION_GENE], '=',
      PRIORITY_ALLOCATION_NAME[PRIORITY_ALLOCATION]);
    writeln (INI_FILE);
    writeln (INI_FILE, 'Application');
    writeln (INI_FILE,
      GENE_NAME[INITIAL_PRIORITY_ASSIGNMENT_GENE], '=',
      INITIAL_PRIORITY_ASSIGNMENT_NAME[INITIAL_PRIORITY_ASSIGNMENT]);
    writeln (INI_FILE, 'RandomNumberSeed', RANDOM_NUMBER_SEED);
  end;
  [* with...do *]
end;
  [* WRITE TO INI FILE *]

function PROCESS_GENETIC_ALGORITHM_INDIVIDUAL : integer;
[*****]
[* This routine performs all processing on the individual [i.e., Selection, *]
[* Crossover, Mutation and Evaluation]. The individual that was processed *]
[* returned by this function. *]
[*****]

procedure RANDOM_GENE [var GENES : GENOTYPE; GENE : RTOS_GENES];
[*****]
[* Replaces the gene of the genotype with a new, randomly selected value. *]
[*****]

begin
  [* RANDOM_GENE *]
  with GENES do
  case GENE of
    TASKING_MODEL_GENE :
      TASKING_MODEL :=
        TASKING_MODELS
          round
            random
              RANDOM_VALUE
                LOWER_BOUND[TASKING_MODEL_GENE],
                UPPER_BOUND[TASKING_MODEL_GENE]
          ;
    TARGET_TIMESLICE_GENE :
      TARGET_TIMESLICE :=
        round
          random
            RANDOM_VALUE
              LOWER_BOUND[TARGET_TIMESLICE_GENE],
              UPPER_BOUND[TARGET_TIMESLICE_GENE]
          ;
    PRIORITY_INHERITANCE_GENE :
      PRIORITY_INHERITANCE_ENABLED :=
        boolean
          round
            random
              RANDOM_VALUE
                LOWER_BOUND[PRIORITY_INHERITANCE_GENE],
                UPPER_BOUND[PRIORITY_INHERITANCE_GENE]
          ;
    PRIORITY_ALLOCATION_GENE :
      PRIORITY_ALLOCATION :=
        PRIORITY_SCHEDULING_POLICIES
          round
            random
              RANDOM_VALUE
                LOWER_BOUND[PRIORITY_ALLOCATION_GENE],
                UPPER_BOUND[PRIORITY_ALLOCATION_GENE]
          ;
    INITIAL_PRIORITY_ASSIGNMENT_GENE :
      INITIAL_PRIORITY_ASSIGNMENT :=
        PRIORITY_ASSIGNMENT_ALGORITHMS
          round
            random
              RANDOM_VALUE
                LOWER_BOUND[INITIAL_PRIORITY_ASSIGNMENT_GENE],
                UPPER_BOUND[INITIAL_PRIORITY_ASSIGNMENT_GENE]
          ;
  end;
  [* case...of *]
end;
  [* RANDOM_GENE *]

procedure INITIALIZE_POPULATION;
[*****]
[* This routine initializes the GA by randomly initializing the genes of *]
[* the population and initializes the long term statistics. *]
[*****]

var
[*****]
[* Loop control variables. *]
[*****]
INDIVIDUAL_LCV, LCV : integer;
GENE : RTOS_GENES;

begin
  [* INITIALIZE POPULATION *]
  for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
  with POPULATION[INDIVIDUAL_LCV] do
  begin
    LOWER_BOUND[TASKING_MODEL_GENE] := integer[low]TASKING_MOEOL;
    UPPER_BOUND[TASKING_MODEL_GENE] := integer[high]TASKING_MOEOL;
    RANGE[TASKING_MODEL_GENE] :=
      UPPER_BOUND[TASKING_MODEL_GENE] -
      LOWER_BOUND[TASKING_MODEL_GENE];
    LOWER_BOUND[TARGET_TIMESLICE_GENE] := MINIMUM_TIMESLICE;
    UPPER_BOUND[TARGET_TIMESLICE_GENE] := SFFFF;
    RANGE[TARGET_TIMESLICE_GENE] :=
      UPPER_BOUND[TARGET_TIMESLICE_GENE] -
      LOWER_BOUND[TARGET_TIMESLICE_GENE];
    LOWER_BOUND[PRIORITY_INHERITANCE_GENE] :=
      integer[low]PRIORITY_INHERITANCE_ENABLED;
    UPPER_BOUND[PRIORITY_INHERITANCE_GENE] :=
      integer[high]PRIORITY_INHERITANCE_ENABLED;
    RANGE[PRIORITY_INHERITANCE_GENE] :=
      UPPER_BOUND[PRIORITY_INHERITANCE_GENE] -
      LOWER_BOUND[PRIORITY_INHERITANCE_GENE];
    LOWER_BOUND[PRIORITY_ALLOCATION_GENE] :=
      integer[low]PRIORITY_ALLOCATION;
    UPPER_BOUND[PRIORITY_ALLOCATION_GENE] :=
      integer[high]PRIORITY_ALLOCATION;
    RANGE[PRIORITY_ALLOCATION_GENE] :=
      UPPER_BOUND[PRIORITY_ALLOCATION_GENE] -
      LOWER_BOUND[PRIORITY_ALLOCATION_GENE];
    LOWER_BOUND[INITIAL_PRIORITY_ASSIGNMENT_GENE] :=
      integer[low]INITIAL_PRIORITY_ASSIGNMENT;
    UPPER_BOUND[INITIAL_PRIORITY_ASSIGNMENT_GENE] :=
      integer[high]INITIAL_PRIORITY_ASSIGNMENT;
    RANGE[INITIAL_PRIORITY_ASSIGNMENT_GENE] :=
      UPPER_BOUND[INITIAL_PRIORITY_ASSIGNMENT_GENE] -
      LOWER_BOUND[INITIAL_PRIORITY_ASSIGNMENT_GENE];
  end;
  [* for...to...do *]
  [* Randomly initialize population *]
  for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
  for GENE := low[RTOS_GENES] to high[RTOS_GENES] do
    RANDOM_GENE [POPULATION[INDIVIDUAL_LCV], GENE];
  end;
end;

```

```

(* Initialize Statistics *)
with PRESENT_GENERATION do
  begin
    BEST := -maxlongint;
    WORST := maxlongint;
    MEAN := 0.0;
    VARIANCE := 0.0;
    STANDARD_DEVIATION := 0.0;
    for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
      HISTOGRAM[LCV] := 0.0;
    end; (* with...do *)
    LONG_TERM := PRESENT_GENERATION;
    PRESENT_DIVERSITY := PRESENT_GENERATION;
    LONG_TERM_DIVERSITY := PRESENT_GENERATION;
    for LCV := 0 to (NUMBER_OF_BEST_TO_REPORT-1) do
      ALL_TIME_BEST[LCV].FITNESS := -maxlongint;
    end; (* INITIALIZE_POPULATION *)
end;

procedure EVALUATE(THIS_INDIVIDUAL : integer);
(* This routine evaluates the fitness for a single individual of the *)
(* population. *)
var
  (* Text file to contain the RTOS application configuration. *)
  INI_FILE : text;
function TASKING_REPORT : real;
(* This routine analyzes the TASKING report file and produces a floating *)
(* point number from 0.0 to 100.0 that corresponds to how well TASKING *)
(* thinks the RTOS parameters tune TASKING to the application. *)
const
  (* The TASKING parameters that are measured are assigned *)
  (* proportions that are used to weight each parameter (of *)
  (* course the total must sum to 1.0). *)
  PERIODIC_FAULTS_PERCENTAGE = 0.95;
  BANDWIDTH_PERCENTAGE = 1.0 - PERIODIC_FAULTS_PERCENTAGE;
var
  (* The text file that corresponds to the TASKING report file. *)
  REPORT : text;
  (* Text string used to parse the TASKING report file. *)
  LINE : string;
  (* Parameters from the TASKING report used to evaluate the RTOS *)
  (* configuration. *)
  BANDWIDTH, PERIODIC_FAULTS : real;
  (* Used to convert strings to real numbers (indicates status). *)
  CODE : integer;
begin
  (* TASKING_REPORT *)
  assign(REPORT, 'TASKING.RPT');
  (*SI-*) reset(REPORT); (*SI+*)
  repeat
    readln(REPORT, LINE);
  until pos('Available CPU Bandwidth', LINE) < 0;
  delete(LINE, 1, pos(' ', LINE));
  while LINE[1] = ' ' do delete(LINE, 1, 1);
  delete(LINE, length(LINE) - 1, 2);
  val(LINE, BANDWIDTH, CODE);

  PERIODIC_FAULTS := 0.0;
  readln(REPORT, LINE);
  if pos('Periodic Event Faults', LINE) < 0 then
    begin
      delete(LINE, 1, pos(' ', LINE));
      while LINE[1] = ' ' do delete(LINE, 1, 1);
      delete(LINE, length(LINE) - 1, 2);
      val(LINE, PERIODIC_FAULTS, CODE);
    end; (* if...then *)

  TASKING_REPORT := PERIODIC_FAULTS_PERCENTAGE * (100 - PERIODIC_FAULTS)
    + BANDWIDTH_PERCENTAGE * BANDWIDTH;
  close(REPORT);
end; (* TASKING_REPORT *)

function APPLICATION_REPORT : real;
(* This routine analyzes the TASKING report file and produces a floating *)
(* point number from 0.0 to 100.0 that corresponds to how well TASKING *)
(* thinks the RTOS parameters tune TASKING to the application. *)
begin
  (* APPLICATION_REPORT *)
  APPLICATION_REPORT := dosexitcode;
end; (* APPLICATION_REPORT *)

begin
  (* EVALUATE *)
  assign(INI_FILE, PROGRAM_NAME + '.INI');
  (*SI-*) rewrite(INI_FILE); (*SI+*)
  WRITE TO INI_FILE [INI_FILE, POPULATION(THIS_INDIVIDUAL)];
  close(INI_FILE);

  exec(PROGRAM_NAME + '.EXE', 'No Report File');
  with POPULATION(THIS_INDIVIDUAL) do
    begin
      FITNESS := RTOS_EVALUATION_PERCENTAGE * TASKING_REPORT +
        APPLICATION_EVALUATION_PERCENTAGE * APPLICATION_REPORT;
      DIVERSITY := longint [TARGET_TIMESLICE and $000FFFF] +
        longint [TASKING_MODEL] shl 16 +
        longint [PRIORITY_INHERITANCE_ENABLED] shl 17 +
        longint [PRIORITY_ALLOCATION] shl 18 +
        longint [INITIAL_PRIORITY_ASSIGNMENT] shl 19;
    end; (* with...do *)
  end; (* EVALUATE *)

  (* For computing correlation. *)
  SUM_OF_SQUARES := real;
  (* Weight used to average for long term statistics. *)
  AVERAGE_WEIGHT := real;
  (* For determining where to put the fitness value (in histogram). *)
  BUCKET : integer;
begin
  (* COMPUTE_POPULATION_STATISTICS *)
  with PRESENT_GENERATION do
    begin
      BEST := -maxlongint;
      WORST := maxlongint;
      end; (* with...do *)
    with PRESENT_GENERATION do
      begin
        MEAN := 0.0;
        for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
          with POPULATION(INDIVIDUAL_LCV) do
            begin
              MEAN := MEAN + FITNESS;
              if FITNESS > PRESENT_GENERATION.BEST then
                PRESENT_GENERATION.BEST := FITNESS;
              if FITNESS > LONG_TERM.BEST then
                begin
                  LONG_TERM.BEST := FITNESS;
                  ALL_TIME_BEST[LOWEST_IN_LIST_OF_BEST_INDIVIDUALS] :=
                    POPULATION(INDIVIDUAL_LCV);
                  LOWEST_IN_LIST_OF_BEST_INDIVIDUALS :=
                    (LOWEST_IN_LIST_OF_BEST_INDIVIDUALS + 1) mod
                    NUMBER_OF_BEST_TO_REPORT;
                end; (* if...then *)
              if FITNESS < PRESENT_GENERATION.WORST then
                PRESENT_GENERATION.WORST := FITNESS;
              if FITNESS < LONG_TERM.WORST then
                LONG_TERM.WORST := FITNESS;
            end;
            MEAN := MEAN / POPULATION_SIZE;
            SUM_OF_SQUARES := 0.0;
            for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
              SUM_OF_SQUARES := SUM_OF_SQUARES +
                sqr(POPULATION(INDIVIDUAL_LCV).FITNESS - MEAN);
            VARIANCE := SUM_OF_SQUARES / POPULATION_SIZE;
            STANDARD_DEVIATION := sqrt(VARIANCE);
            for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do HISTOGRAM[LCV] := 0.0;
            for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
              begin
                BUCKET := round(POPULATION(INDIVIDUAL_LCV).FITNESS);
                if (BUCKET < LOWEST_BUCKET) or (HIGHEST_BUCKET < BUCKET) then
                  FITNESS_OUT_OF_RANGE := true;
                HISTOGRAM[BUCKET] := HISTOGRAM[BUCKET] + 1;
                end; (* for...to...do *)
              SUM := 0.0;
              for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
                SUM := SUM + HISTOGRAM[LCV];
              for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
                HISTOGRAM[LCV] := HISTOGRAM[LCV] /
                  (SUM / (HIGHEST_BUCKET - LOWEST_BUCKET + 1));
            end; (* with...do *)
          end; (* Compute long term statistics *)
          if GENERATION > FITNESS_MAX_AVERAGE_GENERATIONS
            then
              AVERAGE_WEIGHT := 1 / FITNESS_MAX_AVERAGE_GENERATIONS
            else
              AVERAGE_WEIGHT := 1 / GENERATION;
          with LONG_TERM do
            begin
              MEAN := MEAN * (1 - AVERAGE_WEIGHT) +
                PRESENT_GENERATION.MEAN * AVERAGE_WEIGHT;
              VARIANCE := VARIANCE * (1 - AVERAGE_WEIGHT) +
                PRESENT_GENERATION.VARIANCE * AVERAGE_WEIGHT;
              STANDARD_DEVIATION := STANDARD_DEVIATION * (1 - AVERAGE_WEIGHT) +
                PRESENT_GENERATION.STANDARD_DEVIATION *
                AVERAGE_WEIGHT;
              for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
                HISTOGRAM[LCV] := HISTOGRAM[LCV] * (1 - AVERAGE_WEIGHT) +
                  PRESENT_GENERATION.HISTOGRAM[LCV] *
                  AVERAGE_WEIGHT;
            end; (* with...do *)
          end; (* Compute Diversity Statistics *)
          with PRESENT_DIVERSITY do
            begin
              BEST := -maxlongint;
              WORST := maxlongint;
              MEAN := 0.0;
              for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
                with POPULATION(INDIVIDUAL_LCV) do
                  begin
                    MEAN := MEAN + DIVERSITY / MAX_NORMALIZED_DIVERSITY * 100;
                    MEAN := MEAN / POPULATION_SIZE;
                    SUM_OF_SQUARES := 0.0;
                    for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
                      SUM_OF_SQUARES := SUM_OF_SQUARES +
                        sqr(POPULATION(INDIVIDUAL_LCV).DIVERSITY /
                          MAX_NORMALIZED_DIVERSITY * 100 - MEAN);
                    VARIANCE := SUM_OF_SQUARES / POPULATION_SIZE;
                    STANDARD_DEVIATION := sqrt(VARIANCE);
                  end; (* with...do *)
                end; (* Compute long term statistics *)
                if GENERATION > DIVERSITY_MAX_AVERAGE_GENERATIONS
                  then
                    AVERAGE_WEIGHT := 1 / DIVERSITY_MAX_AVERAGE_GENERATIONS
                  else
                    AVERAGE_WEIGHT := 1 / GENERATION;
                with LONG_TERM_DIVERSITY do
                  begin
                    MEAN := MEAN * (1 - AVERAGE_WEIGHT) +
                      PRESENT_DIVERSITY.MEAN * AVERAGE_WEIGHT;
                    VARIANCE := VARIANCE * (1 - AVERAGE_WEIGHT) +
                      PRESENT_DIVERSITY.VARIANCE * AVERAGE_WEIGHT;
                    STANDARD_DEVIATION := STANDARD_DEVIATION * (1 - AVERAGE_WEIGHT) +
                      PRESENT_DIVERSITY.STANDARD_DEVIATION *
                      AVERAGE_WEIGHT;
                    for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do HISTOGRAM[LCV] := 0.0;
                    for INDIVIDUAL := 1 to POPULATION_SIZE do
                      begin
                        BUCKET := round(POPULATION(INDIVIDUAL).DIVERSITY /
                          MAX_NORMALIZED_DIVERSITY * 100);
                        if (BUCKET < LOWEST_BUCKET) or (HIGHEST_BUCKET < BUCKET) then
                          FITNESS_OUT_OF_RANGE := true;
                        HISTOGRAM[BUCKET] := HISTOGRAM[BUCKET] + 1;
                        end; (* for...to...do *)
                      SUM := 0.0;

```

```

for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
  SUM := SUM + HISTOGRAM(LCV);
for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
  HISTOGRAM(LCV) := HISTOGRAM(LCV) /
    (SUM / (HIGHEST_BUCKET - LOWEST_BUCKET + 1));
end; (* with...do *)

GENETIC_ALGORITHM_FITNESS_CONVERGED :=
  abs(LONG_TERM_STANDARD_DEVIATION) < CONVERGENCE_THRESHOLD;

GENETIC_ALGORITHM_DIVERSITY_CONVERGED :=
  abs(LONG_TERM_DIVERSITY_STANDARD_DEVIATION) < CONVERGENCE_THRESHOLD;
end; (* COMPUTE_POPULATION_STATISTICS *)

procedure EVALUATE_POPULATION;
(* This routine evaluates the fitness for each individual in the *)
(* population. *)
var
  (* Loop Control Variables. *)
  INDIVIDUAL, LCV : integer;

begin (* EVALUATE_POPULATION *)
  for INDIVIDUAL := 1 to POPULATION_SIZE do
    EVALUATE(INDIVIDUAL);
  COMPUTE_POPULATION_STATISTICS;
end; (* EVALUATE_POPULATION *)

procedure CROSSOVER_POPULATION;
(* Determines which survivors will be used to reproduce offspring to take *)
(* the place of all non-survivors. *)
var
  (* Each generation, the survivors must be counted so that selection *)
  (* for reproduction can be done using only survivors. *)
  NUMBER_OF_SURVIVORS : integer;

  (* Loop Control Variables. *)
  INDIVIDUAL_LCV, LCV : integer;
  GENE : RTOS_GENES;

  (* Indices of the parents that will be 'crossed' to produce the *)
  (* offspring. *)
  FIRST_PARENT, SECOND_PARENT : integer;

  (* For each parent the nth survivor is chosen, these variables keep *)
  (* track of the index to be chosen. *)
  FIRST_INDEX, SECOND_INDEX, COUNT : integer;

  (* Genes 0..CROSSOVER_POINT-1 are taken from the first parent and *)
  (* the rest are taken from the second parent. *)
  CROSSOVER_POINT : RTOS_GENES;

procedure NATURAL_SELECTION;
(* 'Culls' the population, retaining all individuals that are 'above *)
(* average' and also a percentage of the other individuals based on *)
(* PROBABILITY_OF_CROSSOVER. *)
var
  (* Loop Control Variable. *)
  INDIVIDUAL_LCV : integer;

  (* Total of the fitness of all individuals in the population. *)
  SUM : real;

begin (* NATURAL_SELECTION *)
  SUM := 0.0;
  for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
    SUM := SUM + POPULATION(INDIVIDUAL_LCV).FITNESS;

  for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
    with POPULATION(INDIVIDUAL_LCV) do
      RELATIVE_FITNESS := FITNESS / SUM;

  POPULATION[1].CUMULATIVE_FITNESS := POPULATION[1].RELATIVE_FITNESS;
  for INDIVIDUAL_LCV := 2 to POPULATION_SIZE do
    with POPULATION(INDIVIDUAL_LCV) do
      CUMULATIVE_FITNESS :=
        POPULATION(INDIVIDUAL_LCV-1).CUMULATIVE_FITNESS +
        RELATIVE_FITNESS;

  NUMBER_OF_SURVIVORS := 0;
  for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
    if (POPULATION(INDIVIDUAL_LCV).FITNESS > PRESENT_GENERATION.MEAN) xor
      ((random(maxint) / maxint) <= PROBABILITY_OF_CROSSOVER)
    then
      begin
        POPULATION(INDIVIDUAL_LCV).SURVIVOR := true;
        inc(NUMBER_OF_SURVIVORS);
      end (* if...then *)
    else
      POPULATION(INDIVIDUAL_LCV).SURVIVOR := false;
  end; (* NATURAL_SELECTION *)

begin (* CROSSOVER_POPULATION *)
  NATURAL_SELECTION;
  for INDIVIDUAL_LCV := 1 to POPULATION_SIZE do
    with POPULATION(INDIVIDUAL_LCV) do
      if not SURVIVOR then
        begin
          FIRST_INDEX := random(NUMBER_OF_SURVIVORS + 1);
          SECOND_INDEX := random(NUMBER_OF_SURVIVORS + 1);

          COUNT := 0;
          FIRST_PARENT := 1;
          SECOND_PARENT := 1;
          for LCV := 1 to POPULATION_SIZE do
            if POPULATION(LCV).SURVIVOR then
              begin
                inc(COUNT);
                if COUNT = FIRST_INDEX then FIRST_PARENT := LCV;
                if COUNT = SECOND_INDEX then SECOND_PARENT := LCV;
              end; (* if...then *)

          CROSSOVER_POINT := RTOS_GENES[random(integer(high(RTOS_GENES)))];
          if CROSSOVER_POINT <> low(RTOS_GENES) then
            for GENE := low(RTOS_GENES) to pred(CROSSOVER_POINT) do
              case GENE of
                TASKING_MODEL_GENE :
                  TASKING_MODEL := POPULATION(FIRST_PARENT).TASKING_MODEL;
                  TARGET_TIMESLICE_GENE :
                    TARGET_TIMESLICE :=
                      POPULATION(FIRST_PARENT).TARGET_TIMESLICE;
                    PRIORITY_INHERITANCE_GENE :
                      PRIORITY_INHERITANCE_ENABLED :=
                        POPULATION(FIRST_PARENT).PRIORITY_INHERITANCE_ENABLED;
                    PRIORITY_ALLOCATION_GENE :
                      PRIORITY_ALLOCATION :=
                        POPULATION(FIRST_PARENT).PRIORITY_ALLOCATION;
                    INITIAL_PRIORITY_ASSIGNMENT_GENE :
                      INITIAL_PRIORITY_ASSIGNMENT :=
                        POPULATION(FIRST_PARENT).INITIAL_PRIORITY_ASSIGNMENT;
                  end; (* case...of *)
                end (* if...then *)
              end; (* CROSSOVER_POPULATION *)

procedure MUTATE_POPULATION;
(* This routine performs the population wide mutation on individuals. *)
var
  (* Maximum number of mutations that will be performed for this *)
  (* generation (directly derived from probability of mutation). *)
  MAX_MUTATIONS : integer;

  (* The individual in the population being mutated. *)
  MUTANT_INDIVIDUAL : integer;

  (* The gene in the individual being mutated. *)
  MUTANT_GENE : RTOS_GENES;

  (* Loop Control variable. *)
  LCV : integer;

begin (* MUTATE_POPULATION *)
  MAX_MUTATIONS := round(POPULATION_SIZE * (integer(high(RTOS_GENES)) -
    integer(low(RTOS_GENES)) + 1) *
    PROBABILITY_OF_MUTATION);

  for LCV := 0 to random(MAX_MUTATIONS) do
    begin
      MUTANT_INDIVIDUAL := random(POPULATION_SIZE) + 1;
      MUTANT_GENE := RTOS_GENES[random(integer(high(RTOS_GENES)))];
      RANDOM_GENE(POPULATION(MUTANT_INDIVIDUAL), MUTANT_GENE);
    end; (* For...to...do *)
  end; (* MUTATE_POPULATION *)

begin (* PROCESS_GENETIC_ALGORITHM_INDIVIDUAL *)
  if INDIVIDUAL = POPULATION_SIZE then
    begin
      INDIVIDUAL := 0;
      if GENERATION = 0
      then
        INITIALIZE_POPULATION
      else
        begin
          CROSSOVER_POPULATION;
          MUTATE_POPULATION;
        end; (* if...then...else *)
      end; (* if...then *)

    inc(INDIVIDUAL);
    EVALUATE(INDIVIDUAL);

    PROCESS_GENETIC_ALGORITHM_INDIVIDUAL := INDIVIDUAL;
    if INDIVIDUAL = POPULATION_SIZE then
      begin
        inc(INDIVIDUAL);
        COMPUTE_POPULATION_STATISTICS;
      end; (* if...then *)
    end; (* PROCESS_GENETIC_ALGORITHM_INDIVIDUAL *)

procedure PROCESS_GENETIC_ALGORITHM_GENERATION;
(* This routine performs all processing on the population for an entire *)
(* generation (i.e., Selection/Crossover, Mutation and Evaluation). *)
var INDIVIDUAL_LCV : integer;

begin (* PROCESS_GENETIC_ALGORITHM_GENERATION *)
  INDIVIDUAL_LCV := 1;
  repeat
    PROCESS_GENETIC_ALGORITHM_INDIVIDUAL;
  until (INDIVIDUAL_LCV > POPULATION_SIZE);
end; (* PROCESS_GENETIC_ALGORITHM_GENERATION *)

procedure PRODUCE_GENETIC_ALGORITHM_REPORT;
(* This routine produces a report on the 'best' individuals found. *)
const
  (* General purpose special characters. *)
  CR = chr($0D); (* Carriage Return *)
  LF = chr($0A); (* Line Feed *)

var
  (* Loop Control Variables. *)
  LCV, COUNTER : integer;
  GENE : RTOS_GENES;

  (* Text file for report. *)
  REPORT : text;

  (* Temporary variable to hold the fitness of the individuals in the *)

```

## 9.4 MATH\_GA.PAS

```
(* report. *)
(*.....*)
THMP : real;
(*.....*)
(* Used to extract the genotype from the most popular diversity *)
(* value within the entire population. *)
(*.....*)
CHOICE : longint;

function REPORT_FILENAME : string;
(*.....*)
(* This routine determines the next available GA xx.RPT filename. *)
(*.....*)

const (*.....*)
(* Used to assign a unique filename to the report file. *)
(*.....*)
REPORT_NUMBER : integer = 0;

var (*.....*)
(* Result of I/O operation, used to find first available *)
(* filename. *)
(*.....*)
RESULT : integer;

(*.....*)
(* Filename of file to get report *)
(*.....*)
FILENAME : string(sizeof(namestr) + sizeof(dirstr) + 1);

(*.....*)
(* File used to find first available file. *)
(*.....*)
REPORT : file;

begin (* REPORT_FILENAME *)
  repeat
    FILENAME := 'GA ';
    FILENAME := FILENAME + chr(ord('0') + REPORT_NUMBER div 10) +
      chr(ord('0') + REPORT_NUMBER mod 10);
    FILENAME := FILENAME + '.RPT';
    inc(REPORT_NUMBER);
    assign(REPORT, FILENAME);
    (*SI*) reset(REPORT); (*SI*)
    RESULT := loresult;
    if RESULT = 0 then
      close(REPORT);
    until RESULT = 2;
    REPORT_FILENAME := FILENAME;
  end; (* REPORT_FILENAME *)

begin (* PRODUCE_GENETIC_ALGORITHM_REPORT *)
  assign(REPORT, REPORT_FILENAME);
  (*SI*) rewrite(REPORT); (*SI*)
  if FITNESS_OUT_OF_RANGE then
    writeln(REPORT, ' Fatal Error! Fitness out of range [0..100]!');
  if GENETIC_ALGORITHM_FITNESS_CONVERGED then
    writeln(REPORT, ' Fitness converged, ');
  if GENETIC_ALGORITHM_DIVERSITY_CONVERGED then
    writeln(REPORT, ' Diversity converged, ');
  writeln(REPORT, ' ', GENERATION, ' Generations');
  for COUNTER := 0 to (NUMBER_OF_BEST_TO_REPORT-1) do
    begin
      THMP := ALL_TIME_BEST[(LOWEST_IN_LIST_OF_BEST_INDIVIDUALS + COUNTER)
        mod
        NUMBER_OF_BEST_TO_REPORT].FITNESS;
      if THMP <> -maxlongint then
        begin
          writeln(REPORT, ' ');
          writeln(REPORT, ' ', ' ', ' ');
          writeln(REPORT, ' ');
          writeln(REPORT, ' Rank: ', NUMBER_OF_BEST_TO_REPORT - COUNTER);
          writeln(REPORT, ' Fitness: ', THMP : 3 : 3);
          WRITE TO INI FILE(REPORT, ALL_TIME_BEST[
            (LOWEST_IN_LIST_OF_BEST_INDIVIDUALS + COUNTER) mod
            NUMBER_OF_BEST_TO_REPORT]);
          writeln(REPORT, ' ');
        end; (* if...then *)
      end; (* for...to...do *)
    close(REPORT);
  end; (* PRODUCE_GENETIC_ALGORITHM_REPORT *)

begin (* GENETIC_ALGORITHM *)
  randomize;
  MAX_NORMALIZED_DIVERSITY := $003FFFF; (* 22 bits are used *)

  (* make sure application executable exists in this directory! *)
  assign(EXECUTABLE_FILE, PROGRAM_NAME + '.EXE');
  (*SI*) reset(EXECUTABLE_FILE); (*SI*)
  if loresult <> 0 then
    begin
      writeln;
      writeln('RTOS_GA: TRSKINS application does not exist, ( ',
        PROGRAM_NAME, '.EXE)');
      halt(1);
    end; (* if...then *)
  end. (* GENETIC_ALGORITHM *)

unit MATH_GA;
(*.....*)
(* Genetic Algorithm (GA) specification to determine the maximum value for *)
(* a multi-dimensional function. This specification is merely used to test *)
(* the basic operation of the GA. *)
(*.....*)

(*
  Compiler Options (Ver. 7.0)
  (*SA+ Word Alignment *)
  (*SE- Short Circuit Boolean Evaluation *)
  (*SD+ Debug Code Generation ON (Sort of) *)
  (*S+ Requires /V option to TPC to activate *)
  (*SL+ Local Debug symbols ON (Sort of) *)
  (*S+ Requires /V option to TPC to activate *)
  (*SP- Far calls only as needed *)
  (*SI- I/O Checking OFF *)
  (*SO+ Hardware 80x87 Used (if available) *)
  (*SO- Overlays NOT allowed *)
  (*SP- Standard 'string' parameters *)
  (*SQ- Overflow Checking OFF *)
  (*SR- Range Checking OFF *)
  (*SS- Stack Checking OFF *)
  (*SV- Var-string Checking OFF *)
  (*ST+ Force Typed 'S' references *)
  (*SX+ Enable Extended syntax *)
  *)

interface

const (*.....*)
(* The following defines the number of individuals in the population. *)
(*.....*)
POPULATION_SIZE = 175;

(*.....*)
(* Genetic operator probabilities; these define the speed that the GA *)
(* converges as well as the probability of false convergence. *)
(*.....*)
PROBABILITY_OF_CROSSOVER = 0.25;
PROBABILITY_OF_MUTATION = 0.03;
CONVERGENCE_THRESHOLD = 6;

(*.....*)
(* Long term statistic are computing using data from (up to) this *)
(* many previous generations. *)
(*.....*)
FITNESS_MAX_AVERAGE_GENERATIONS = 175;
DIVERSITY_MAX_AVERAGE_GENERATIONS = 100;

(*.....*)
(* Number of 'best' individuals to maintain for reporting. *)
(*.....*)
NUMBER_OF_BEST_TO_REPORT = 5;

(*.....*)
(* Defines the bounds of the histogram, values should be guaranteed *)
(* to be beyond the possible fitness values (this is verified and *)
(* reported any fitness is beyond this range). *)
(*.....*)
MINIMUM_FITNESS = 0.0; LOWEST_BUCKET = round(MINIMUM_FITNESS);
MAXIMUM_FITNESS = 100.0; HIGHEST_BUCKET = round(MAXIMUM_FITNESS);

const (*.....*)
(* Problem specific parameter. This is the number of dimensions of *)
(* the function to be searched. *)
(*.....*)
NUMBER_OF_VARIABLES = 2;

type (*.....*)
(* Problem specific parameter. This structure defines the 'genes' of *)
(* the individual that will make up the population. *)
(*.....*)
GENOTYPE = record
  (*.....*)
  (* The actual problem 'genes', for this problem they are all the *)
  (* same. In general this will not be the case. *)
  (*.....*)
  GENE : array [0..NUMBER_OF_VARIABLES-1] of real;
  (*.....*)
  (* The genes are allowed to range within these bounds so that each *)
  (* individual created is guaranteed to produce genes that are *)
  (* reasonable for the problem (i.e. all individuals are viable). *)
  (*.....*)
  UPPER_BOUND, LOWER_BOUND : array [0..NUMBER_OF_VARIABLES-1] of real;
  (*.....*)
  (* Genetic Algorithm bookkeeping variables used to determine which *)
  (* individuals are allowed to 'live' in the next generation. *)
  (*.....*)
  (* The individual's fitness evaluation parameters. *)
  (*.....*)
  FITNESS, RELATIVE_FITNESS, CUMULATIVE_FITNESS : real;
  (*.....*)
  (* The individual's diversity parameter (there is a 1:1 mapping *)
  (* of genotype values and diversity value). *)
  (*.....*)
  DIVERSITY : real;
  (*.....*)
  (* Flag to indicate whether the individual will 'live' in the *)
  (* next generation. *)
  (*.....*)
  SURVIVOR : boolean;
end; (* GENOTYPE *)

(*.....*)
(* This data structure is used to determine how well the GA is doing. *)
(* The only criteria used is convergence to a single solution (or to a *)
(* single 'fitness' value). The GA is assumed to be done when the *)
(* population has converged regardless of the actual fitness value *)
(* achieved. *)
(*.....*)
STATISTICS = record
  (*.....*)
  (* Basic values for the population... *)
  (*.....*)
  BEST, MEAN, WORST : real;
  (*.....*)
  (* Statistics based on the entire population... *)
  (*.....*)
  VARIANCE, STANDARD_DEVIATION : real;
  (*.....*)
  (* Normalized histogram buckets. *)
  (*.....*)
  HISTOGRAM : array [round(MINIMUM_FITNESS - 0.5) ..
    round(MAXIMUM_FITNESS + 0.5)] of real;
end; (* STATISTICS *)

const (*.....*)
(* This is the current generation that has completed processing. *)
(*.....*)
GENERATION : longint = 0;

```

```

[* This is the current individual that has complete processing. *]
[.....]
INDIVIDUAL : integer = POPULATION_SIZE;
[.....]
[* This is the maximum value for the genotype diversity, it is used *]
[* to scale the diversity display. *]
MAX_NORMALIZED_DIVERSITY : real = 0.0;
[.....]
[* Indicates the status of the GA, can be used to terminate the GA. *]
GENETIC_ALGORITHM_FITNESS_CONVERGED : boolean = false;
GENETIC_ALGORITHM_DIVERSITY_CONVERGED : boolean = false;

var
[.....]
[* Statistics are available for each generation that is processed, *]
[* the data is independent from generation to generation. *]
PRESENT_GENERATION : STATISTICS;
[.....]
[* Cumulative statistics are available for the most recent *]
[* generations (see FITNESS_MAX_AVERAGE_GENERATIONS). *]
LONG_TERM : STATISTICS;
[.....]
[* Diversity statistics for the present generation. *]
PRESENT_DIVERSITY : STATISTICS;
[.....]
[* Cumulative diversity statistics for the most recent generations *]
[* (see DIVERSITY_MAX_AVERAGE_GENERATIONS). *]
LONG_TERM_DIVERSITY : STATISTICS;
[.....]
[* These are all of the genes of the population. *]
POPULATION : array [0..POPULATION_SIZE] of GENOTYPE;

function PROCESS_GENETIC_ALGORITHM_INDIVIDUAL : integer;
[.....]
[* This routine performs all processing on the individual (i.e., Selection, *]
[* Crossover, Mutation and Evaluation). The individual that was processed *]
[* returned by this function. *]
[.....]

procedure PROCESS_GENETIC_ALGORITHM_GENERATION;
[.....]
[* This routine performs all processing on the population for an entire *]
[* generation (i.e., Selection, Crossover, Mutation and Evaluation). *]
[.....]

procedure PRODUCE_GENETIC_ALGORITHM_REPORT;
[.....]
[* This routine produces a report on the 'best' individuals found. *]
[.....]

implementation
uses dos;

var
[.....]
[* List of all time best individuals that will be reported. *]
ALL_TIME_BEST : array [0..(NUMBER_OF_BEST_TO_REPORT-1)] of GENOTYPE;

const
[.....]
[* The index of the lowest in the list is maintained so that the list *]
[* can be kept as a circular buffer very efficiently. *]
LOWEST_IN_LIST_OF_BEST_INDIVIDUALS : integer = 0;
[.....]
[* Flag that indicates that at least one fitness value was outside of *]
[* the allowable range [0.0 through 100.0], used in report. *]
FITNESS_OUT_OF_RANGE : boolean = false;

var
[.....]
[* INI file that contains the bounds for each variable. *]
INI_FILE : text;
[.....]
[* Loop control variables. *]
VARIABLE : integer;
[.....]
[* Bounds for each variable (from INI file). *]
LOWER, UPPER : real;

function RANDOM_VALUE [LOWER, UPPER : real] : real;
[.....]
[* Produces a random floating point number between LOWER and UPPER. *]
[.....]

begin [* RANDOM_VALUE *]
RANDOM_VALUE := [random]maxint / maxint * [UPPER - LOWER] + LOWER;
end; [* RANDOM_VALUE *]

function PROCESS_GENETIC_ALGORITHM_INDIVIDUAL : integer;
[.....]
[* This routine performs all processing on the individual (i.e., Selection, *]
[* Crossover, Mutation and Evaluation). The individual that was processed *]
[* returned by this function. *]
[.....]

procedure INITIALIZE_POPULATION;
[.....]
[* This routine initializes the GA by determining the bounds of the math *]
[* problem to be solved (via MATH_GA.INI). It also randomly initializes *]
[* the population and initializes the long term statistics. *]
[.....]

var
[.....]
[* Loop control variables. *]
INDIVIDUAL, LCV : integer;

begin [* INITIALIZE_POPULATION *]
[* Randomly initialize population *]
for INDIVIDUAL := 1 to POPULATION_SIZE do
for VARIABLE := 0 to NUMBER_OF_VARIABLES-1 do
with POPULATION[INDIVIDUAL] do
GENE[VARIABLE] := RANDOM_VALUE [LOWER_BOUND[VARIABLE],
UPPER_BOUND[VARIABLE]];

[* Initialize Statistics *]
with PRESENT_GENERATION do
begin
BEST := -maxlongint;
WORST := maxlongint;

MEAN := 0.0;
VARIANCE := 0.0;
STANDARD_DEVIATION := 0.0;
for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
HISTOGRAM[LCV] := 0.0;
end; [* with...do *]
LONG_TERM := PRESENT_GENERATION;
PRESENT_DIVERSITY := PRESENT_GENERATION;
for LCV := 0 to (NUMBER_OF_BEST_TO_REPORT-1) do
ALL_TIME_BEST[LCV].FITNESS := -maxlongint;
end; [* INITIALIZE_POPULATION *]

procedure EVALUATE [THIS_INDIVIDUAL : integer];
[.....]
[* This routine evaluates the fitness for a single individual of the *]
[* population. *]
[.....]

begin [* EVALUATE *]
with POPULATION [THIS_INDIVIDUAL] do
begin
FITNESS := 21.5 +
GENE[0] * sin[4 * pi * GENE[0]] +
GENE[1] * sin[10 * pi * GENE[1]];
end; [* with...do *]
end; [* EVALUATE *]

procedure COMPUTE_POPULATION_STATISTICS;
[.....]
[* This routine computes both the present generation and long term *]
[* statistics. *]
[.....]

var
[.....]
[* Loop Control Variables. *]
INDIVIDUAL, LCV : integer;
[.....]
[* For computing variance and correlation. *]
SUM_OF_SQUARES : real;
[.....]
[* For computing correlation. *]
SUM : real;
[.....]
[* Weight used to average for long term statistics. *]
AVERAGE_WEIGHT : real;
[.....]
[* For determining where to put the fitness value (in histogram) *]
BUCKET : integer;
[.....]
[* Loop Control Variable used to compute diversity value for each *]
[* genotype value (i.e., individual). *]
GENE_LCV : integer;

begin [* COMPUTE_POPULATION_STATISTICS *]
with PRESENT_GENERATION do
begin
BEST := -maxlongint;
WORST := maxlongint;
end; [* with...do *]

with PRESENT_GENERATION do
begin
MEAN := 0.0;
for INDIVIDUAL := 1 to POPULATION_SIZE do
with POPULATION[INDIVIDUAL] do
begin
MEAN := MEAN + FITNESS;
if FITNESS > PRESENT_GENERATION.BEST then
PRESENT_GENERATION.BEST := FITNESS;
if FITNESS > LONG_TERM.BEST then
begin
LONG_TERM.BEST := FITNESS;
ALL_TIME_BEST[LOWEST_IN_LIST_OF_BEST_INDIVIDUALS] :=
POPULATION[INDIVIDUAL];
LOWEST_IN_LIST_OF_BEST_INDIVIDUALS :=
(LOWEST_IN_LIST_OF_BEST_INDIVIDUALS + 1) mod
NUMBER_OF_BEST_TO_REPORT;
end; [* if...then *]
if FITNESS < PRESENT_GENERATION.WORST then
PRESENT_GENERATION.WORST := FITNESS;
if FITNESS < LONG_TERM.WORST then
LONG_TERM.WORST := FITNESS;
end;
MEAN := MEAN / POPULATION_SIZE;

SUM_OF_SQUARES := 0.0;
for INDIVIDUAL := 1 to POPULATION_SIZE do
SUM_OF_SQUARES := SUM_OF_SQUARES +
sq[POPULATION[INDIVIDUAL].FITNESS - MEAN];
VARIANCE := SUM_OF_SQUARES / POPULATION_SIZE;

STANDARD_DEVIATION := sqrt[VARIANCE];

for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do HISTOGRAM[LCV] := 0.0;
for INDIVIDUAL := 1 to POPULATION_SIZE do
begin
BUCKET := round[POPULATION[INDIVIDUAL].FITNESS];
if [BUCKET < LOWEST_BUCKET] or [HIGHEST_BUCKET < BUCKET] then
FITNESS_OUT_OF_RANGE := true;
HISTOGRAM[BUCKET] := HISTOGRAM[BUCKET] + 1;
end; [* for...to...do *]
SUM := 0.0;
for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
SUM := SUM + HISTOGRAM[LCV];
for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
HISTOGRAM[LCV] := HISTOGRAM[LCV] /
(SUM / (HIGHEST_BUCKET - LOWEST_BUCKET + 1));
end; [* with...do *]

[* Compute long term statistics *]
if GENERATION > FITNESS_MAX_AVERAGE_GENERATIONS
then
AVERAGE_WEIGHT := 1 / FITNESS_MAX_AVERAGE_GENERATIONS
else
AVERAGE_WEIGHT := 1 / GENERATION;

with LONG_TERM do
begin
MEAN := MEAN * [1 - AVERAGE_WEIGHT] +
PRESENT_GENERATION.MEAN * AVERAGE_WEIGHT;
VARIANCE := VARIANCE * [1 - AVERAGE_WEIGHT] +
PRESENT_GENERATION.VARIANCE * AVERAGE_WEIGHT;
STANDARD_DEVIATION := STANDARD_DEVIATION * [1 - AVERAGE_WEIGHT] +
PRESENT_GENERATION.STANDARD_DEVIATION *
AVERAGE_WEIGHT;
for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
HISTOGRAM[LCV] := HISTOGRAM[LCV] * [1 - AVERAGE_WEIGHT] +

```



```

                PRSSENT_GENERATION.HISTOGRAM(LCV) *
                AVERAGE_WEIGHT;
end; (* with...do *)

(* Compute Diversity statistics *)
with PRSSENT_DIVERSITY do
begin
    BEST := -maklongint;
    WORST := maklongint;
    MEAN := 0.0;
    for INDIVIDUAL := 1 to POPULATION_SIZE do
        with POPULATION(INDIVIDUAL) do
            begin
                DIVERSITY := 0.0;
                for GENE_LCV := 0 to NUMBER_OF_VARIABLES-1 do
                    DIVERSITY := DIVERSITY + GENE(GENE_LCV);
                DIVERSITY := 100 * DIVERSITY / MAX_NORMALIZED_DIVERSITY;
                MEAN := MEAN + DIVERSITY;
            end;
            MEAN := MEAN / POPULATION_SIZE;

            SUM_OF_SQUARES := 0.0;
            for INDIVIDUAL := 1 to POPULATION_SIZE do
                SUM_OF_SQUARES := SUM_OF_SQUARES +
                    sq(POPULATION(INDIVIDUAL).DIVERSITY - MEAN);
            VARIANCE := SUM_OF_SQUARES / POPULATION_SIZE;

            STANDARD_DEVIATION := sqrt(VARIANCE);
        end; (* with...do *)
    end;

    (* Compute long term statistics *)
    if GENERATION > DIVERSITY_MAX_AVERAGE_GENERATIONS
    then
        AVERAGE_WEIGHT := 1 / DIVERSITY_MAX_AVERAGE_GENERATIONS
    else
        AVERAGE_WEIGHT := 1 / GENERATION;
    end;

with LONG_TERM_DIVERSITY do
begin
    MEAN := MEAN * (1 - AVERAGE_WEIGHT) +
        PRSSENT_DIVERSITY.MEAN * AVERAGE_WEIGHT;
    VARIANCE := VARIANCE * (1 - AVERAGE_WEIGHT) +
        PRSSENT_DIVERSITY.VARIANCE * AVERAGE_WEIGHT;
    STANDARD_DEVIATION := STANDARD_DEVIATION * (1 - AVERAGE_WEIGHT) +
        PRSSENT_DIVERSITY.STANDARD_DEVIATION *
        AVERAGE_WEIGHT;
    for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do HISTOGRAM(LCV) := 0.0;
    for INDIVIDUAL := 1 to POPULATION_SIZE do
        begin
            BUCKET := round(POPULATION(INDIVIDUAL).DIVERSITY);
            if (BUCKET < LOWEST_BUCKET) or (HIGHEST_BUCKET < BUCKET) then
                FITNESS_OUT_OF_RANGE := true;
                HISTOGRAM(BUCKET) := HISTOGRAM(BUCKET) + 1;
            end; (* for...to...do *)
        end;
    SUM := 0.0;
    for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
        SUM := SUM + HISTOGRAM(LCV);
    end;
    for LCV := LOWEST_BUCKET to HIGHEST_BUCKET do
        HISTOGRAM(LCV) := HISTOGRAM(LCV) /
            (SUM / (HIGHEST_BUCKET - LOWEST_BUCKET + 1));
    end; (* with...do *)

GENETIC_ALGORITHM_FITNESS_CONVERGED :=
    abs(LONG_TERM.STANDARD_DEVIATION) < CONVERGENCE_THRESHOLD;

GENETIC_ALGORITHM_DIVERSITY_CONVERGED :=
    abs(LONG_TERM.DIVERSITY.STANDARD_DEVIATION) < CONVERGENCE_THRSHOLD;
end; (* COMPUTE_POPULATION_STATISTICS *)

procedure EVALUATE_POPULATION;
(* This routine evaluates the fitness for each individual in the
   * population.
   *)
var
    (* Loop Control variables.
    *)
    INDIVIDUAL, LCV : integer;
begin (* EVALUATE_POPULATION *)
    for INDIVIDUAL := 1 to POPULATION_SIZE do
        EVALUATE(INDIVIDUAL);
        COMPUTE_POPULATION_STATISTICS;
    end; (* EVALUATE_POPULATION *)
end;

procedure CROSSOVER_POPULATION;
(* Determines which survivors will be used to reproduce offspring to take
   * the place of all non-survivors.
   *)
var
    (* Each generation, the survivors must be counted so that selection
    * for reproduction can be done using only survivors.
    *)
    NUMBER_OF_SURVIVORS : integer;
    (* Loop Control Variables.
    *)
    INDIVIDUAL, LCV : integer;
    (* Indices of the parents that will be 'crossed' to produce the
    * offspring.
    *)
    FIRST_PARENT, SECOND_PARENT : integer;
    (* For each parent the Nth survivor is chosen, these variables keep
    * track of the index to be chosen.
    *)
    FIRST_INDEX, SECOND_INDEX, COUNT : integer;
    (* Genes 0..CROSSOVER_POINT-1 are taken from the first parent and
    * the rest are taken from the second parent.
    *)
    CROSSOVER_POINT : integer;
begin
    inc(INDIVIDUAL);
    EVALUATE(INDIVIDUAL);
    PROCESS_GENETIC_ALGORITHM_INDIVIDUAL := INDIVIDUAL;
    if INDIVIDUAL = POPULATION_SIZE then
        begin
            inc(GENERATION);
            COMPUTE_POPULATION_STATISTICS;
            end; (* if...then *)
    end; (* PROCESS_GENETIC_ALGORITHM_INDIVIDUAL *)
end;

procedure PROCESS_GENETIC_ALGORITHM_GENERATION;
(* This routine performs all processing on the population for an entire
   * generation [i.e., Selection/Crossover, Mutation and Evaluation].
   *)
var INDIVIDUAL_LCV : integer;
begin (* PROCESS_GENETIC_ALGORITHM_GENERATION *)
    INDIVIDUAL_LCV := 1;
    repeat
        PROCESS_GENETIC_ALGORITHM_INDIVIDUAL;
        inc(INDIVIDUAL_LCV);
    until (INDIVIDUAL_LCV > POPULATION_SIZE);
end;

SUM := 0.0;
for INDIVIDUAL := 1 to POPULATION_SIZE do
    SUM := SUM + POPULATION(INDIVIDUAL).FITNESS;
end;

for INDIVIDUAL := 1 to POPULATION_SIZE do
    with POPULATION(INDIVIDUAL) do
        RSLATIVE_FITNESS := FITNESS / SUM;
        POPULATION(1).CUMULATIVE_FITNESS := POPULATION(1).RSLATIVE_FITNESS;
        for INDIVIDUAL := 2 to POPULATION_SIZE do
            with POPULATION(INDIVIDUAL) do
                CUMULATIVE_FITNESS := POPULATION(INDIVIDUAL-1).CUMULATIVE_FITNESS +
                    RSLATIVE_FITNESS;
            end;
        end;
    end;

NUMBER_OF_SURVIVORS := 0;
for INDIVIDUAL := 1 to POPULATION_SIZE do
    if (POPULATION(INDIVIDUAL).FITNESS > PRSSENT_GENERATION.MEAN) or
        ((random(maxint) / maxint) <= PROBABILITY_OF_CROSSOVER)
    then
        begin
            POPULATION(INDIVIDUAL).SURVIVOR := true;
            inc(NUMBER_OF_SURVIVORS);
            end; (* if...then *)
        else
            POPULATION(INDIVIDUAL).SURVIVOR := false;
        end; (* NATURAL_SELECTION *)
    end;

begin (* CROSSOVER_POPULATION *)
    NATURAL_SELECTION;
    for INDIVIDUAL := 1 to POPULATION_SIZE do
        with POPULATION(INDIVIDUAL) do
            if not SURVIVOR then
                begin
                    FIRST_INDEX := random(NUMBER_OF_SURVIVORS) + 1;
                    SECOND_INDEX := random(NUMBER_OF_SURVIVORS) + 1;

                    COUNT := 0;
                    FIRST_PARENT := 1;
                    SECOND_PARENT := 1;
                    for LCV := 1 to POPULATION_SIZE do
                        if POPULATION(LCV).SURVIVOR then
                            begin
                                inc(COUNT);
                                if COUNT = FIRST_INDEX then FIRST_PARENT := LCV;
                                if COUNT = SECOND_INDEX then SECOND_PARENT := LCV;
                                end; (* if...then *)
                            end;
                    end;
                    CROSSOVER_POINT := random(NUMBER_OF_VARIABLES);
                    for LCV := 0 to CROSSOVER_POINT-1 do
                        GENE(LCV) := POPULATION(FIRST_PARENT).GENE(LCV);
                    for LCV := CROSSOVER_POINT to NUMBER_OF_VARIABLES-1 do
                        GENE(LCV) := POPULATION(SECOND_PARENT).GENE(LCV);
                    end; (* if...then *)
                end; (* CROSSOVER_POPULATION *)
            end;
        end;
    end;

procedure MUTATE_POPULATION;
(* This routine performs the population wide mutation on individuals.
   *)
var
    (* Maximum number of mutations that will be performed for this
    * generation (directly derived from probability of mutation).
    *)
    MAX_MUTATIONS : integer;
    (* The individual in the population being mutated.
    *)
    MUTANT_INDIVIDUAL : integer;
    (* The gene in the individual being mutated.
    *)
    MUTANT_GENE : integer;
    (* Loop Control variable.
    *)
    LCV : integer;
begin (* MUTATE_POPULATION *)
    MAX_MUTATIONS := round(POPULATION_SIZE * NUMBER_OF_VARIABLES *
        PROBABILITY_OF_MUTATION);
    for LCV := 0 to random(MAX_MUTATIONS) do
        begin
            MUTANT_INDIVIDUAL := random(POPULATION_SIZE) + 1;
            MUTANT_GENE := random(NUMBER_OF_VARIABLES);
            with POPULATION(MUTANT_INDIVIDUAL) do
                GENE(MUTANT_GENE) := RANDOM_VALUE
                    (LOWER_BOUND(MUTANT_GENE),
                     UPPER_BOUND(MUTANT_GENE));
            end; (* for...to...do *)
        end; (* MUTATE_POPULATION *)
    end;
begin (* PROCESS_GENETIC_ALGORITHM_INDIVIDUAL *)
    if INDIVIDUAL = POPULATION_SIZE then
        begin
            INDIVIDUAL := 0;
            if GENERATION = 0
            then
                INITIALIZE_POPULATION
            else
                CROSSOVER_POPULATION;
                MUTATE_POPULATION;
                end; (* if...then...else *)
            end; (* if...then *)
        inc(INDIVIDUAL);
        EVALUATE(INDIVIDUAL);
        PROCESS_GENETIC_ALGORITHM_INDIVIDUAL := INDIVIDUAL;
        if INDIVIDUAL = POPULATION_SIZE then
            begin
                inc(GENERATION);
                COMPUTE_POPULATION_STATISTICS;
                end; (* if...then *)
            end; (* PROCESS_GENETIC_ALGORITHM_INDIVIDUAL *)
        end;
    end;
end;

procedure NATURAL_SELECTION;
(* 'Culls' the population, retaining all individuals that are 'above
   * average' and also a percentage of the other individuals based on the
   * PROBABILITY_OF_CROSSOVER.
   *)
var
    (* Loop Control Variable.
    *)
    INDIVIDUAL : integer;
    (* Total of the fitness of all individuals in the population.
    *)
    SUM : real;
begin (* NATURAL_SELECTION *)

```

```

end: (* PROCESS_GENETIC_ALGORITHM_GENERATION *)

procedure PRODUCE_GENETIC_ALGORITHM_REPORT;
(* This routine produces a report on the 'best' individuals found. *)
const
  (* General purpose special characters. *)
  CR = chr(SOD); (* Carriage Return *)
  LF = chr(SOA); (* Line Feed *)
var
  (* Loop Control Variables. *)
  LCV, COUNTER : integer;
  (* Text file for report. *)
  REPORT : text;
function REPORT_FILENAME : string;
(* This routine determines the next available GA_xx.RPT filename. *)
const
  (* Used to assign a unique filename to the report file. *)
  REPORT_NUMBER : integer = 0;
var
  (* Result of I/O operation, used to find first available *)
  (* filename. *)
  RESULT : integer;
  (* Filename of file to get report. *)
  FILENAME : string[ sizeof(namestr) + sizeof(dirstr) + 1];
  (* File used to find first available file. *)
  REPORT : file;
begin (* REPORT_FILENAME *)
  repeat
    FILENAME := 'GA ';
    FILENAME = FILENAME + chr(ord('0') + REPORT_NUMBER div 10) +
      chr(ord('0') + REPORT_NUMBER mod 10);
    FILENAME := FILENAME + '.RPT';
    assign(REPORT, FILENAME);
    (*SI*) reset(REPORT); (*SI*)
    RESULT := ioreult;
    if RESULT = 0 then
      close(REPORT);
    until RESULT = 2;
    REPORT_FILENAME := FILENAME;
  end; (* REPORT_FILENAME *)
begin (* PRODUCE_GENETIC_ALGORITHM_REPORT *)
  assign(REPORT, REPORT_FILENAME);
  (*SI*) rewrite(REPORT); (*SI*)
  if FITNESS_OUT_OF_RANGE then
    writeln(REPORT, 'Fatal Error! Fitness out of range [0..100]!');
  writeln;
  if GENETIC_ALGORITHM_FITNESS_CONVERGED then
    write(REPORT, 'Algorithm converged. ');
    writeln(REPORT, GENERATION, ' Generations');
    write(REPORT, ' Rank Fitness');
    for LCV := 0 to NUMBER_OF_VARIABLES-1 do
      write(REPORT, ' Var[', LCV, ']: ');
    for COUNTER := (NUMBER_OF_BEST_TO_REPORT-1) downto 0 do
      with ALL_TIME_BEST[(LOWEST_IN_LIST_OF_BEST_INDIVIDUALS + COUNTER) mod
        NUMBER_OF_BEST_TO_REPORT] do
        if FITNESS <> -maxlongint then
          begin
            writeln(REPORT);
            write(REPORT, NUMBER_OF_BEST_TO_REPORT - COUNTER : 6,
              FITNESS : 10 : 3);
            for LCV := 0 to NUMBER_OF_VARIABLES-1 do
              write(REPORT, GEN[LCV] : 10 : 3);
            end; (* with...do *)
          end;
        close(REPORT);
      end; (* PRODUCE_GENETIC_ALGORITHM_REPORT *)
end; (* GENETIC_ALGORITHM *)
(* Process INI file *)
assign(INI_FILE, 'MATH_GA.INI');
(*SI*) reset(INI_FILE); (*SI*)
if ioreult <> 0 then
  begin
    writeln;
    writeln('MATH_GA: Unable to open INI file [MATH_GA.INI]');
    halt[1];
  end; (* if...then *)
readln(INI_FILE, randseed);
for VARIABLE := 0 to NUMBER_OF_VARIABLES-1 do
  begin
    readln(INI_FILE, LOWER, UPPER);
    for INDIVIDUAL := 1 to POPULATION_SIZE do
      with POPULATION[INDIVIDUAL] do
        begin
          LOWER_BOUND[VARIABLE] := LOWER;
          UPPER_BOUND[VARIABLE] := UPPER;
        end; (* For...to...do *)
      MAX_NORMALIZED_DIVERSITY := MAX_NORMALIZED_DIVERSITY + UPPER;
    end; (* for...to...do *)
  end;
close(INI_FILE);
end; (* GENETIC_ALGORITHM *)

```

# 10. Appendix C Support Software Source Code

## 10.1 DINING.PAS

```
.....
LCV : integer;

procedure UPDATE_DISPLAY(NUM_THINKING, NUM_HUNGRY, NUM_EATING : integer);
.....
(* This routine updates all of the bar graphs and averages, it is called *)
(* whenever the philosopher has changed activities. *)
.....
const
.....
(* Maintains the averages of philosopher states over all updates. *)
AVG_THINKING : real = 0;
AVG_HUNGRY : real = 0;
AVG_EATING : real = 0;

procedure BAR_GRAPH(NUM : integer; var AVERAGE : real);
.....
(* This routine draws a bar graph which is proportional to the NUM *)
(* parameter passed in. In addition the average bar position is updated *)
(* to reflect the new bar graph. *)
.....
const
.....
(* Used to reduce 'jitter' on averages. *)
WEIGHT = 15;

var
.....
(* Temporary variable to hold the percent graphed and displayed. *)
PERCENT : real;

.....
(* Loop Control Variable to draw bar graph. *)
LCV : integer;

begin (* BAR_GRAPH *)
AVERAGE := (AVERAGE * WEIGHT + NUM) / (WEIGHT + 1);
PERCENT := (AVERAGE / NUMBER_OF_PHILOSOPHERS) * 100;
write(round(PERCENT) : 4, '% ');
PERCENT := (NUM / NUMBER_OF_PHILOSOPHERS) * BAR_SIZE;
for LCV := 1 to round(PERCENT) do write('#');
write(' | : BAR_SIZE + 1 - round(PERCENT));
PERCENT := (AVERAGE / NUMBER_OF_PHILOSOPHERS) * BAR_SIZE;
gotoxy(16 + round(PERCENT), Wherey); write('█');
end; (* BAR_GRAPH *)

begin (* UPDATE_DISPLAY *)
gotoxy(1, 10); write('Thinking '); BAR_GRAPH(NUM_THINKING, AVG_THINKING);
gotoxy(1, 12); write('Hungry '); BAR_GRAPH(NUM_HUNGRY, AVG_HUNGRY);
gotoxy(1, 14); write('Eating '); BAR_GRAPH(NUM_EATING, AVG_EATING);
end; (* UPDATE_DISPLAY *)

procedure PHILOSOPHER_TASK(TASK_ID : TASK_IDS; PRIORITY : USER_PRIORITIES; far:
.....
(* This task represents each philosopher. Each one will act completely *)
(* independently, i.e., they will not cooperate. Of course 'no man (or *)
(* philosopher) is an island', all philosophers must interact with each *)
(* other but not cooperatively. *)
.....
var
.....
(* In order to perform the simulation each philosopher must know *)
(* his/her seat at the table. This variable does that. *)
PHILOSOPHER_NUMBER : integer;

procedure THINK(PHILOSOPHER_NUMBER : integer);
.....
(* This routine models the thinking activity of a philosopher, all that *)
(* it really does is give the philosopher a place to be when it is in the *)
(* THINKING state. *)
.....
var
.....
(* Loop Control Variable and dummy to 'do work'. *)
THINK_COUNT, DUMMY : longint;

begin (* THINK *)
(* Do some non-blocking work *)
for THINK_COUNT := 1 to DELAY_COUNT do
(* Nothing really *) DUMMY := THINK_COUNT;

(* Do some blocking work *)
WAIT FOR DELAY(DELAY);
end; (* THINK *)

procedure EAT(PHILOSOPHER_NUMBER : integer);
.....
(* This routine models the eating activity of a philosopher, all that it *)
(* really does is give the philosopher a place to be when it is in the *)
(* EATING state. *)
.....
var
.....
(* Loop Control Variable and dummy to 'do work'. *)
EAT_COUNT, DUMMY : longint;

begin (* EAT *)
(* Do some non-blocking work *)
for EAT_COUNT := 1 to DELAY_COUNT do
(* Nothing really *) DUMMY := EAT_COUNT;

(* Do some blocking work *)
WAIT FOR DELAY(DELAY);
end; (* EAT *)

function LEFT_NEIGHBOR(PHILOSOPHER_NUMBER : integer) : integer;
.....
(* This function calculates the place setting (i.e., philosopher number) *)
(* of the philosopher to the left. *)
.....
begin (* LEFT_NEIGHBOR *)
LEFT_NEIGHBOR := (PHILOSOPHER_NUMBER + (NUMBER_OF_PHILOSOPHERS - 1)) mod
NUMBER_OF_PHILOSOPHERS;
end; (* LEFT_NEIGHBOR *)

function RIGHT_NEIGHBOR(PHILOSOPHER_NUMBER : integer) : integer;
.....
(* This function calculates the place setting (i.e., philosopher number) *)
(* of the philosopher to the right. *)
.....
begin (* RIGHT_NEIGHBOR *)
RIGHT_NEIGHBOR := (PHILOSOPHER_NUMBER + 1) mod NUMBER_OF_PHILOSOPHERS;
.....
.....
program DINING_PHILOSOPHERS(input, output);
.....
(* This is a solution to the 'classic' multi-tasking problem called *)
(* 'Dining Philosophers'. This solution was taken from 'Modern Operating *)
(* Systems' by Andrew Tanenbaum. *)
.....
(* Compiler Options (Ver. 7.0) *)
(*$A+ Word Alignment *)
(*$B+ Short Circuit Boolean Evaluation *)
(*$D+ Debug Code Generation ON (Sort of) *)
(* Requires /V option to TPC to activate *)
(*$L+ Local Debug symbols ON (Sort of) *)
(* Requires /V option to TPC to activate *)
(*$F+ Far calls only as needed *)
(*$G+ Generic 80x86 code only *)
(*$I+ I/O Checking OFF *)
(*$M 1024,150000,655360 Memory (Stack, Minheap, Maxheap) *)
(*$N+ Software Emulation of 80x87 *)
(*$E+ No 80x87 run-time emulation *)
(*$O+ Overlays NOT allowed *)
(*$P+ Standard 'string' parameters *)
(*$Q+ Overflow Checking OFF *)
(*$R+ Range Checking OFF *)
(*$T+ Stack Checking OFF *)
(*$S+ Force Typed 'Q' references *)
(*$V+ Var-string Checking ON *)
(*$X+ Disable Extended syntax *)
.....
uses crt, TASKING;

const
.....
(* In order to have the program run consistently across computers *)
(* widely varying processing power, the philosophers must include time *)
(* to 'delay' when eating and thinking. *)
.....
DELAY : TIME =
(
DAYS : 0;
HOURS : 0;
MINUTES : 0;
SECONDS : 1;
MILLISECONDS : 0;
);

.....
(* General purpose special characters. *)
CR = chr(13); (* Carriage Return *)
LF = chr(10); (* Line Feed *)

.....
(* This is used to allow the philosophers to all eat and think for *)
(* the same amount of time. *)
DELAY_COUNT = 100000;

.....
(* These are used to maintain/limit the number of philosophers that *)
(* the user has indicated are present. *)
NUMBER_OF_PHILOSOPHERS = integer * 20; (* Default value *)
MAX_NUMBER_OF_PHILOSOPHERS = MAXIMUM_NUMBER_OF_TASKS - 1;

.....
(* The size of the activity percent bar graph. *)
BAR_SIZE = 60;

.....
(* This software event is broadcasted to all philosophers so that all *)
(* are ready to eat at the same instant. *)
DINNER_IS_SERVED = EVENT * UNSIGNALLED;

.....
(* Controls mutually exclusive access to the 'forks' on the table. *)
MUTEX : BINARY_SEMAPHORE = 1;

type
.....
(* Philosophers only want to think and eat, being hungry is necessary *)
(* before eating is possible. *)
ACTIVITIES = (THINKING, HUNGRY, EATING);

var
.....
(* Table of philosopher activities (one for each guest). *)
ACTIVITY array[0..(MAX_NUMBER_OF_PHILOSOPHERS-1)] of ACTIVITIES;

.....
(* Table of semaphores for philosophers to wait on if they need but *)
(* cannot get forks. *)
NEED_FORKS : array[0..MAX_NUMBER_OF_PHILOSOPHERS-1] of BINARY_SEMAPHORE;

.....
(* Statistics for bar graph display of philosopher activities. *)
NUM_THINKING, NUM_EATING : integer;

.....
(* These are the task identifiers used in the application. *)
CONTROL, GUEST : TASK_IDS;

.....
(* Generic task attribute variable used to create all tasks. *)
TASK_ATTR : TASK_ATTRIBUTES;

.....
(* For command line timeslice string conversion. *)
CODE, TEMP : integer; VALUE : longint;

.....
(* For user guest number prompt. *)
CH : char;

.....
(* Loop Control Variable for philosopher table initialization. *)
```

```

end; (* RIGHT_NEIGHBOR *)

procedure CHECK_FORKS(PHILOSOPHER_NUMBER : integer);
(* This procedure models the action of the philosopher looking at the
 * table to see if there is a fork on each side of his plate (of course *)
 * This is prefaced with the fact that the philosopher must be HUNGRY. *)
begin (* CHECK_FORKS *)
  if (ACTIVITY[PHILOSOPHER_NUMBER] = NUMGRY) and
    (ACTIVITY[LEFT_NEIGHBOR(PHILOSOPHER_NUMBER)] <> EATING) and
    (ACTIVITY[RIGHT_NEIGHBOR(PHILOSOPHER_NUMBER)] <> EATING) then
    begin
      ACTIVITY[PHILOSOPHER_NUMBER] := EATING;
      inc(NUM_EATING);
      UPDATE_DISPLAY(NUM_THINKING, NUMBER_OF_PHILOSOPHERS - NUM_THINKING
                    - NUM_EATING, NUM_EATING);
      SIGNAL_BINARY_SEMAPHORS(NEED_FORKS[PHILOSOPHER_NUMBER]);
    end; (* If...then *)
  end; (* CHECK_FORKS *)
end;

procedure GET_FORKS(PHILOSOPHER_NUMBER : integer);
(* This routine models the act of a philosopher (who is hungry) trying to
 * eat. Before he can eat he must get two forks. *)
begin (* GET_FORKS *)
  WAIT_ON_BINARY_SEMAPHORS(MUTEX);

  ACTIVITY[PHILOSOPHER_NUMBER] := HUNGRY;
  dec(NUM_THINKING);
  UPDATE_DISPLAY(NUM_THINKING, NUMBER_OF_PHILOSOPHERS - NUM_THINKING
                - NUM_EATING, NUM_EATING);
  CHECK_FORKS(PHILOSOPHER_NUMBER);

  SIGNAL_BINARY_SEMAPHORE(MUTEX);
  WAIT_ON_BINARY_SEMAPHORE(NEED_FORKS[PHILOSOPHER_NUMBER]);
end; (* GET_FORKS *)

procedure PUT_FORKS(PHILOSOPHER_NUMBER : integer);
(* This routine models the act of a philosopher (who has finished eating) *)
(* putting his forks back on the table. *)
begin (* PUT_FORKS *)
  WAIT_ON_BINARY_SEMAPHORS(MUTEX);

  dec(NUM_EATING);
  ACTIVITY[PHILOSOPHER_NUMBER] := THINKING;
  inc(NUM_THINKING);
  UPDATE_DISPLAY(NUM_THINKING, NUMBER_OF_PHILOSOPHERS - NUM_THINKING
                - NUM_EATING, NUM_EATING);
  CHECK_FORKS(LEFT_NEIGHBOR(PHILOSOPHER_NUMBER));
  CHECK_FORKS(RIGHT_NEIGHBOR(PHILOSOPHER_NUMBER));

  SIGNAL_BINARY_SEMAPHORS(MUTEX);
end; (* PUT_FORKS *)

begin (* PHILOSOPHER TASK *)
  (* Determine philosopher number (must be 0 to N-1) *)
  PHILOSOPHER_NUMBER := TASK_ID - 2;

  (* Wait so that all philosophers start at the same instant *)
  WAIT_ON_EVENT(DINNER_IS_SERVED);

  (* Do philosopher stuff... (forever) *)
  repeat
    THINK(PHILOSOPHER_NUMBER);
    GET_FORKS(PHILOSOPHER_NUMBER);
    EAT(PHILOSOPHER_NUMBER);
    PUT_FORKS(PHILOSOPHER_NUMBER);
  until false;
end; (* PHILOSOPHER_TASK *)

procedure CONTROL_TASK(TASK_ID : TASK_IDS; PRIORITY : USER_PRIORITIES);
(* This is just a basic bookkeeping task that sets up the dinner, announces *)
(* that dinner is served and then terminates the simulation upon the user *)
(* request. *)
var (* Character to hold command from user (i.e., Start or Quit). *)
  COMMAND : char;

begin (* CONTROL_TASK *)
  (* Display Start-up Message *)
  writeln(CR, LF, 'S' = Start' : 29, 'Q' = Quit' : 30);
  repeat COMMAND := WAIT_ON_READKEY until COMMAND in ['s', 'S', 'q', 'Q'];

  if COMMAND in ['s', 'S'] then
    begin
      gotoxy(1, 6); writeln('Guests = ', NUMBER_OF_PHILOSOPHERS);
      gotoxy(1, 8); writeln('Activity Avg 1');
      gotoxy(16, 9);
      write('|'); for LCV := 1 to BAR_SIZE do write('-'); write('|');
      gotoxy(16+BAR_SIZE div 2, 9); write('|');
      gotoxy(16, 8); write('0');
      gotoxy(16+BAR_SIZE div 2, 8); write('50');
      gotoxy(16+BAR_SIZE, 8); write('100');
      gotoxy(16, 11); write('|', '|', '| : BAR_SIZE + 1);
      gotoxy(16, 13); write('|', '|', '| : BAR_SIZE + 1);
      gotoxy(16, 15);
      write('|'); for LCV := 1 to BAR_SIZE do write('-'); write('|');

      UPDATE_DISPLAY(NUM_THINKING, NUMBER_OF_PHILOSOPHERS - NUM_THINKING
                    - NUM_EATING, NUM_EATING);
      BROADCAST_EVENT(DINNER_IS_SERVED);
      repeat until WAIT_ON_READKEY in ['q', 'Q'];
    end; (* if...then *)
    writeln(CR, LF);
    halt(0);
  end; (* CONTROL_TASK *)

function UP_STRING(S : string) : string;
(* This routine returns a string corresponding to the upper case value of *)
(* the string passed to it. *)
var LCV : integer;

begin (* UP_STRING *)
  UP_STRING[0] := S[0];
  for LCV := 1 to length(S) do
    UP_STRING[LCV] := upcase(S[LCV]);
  end; (* UP_STRING *)

begin (* DINING_PHILOSOPHERS *)
  (* Parse parameters to setup TASKING environment *)
  if paramcount > 0 then
    if UP_STRING(paramstr(1)) = '-PSEMPT'
    then
      TASKING_CONFIGURATION.TASKING_MODEL := PSEMPTIVE
    else
      begin
        writeln('TASKS: Invalid parameter: [' , paramstr(1), ']');
        halt(1);
      end; (* if...then...else *)
    if paramcount > 2 then
      if UP_STRING(paramstr(2)) = '-TIMESLICE'
      then
        begin
          val(paramstr(3), VALUE, CODE);
          if CODE <> 0 then
            begin
              writeln('TASKS: Invalid number: [' , paramstr(3), ']');
              halt(1);
            end; (* if...then *)
            TASKING_CONFIGURATION.TARGET_TIMESLICE := VALUE;
          end (* if...then *)
        else
          begin
            writeln('TASKS: Invalid parameter: [' , paramstr(2), ']');
            halt(1);
          end; (* if...then...else *)
        end;

      (* A good idea during application development, bad idea after that *)
      TASKING_CONFIGURATION.STATISTICS := ALL_STATISTICS;

      (* Create TASKING startup attributes for all tasks *)
      with TASK_ATTR do
        begin
          PRIORITY := 20;
          STACK_WORDS_NEEDED := 500;
          (* Use default error handlers, carried through all task creations *)
          ERROR_HANDLERS(TASK_ALREADY_ACTIVE) := nil;
          ERROR_HANDLERS(INSUFFICIENT_RESOURCES) := nil;
          ERROR_HANDLERS(TASK_IS_NOT_ACTIVE) := nil;
          ERROR_HANDLERS(TASK_ALREADY_SUSPENDED) := nil;
          ERROR_HANDLERS(ILLEGAL_TASK_ID) := nil;
          ERROR_HANDLERS(ILLEGAL_OPERATION) := nil;
        end; (* with...do *)

      (* Create the control task *)
      CREATE(CONTROL, TASK_ATTR, CONTROL_TASK);

      (* Inquire as to the number of philosophers to simulate *)
      write('Number of guests > ');
      reset(input);
      TEMP := 0;
      repeat
        CH := readkey;
        write(CH);
        if CH in '0..'9' then
          TEMP := TEMP * 10 + (ord(CH) - ord('0'));
        until not (CH in '0..'9');
        if (CH = CR) and (TEMP <= MAX_NUMBER_OF_PHILOSOPHERS)
        then
          if TEMP < 0
          then
            NUMBER_OF_PHILOSOPHERS := TEMP
          else (* Remember Pascal's dangling else problem *)
            begin
              writeln(CR, LF, 'Error' 'Invalid digit or number');
              halt(1);
            end; (* if...then...else *)
          end;

          clrscr;

          (* Initialize bookkeeping structures *)
          NUM_THINKING := NUMBER_OF_PHILOSOPHERS;
          NUM_EATING := 0;
          for LCV := 0 to (NUMBER_OF_PHILOSOPHERS - 1) do
            begin
              ACTIVITY[LCV] := THINKING;
              NEED_FORKS[LCV] := 0;
            end; (* for...to...do *)

          (* Modify the startup attributes for philosophers *)
          with TASK_ATTR do
            begin
              PRIORITY := 10;
              STACK_WORDS_NEEDED := 500;
            end; (* with...do *)

          (* Create the philosopher tasks *)
          for LCV := 1 to NUMBER_OF_PHILOSOPHERS do
            CREATE(GUEST, TASK_ATTR, PHILOSOPHER_TASK);
          end; (* DINING_PHILOSOPHERS *)
        end;
      end;
    end;
  end;
end;

```

# 10.2 TSK-BNCH.PAS

```

program TASKING_BENCHMARK(input, output);
[*****]
[* This program executes a fixed TASKING application in a controlled manner *]
[* in an attempt to characterize the performance, and therefore aid in the *]
[* selection of parameters, for TASKING. *]
[*****]
Compiler Options (Ver. 7.0)
[*SA+ Word Alignment *]
[*SB- Short Circuit Boolean Evaluation *]
[*SD+ Debug Code Generation ON (Sort of) *]
[*SL+ Requires /V option to TPC to activate *]
[* Local Debug symbols ON (Sort of) *]
[* Requires /V option to TPC to activate *]
[*SF- Far calls only as needed *]
[*SI- I/O Checking OFF *]
[*SM $8000,50000,100000 *]
[*$M+ Memory (Stack, Minheap, Maxheap) *]
[*$H+ Hardware $0x8? Required *]
[*$E+ $0x8? Run-Time Emulation Allowed *]
[*$O- Overlays NOT allowed *]
[*$Q+ Overflow Checking ON *]
[*$R+ Range Checking ON *]
[*$S+ Stack Checking ON *]
[*$V+ Var-string Checking ON *]
[*$X+ Disable Extended syntax *]
[*****]
uses dos, crt, graph, GRAPHICS, BMP_UTIL;

const [*****]
[* Duration parameter to TASK-A.EXE, controls the length of execution *]
[* and therefore the resolution of the benchmark. In general, the *]
[* longer the better (minimum reliable parameter is -20). *]
[*****]
MIN_DURATION = 2;
DURATION integer = MIN_DURATION;

[* Flag to allow data files to be created automatically (without *]
[* prompts to the user). *]
[*****]
AUTOMATIC_MODE : boolean = false;

[*****]
[* Number of data elements to collect, TASKING timeslice will be *]
[* from 1 mSec to this parameter (in mSec). *]
[*****]
MAX_DATA = 250;

[*****]
[* Graphic screen offsets and control points. *]
[*****]
TITLE OFFSET = 26;
MIN_LINE_X = 24;
MIN_LINE_Y = 5;

[*****]
[* Basic screen colors. *]
[*****]
BACKGROUND_COLOR = white;
BORDER_COLOR = green;
TITLE_COLOR = blue;
LEGEND_COLOR = red;
AXIS_COLOR = darkgray;
GRID_COLOR = lightgray;
SCALE_COLOR = darkgray;

[*****]
[* General screen output control characters. *]
[*****]
CR = chr($0D);
LF = chr($0A);

type [*****]
[*****]
GRAPHS =
(
RELATIVE_PERFORMANCE,
CONTEXT_SWITCH_OVERHEAD,
CONTEXT_SWITCH_TIME,
CONTEXT_SWITCH_TIME_AVG,
GENERAL_OVERHEAD,
CONTEXT_SWITCH_PERCENTAGE
);

[* This data structure is used to hold all of the TASKING data for a *]
[* particular run of TASKS-A.EXE. *]
[*****]
DATA_RECORD = record
CONTEXT_SWITCHES : longint;
ABSOLUTE_TIME : double;
EFFECTIVE_TIME : double;
ACTUAL_TIMESLICE : double;
POINTS : array [GRAPHS] of double;
end; (* DATA_RECORD *)

const [*****]
[* Colors of graphs. *]
[*****]
GRAPH_COLOR : array [GRAPHS] of word =
(
(Relative Performance) red,
(Context Switch Overhead) lightmagenta,
(Context Switch Time) blue,
(Context Switch Time Average) lightblue,
(General Overhead) green,
(Context Switch Percentage) black
);

[*****]
[* Legends of graphs. *]
[*****]
LEGENDS : array [GRAPHS] of string[30] =
(
(Relative Performance) 'Relative Performance',
(Context Switch Overhead) 'Context Switch Overhead',
(Context Switch Time) 'Context Switch Time',
(Context Switch Time Average) 'Context Switch Time Average',
(General Overhead) 'General Overhead',
(Context Switch Percentage) 'Context Switch Percentage'
);

[*****]
[* Y-Axis Labels of graphs. *]
[*****]
LABELS : array [GRAPHS] of string[20] =
(
(Relative Performance) 'Percent',
(Context Switch Overhead) 'Percent',
(Context Switch Time) 'Microseconds',
(Context Switch Time Average) 'Microseconds',
(General Overhead) 'Percent',

```

```

(Context Switch Percentage) 'Percent'
);
[*****]
[*****]
[* BMP filename of graphs. *]
[*****]
FILENAME : array [GRAPHS] of string[20] =
(
(Relative Performance) 'REL-PERF.BMP',
(Context Switch Overhead) 'CTX-OVHD.BMP',
(Context Switch Time) 'CTX-TIME.BMP',
(Context Switch Time Average) 'CTX-TAVG.BMP',
(General Overhead) 'GEN-OVHD.BMP',
(Context Switch Percentage) 'CTS-PRCT.BMP'
);

var [*****]
[* Collection of data records for all runs of TASKS-A.EXE that will *]
[* be made. *]
[*****]
DATA : array[1..MAX_DATA] of DATA_RECORD;

[*****]
[* Baseline data from the runs of TASKS-A.EXE and TASKS-B.EXE. *]
[*****]
BASE_ABSOLUTE_TIME : double;
MINIMUM_ABSOLUTE_TIME : double;
MINIMUM_EFFECTIVE_TIME : double;
AVERAGE_SWITCH_TIME : double;

[*****]
[* Used to convert command line parameters from strings to integers. *]
[*****]
COMMAND_LINE : string; CODE : integer;

procedure GENERATE_RAW_DATA(DURATION : integer);
[*****]
[* This routine generates all of the TASKING.RPT files by executing the *]
[* benchmark program with the specified duration parameter. *]
[*****]
var [*****]
[* Loop Control Variable used to generate all of the TASKING.RPT data *]
[* files. *]
[*****]
LCV : longint;

[*****]
[* Used to rename the TASKING.RPT file to the specific data file. *]
[*****]
LCV_STR : string(9);

[*****]
[* Text file for parsing the baseline TASKING.RPT files. *]
[*****]
REPORT : text;
REPORT_NAME : string(11);

[*****]
[* Used to pass the duration parameter to the subprogram command line. *]
[*****]
DUR_STR : string(5);

[*****]
[* Time variables used to compute the overhead associated with the *]
[* subprogram execution (so that it does not affect the results). *]
[*****]
HOUR, MIN, SEC, SEC100 : word;
HOUR_2, MIN_2, SEC_2, SEC100_2 : word;
ELAPSED_TIME : double;
OVERHEAD, TEMP : longint;

[*****]
[* Command line sent to DOS to execute the subprograms. *]
[*****]
COMMAND_LINE : string(80);

[*****]
[* User response to command prompts. *]
[*****]
ANSWER : char;

[*****]
[* String buffers to parse the TASKING.RPT file. *]
[*****]
BUFFER, LINE : string;

[*****]
[* Flag used to generate all data files automatically. *]
[*****]
GENERATE_ALL : boolean;

begin (* GENERATE_RAW_DATA *)
GENERATE_ALL := false;
str(DURATION, DUR_STR);

write(CR, ' Base Time = ? > ');
if AUTOMATIC_MODE
then
COMMAND_LINE := ''
else
readln(COMMAND_LINE);
val(COMMAND_LINE, BASE_ABSOLUTE_TIME, CODE);
if CODE <> 0 then
begin
[* Calculate 'exec' overhead *]
write(CR, ' Base Time = ?(:crlf);
OVERHEAD := 0;
for LCV := 1 to 3 do
repeat
gettime(HOUR, MIN, SEC, SEC100);
exec(getenv('COMSPEC'), '/C');
gettime(HOUR_2, MIN_2, SEC_2, SEC100_2);
TEMP := (longint(HOUR_2) - longint(HOUR)) * 360000;
TEMP := TEMP + (longint(MIN_2) - longint(MIN)) * 6000;
TEMP := TEMP + (longint(SEC_2) - longint(SEC)) * 100;
TEMP := TEMP + (longint(SEC100_2) - longint(SEC100)) - OVERHEAD;
if TEMP > OVERHEAD then OVERHEAD := TEMP;
until OVERHEAD < 0;

[* Calculate 'base' time *]
COMMAND_LINE := DUR_STR + ' -PREMPT -TIMESLICE 60000000';
gettime(HOUR, MIN, SEC, SEC100);
exec('TASKS-B.EXE', COMMAND_LINE);
gettime(HOUR_2, MIN_2, SEC_2, SEC100_2);
ELAPSED_TIME := (longint(HOUR_2) - longint(HOUR)) * 60 * 60;
ELAPSED_TIME := ELAPSED_TIME + (longint(MIN_2) - longint(MIN)) * 60;
ELAPSED_TIME := ELAPSED_TIME + (longint(SEC_2) - longint(SEC));
ELAPSED_TIME := ELAPSED_TIME + (longint(SEC100_2) - longint(SEC100))
OVERHEAD) / 100;

BASE_ABSOLUTE_TIME := ELAPSED_TIME;
end; (* if...then *)
writeln(CR, ' Base Time = ', BASE_ABSOLUTE_TIME 6 : 2, ' sec'(:crlf);

write(CR, ' Minimum Absolute Time = ? > ');
if AUTOMATIC_MODE
then
COMMAND_LINE := ''
else

```

```

    readln(COMMAND_LINE);
val (COMMAND_LINE, MINIMUM_ABSOLUTE_TIME, CODE);
if CODE <> 0 then
begin
    (* Calculate Minimum absolute time *)
    write(CR, ' Minimum Absolute Time = ?'); clreol;
    COMMAND_LINE := DUR_STR + ' -PREEMPT -TIMESLICE 6000000';
    exec('TASKS-A.EXE', COMMAND_LINE);
    assign(REPORT, 'TASKING.RPT');
    (*$I-*) reset(REPORT); (*$I+*)
    if ioreult = 0 then
    begin
        while not eof(REPORT) do
        begin
            readln(REPORT, LINE);
            if pos('Absolute =', LINE) <> 0 then
            begin
                BUFFER := copy(LINE, pos('=', LINE) + 1, pos('seconds',
                    LINE) - pos('=', LINE) - 2);
                val (BUFFER, MINIMUM_ABSOLUTE_TIME, CODE);
                end; (* if...then *)
            end; (* while...do *)
            close(REPORT);
        end; (* if...then *)
        end; (* if...then *)
    write(CR, ' Minimum Absolute Time = ', MINIMUM_ABSOLUTE_TIME : 6 : 2,
        ' sec'); clreol;
    write(CR, ' Minimum Effective Time = ?');
    if AUTOMATIC_MODE
    then
        COMMAND_LINE := ''
    else
        readln(COMMAND_LINE);
    val (COMMAND_LINE, MINIMUM_EFFECTIVE_TIME, CODE);
    if CODE <> 0 then
    begin
        (* Calculate Minimum effective time *)
        write(CR, ' Minimum Effective Time = ?'); clreol;
        assign(REPORT, 'TASKING.RPT');
        (*$I-*) reset(REPORT); (*$I+*)
        if ioreult <> 0 then
        begin
            COMMAND_LINE := DUR_STR + ' -PREEMPT -TIMESLICE 6000000';
            exec('TASKS-A.EXE', COMMAND_LINE);
            assign(REPORT, 'TASKING.RPT');
            (*$I-*) reset(REPORT); (*$I+*)
            end; (* if...then *)
            if ioreult = 0 then
            begin
                while not eof(REPORT) do
                begin
                    readln(REPORT, LINE);
                    if pos('Effective =', LINE) <> 0 then
                    begin
                        BUFFER := copy(LINE, pos('=', LINE) + 1, pos('seconds',
                            LINE) - pos('=', LINE) - 2);
                        val (BUFFER, MINIMUM_EFFECTIVE_TIME, CODE);
                        end; (* if...then *)
                    end; (* while...do *)
                    close(REPORT);
                end; (* if...then *)
                end; (* if...then *)
            write(CR, ' Minimum Effective Time = ', MINIMUM_EFFECTIVE_TIME : 6 : 3,
                ' sec'); clreol;
            assign(REPORT, 'BASELINE');
            rewrite(REPORT);
            write(CR, ' DURATION');
            write(CR, ' BASE ABSOLUTE TIME : 6 : 2);
            write(CR, ' MINIMUM ABSOLUTE TIME : 6 : 2);
            write(CR, ' MINIMUM EFFECTIVE TIME : 6 : 3);
            close(REPORT);
        for LCV := MAX_DATA downto 1 do
        begin
            if keypressed then
            begin
                write('User break!');
                halt(1);
            end; (* if...then *)
            str(LCV, LCV_STR);
            REPORT_NAME := 'TASKING.' + LCV_STR;
            assign(REPORT, REPORT_NAME);
            (*$I-*) reset(REPORT); (*$I+*)
            if ioreult <> 0
            then
            begin
                if AUTOMATIC_MODE then
                    GENERATE_ALL := true;
                if not GENERATE_ALL then
                repeat
                    write(CR, ' File not found: (' , REPORT_NAME,
                        ') , (A)bot, (C)alculate, (S)kip, (G)enerate all > ');
                    ANSWER := upcase(readkey);
                    GENERATE_ALL := ANSWER = 'G';
                    until ANSWER in ('C', 'S', 'G', 'A');
                    if ANSWER = 'A' then
                    begin
                        write(CR, ' User break!');
                        halt(1);
                    end; (* if...then *)
                    if GENERATE_ALL or (ANSWER = 'C') then
                    begin
                        COMMAND_LINE := DUR_STR + ' -PREEMPT -TIMESLICE ' + LCV_STR;
                        write(CR, ' Generating Data File (' , LCV, ')'); clreol;
                        swapvectors;
                        exec('TASKS-A.EXE', COMMAND_LINE);
                        exec(getenv('COMSPEC'), '/C " + 'RENAME TASKING.RPT ' +
                            REPORT_NAME + ' > NUL"');
                        swapvectors;
                        write(CR); clreol;
                        end; (* if...then *)
                    end; (* if...then *)
                else
                    close(REPORT);
                end; (* for...to...do *)
            write(CR, ' Generated Data File(s)'); clreol;
            write(CR, ' Generated RAW DATA ');
            end; (* GENERATE_RAW_DATA *)
        procedure PARSE_RAW_DATA;
        (* Parses the TASKING.* data files for the performance information. *)
        var
        (* Search data structure for finding TASKING.* files. *)
        SEARCH : searchrec;
        (* Loop Control Variable to access all data files. *)
        LCV : integer;
        (* Text file for parsing the TASKING.RPT files. *)
        REPORT : text;
        (* String buffers to parse the TASKING.RPT file. *)
        BUFFER, LINE : string;
        (* Parameters that are extracted from the TASKING.* files. *)
        TIMESLICE : longint;
        SWITCHES : longint;
        TIME1, TIME2 : double;
        (* Used to convert strings to numbers. *)
        CODE : word;
        begin (* PARSE_RAW_DATA *)
            write(' Parsing Data Files');
            for LCV := 1 to MAX_DATA do
            with DATA(LCV) do
            begin
                CONTEXT_SWITCHES := 0;
                ACTUAL_TIMESLICE := 0;
                POINTS[CONTEXT_SWITCH_TIME_AVG] := 0;
                end; (* with...do *)
                findfirst('TASKING.*', anyfile, SEARCH);
                while doserror = 0 do
                begin
                    TIMESLICE := 0;
                    assign(REPORT, SEARCH.NAME);
                    (*$I-*) reset(REPORT); (*$I+*)
                    if ioreult = 0 then
                    begin
                        write(CR, ' Parsing Data Files: (' , SEARCH.NAME, ')'); clreol;
                        while not eof(REPORT) do
                        begin
                            readln(REPORT, LINE);
                            if pos('Achieved Time Slice =', LINE) <> 0 then
                            begin
                                BUFFER := copy(LINE, pos('=', LINE) + 3, pos('usec', LINE)
                                    - pos('=', LINE) - 4);
                                val (BUFFER, TIMESLICE, CODE);
                                TIMESLICE := round(TIMESLICE / 1000);
                                end;
                                if pos('Context Switches =', LINE) <> 0 then
                                begin
                                    BUFFER := copy(LINE, pos('=', LINE) + 2, pos('(', LINE)
                                        - pos('=', LINE) - 3);
                                    val (BUFFER, SWITCHES, CODE);
                                    end;
                                    if pos('Effective =', LINE) <> 0 then
                                    begin
                                        BUFFER := copy(LINE, pos('=', LINE) + 1, pos('seconds',
                                            LINE) - pos('=', LINE) - 2);
                                        val (BUFFER, TIME1, CODE);
                                        end;
                                        if pos('Absolute =', LINE) <> 0 then
                                        begin
                                            BUFFER := copy(LINE, pos('=', LINE) + 1, pos('seconds',
                                                LINE) - pos('=', LINE) - 2);
                                            val (BUFFER, TIME2, CODE);
                                            end;
                                            end; (* while...do *)
                                            if (TIMESLICE <= MAX_DATA) and (TIMESLICE <> 0)
                                            then
                                            if DATA(TIMESLICE).CONTEXT_SWITCHES <> 0
                                            then
                                            with DATA(TIMESLICE) do
                                            begin
                                                (* Duplicate Data Record *)
                                                CONTEXT_SWITCHES := (CONTEXT_SWITCHES + SWITCHES) div
                                                    2;
                                                EFFECTIVE_TIME := (EFFECTIVE_TIME + TIME1) / 2;
                                                ABSOLUTE_TIME := (ABSOLUTE_TIME + TIME2) / 2;
                                                end (* if...then *)
                                                else
                                                with DATA(TIMESLICE) do
                                                begin
                                                    CONTEXT_SWITCHES := SWITCHES;
                                                    EFFECTIVE_TIME := TIME1;
                                                    ABSOLUTE_TIME := TIME2;
                                                    end (* with...do *)
                                                end
                                                else
                                                begin
                                                    (* Timeslice out of bounds! Ignore it! *)
                                                    end; (* if...then...else *)
                                                close(REPORT);
                                                end; (* if...then *)
                                                findnext(SEARCH);
                                                end; (* while...do *)
                                            write(CR, ' Parsed Data Files'); clreol;
                                            end; (* PARSE_RAW_DATA *)
        procedure CALCULATE_PERFORMANCE_CURVES;
        (* Calculates the performance values for all of the TASKING.* files. *)
        var
        (* Loop Control Variable to access all data files. *)
        LCV : integer;
        (* Number of elements that are included in the average. *)
        NUM_AVG : integer;
        (* Temporary data records for creating the curves. *)
        PREVIOUS, FIRST : DATA_RECORD;
        (* Parsed data output file (for importing into a spreadsheet). *)
        PARSE_FILE : text;
        begin (* CALCULATE_PERFORMANCE_CURVES *)
            write(' Calculating Performance Curves');
            FIRST.CONTEXT_SWITCHES := 0;
            assign(PARSE_FILE, 'PARSE.OUT');
            rewrite(PARSE_FILE);
            for LCV := 1 to MAX_DATA do
            if DATA(LCV).CONTEXT_SWITCHES <> 0 then
            with DATA(LCV) do
            begin
                write(CR, ' Calculating Performance Curves: (TASKING.', LCV, ')');
                clreol;
                ACTUAL_TIMESLICE := ABSOLUTE_TIME * 1000 / CONTEXT_SWITCHES;
                POINTS[RELATIVE_PERFORMANCE] :=
                    (BASE_ABSOLUTE_TIME - ABSOLUTE_TIME) /
                    (BASE_ABSOLUTE_TIME * 100;
                POINTS[CONTEXT_SWITCH_TIME] := abs(EFFECTIVE_TIME * 1e6 -
                    MINIMUM_EFFECTIVE_TIME * 1e6) /

```

```

CONTEXT_SWITCHES := POINTS(CONTEXT_SWITCH_TIME) /
POINTS(CONTEXT_SWITCH_TIME_AVG) := POINTS(CONTEXT_SWITCH_TIME) /
POINTS(CONTEXT_SWITCH_OVERHEAD) := 100 -
((EFFECTIVE_TIME * 1e6 /
CONTEXT_SWITCHES
) - POINTS(CONTEXT_SWITCH_TIME)) /
(EFFECTIVE_TIME * 1e6 /
CONTEXT_SWITCHES) * 100;
POINTS(GENERAL_OVERHEAD) := POINTS(RELATIVE_PERFORMANCE) +
POINTS(CONTEXT_SWITCH_OVERHEAD);
POINTS(CONTEXT_SWITCH_PERCENTAGE) :=
(POINTS(CONTEXT_SWITCH_TIME) / 1000) /
ACTUAL_TIMESLICE * 100;
if FIRST.CONTEXT_SWITCHES = 0 then
FIRST := DATA[LCV];
writeln(PARSE_FILE,
ACTUAL_TIMESLICE : 3 : 3, ',',
POINTS(RELATIVE_PERFORMANCE) : 3 : 3, ',',
POINTS(CONTEXT_SWITCH_OVERHEAD) : 3 : 3, ',',
POINTS(CONTEXT_SWITCH_TIME) : 3 : 3, ',',
POINTS(GENERAL_OVERHEAD) : 3 : 3, ',',
POINTS(CONTEXT_SWITCH_PERCENTAGE) : 3 : 3
);
end; (* with...do *)
close(PARSE_FILE);
PREVIOUS := FIRST;
for LCV := 1 to MAX_DATA do
with DATA[LCV] do
with DATA[LCV] do
if (CONTEXT_SWITCHES < 0) then
begin
POINTS(CONTEXT_SWITCH_TIME_AVG) :=
(PREVIOUS.POINTS(CONTEXT_SWITCH_TIME_AVG) +
POINTS(CONTEXT_SWITCH_TIME_AVG)) / 2;
PREVIOUS := DATA[LCV];
end; (* if...then *)
NUM_AVG := 0;
AVERAGE_SWITCH_TIME := 0;
for LCV := 1 to MAX_DATA do
with DATA[LCV] do
if (CONTEXT_SWITCHES < 0) then
begin
AVERAGE_SWITCH_TIME := AVERAGE_SWITCH_TIME +
POINTS(CONTEXT_SWITCH_TIME_AVG);
inc(NUM_AVG);
end; (* if...then *)
if NUM_AVG > 0 then
AVERAGE_SWITCH_TIME := AVERAGE_SWITCH_TIME / NUM_AVG;
write(CR, ' Calculated Performance Curves'); clrscr;
writeln;
end; (* CALCULATE PERFORMANCE CURVES *)
procedure GRAPH_DATA;
(*.....*)
(* This routine creates the axis and draws the graphs on the screen. The *)
(* user is allowed only two options 'Q' to quit the program and 'G' to grab *)
(* the screen image as a Windows Bit-Map (BMP) file. *)
(*.....*)
const (*.....*)
(* The length of the legend lines. *)
LEGEND_LENGTH = 30;
var (*.....*)
(* Graphic video driver parameters. *)
MODE, DRIVER : integer;
(*.....*)
(* Loop Control Variable to process the data points. *)
LCV : longint;
(*.....*)
(* Graph critical screen positions. *)
MAX_LINE_X, MAX_LINE_Y, MID_LINE_Y, MID_LINE_X : integer;
MAX_X, MAX_Y, MID_X, MID_Y : integer;
(*.....*)
(* Flag used to detect the first point to be graphed. *)
STARTED : boolean;
(*.....*)
(* Flag used to allow the user to exit the program. *)
USER_QUIT : boolean;
(*.....*)
(* Used to create the appropriate 'tick' marks on the graph axis. *)
SCALE_FACTOR : integer;
THE_GRAPH : GRAPHS;
procedure DRAW_GRAPH (GRAPH : GRAPHS);
(*.....*)
(* This routine draws the graph onto the screen. *)
(*.....*)
var (*.....*)
(* String used to display text on the screen. *)
TEMP_STR : string;
(*.....*)
(* Loop Control Variable to access all data points. *)
LCV : integer;
(*.....*)
(* Graph scales are dynamically calculated based on the data to be *)
(* displayed by the graph. *)
MAX_Y_SCALE, MIN_Y_SCALE : double;
MAX_X_SCALE : integer;
MID_Y_SCALE : integer;
function REAL_STR(R: real; F1, F2 : integer): string;
(*.....*)
(* Convert real type to a string. *)
(*.....*)
var S : string[11];
begin (* REAL_STR *)
str(R : F1 : F2, S);
REAL_STR := S;
end; (* REAL_STR *)
begin (* DRAW_GRAPH *)
setgraphmode(MODE);
MAX_X := getmaxx;
MAX_Y := getmaxy;
MID_X := MAX_X div 2;
MID_Y := MAX_Y div 2;
MAX_LINE_X := MAX_X - 3 * MIN_LINE_X;
MAX_LINE_Y := MAX_Y - 2 * TITLE_OFFSET - 5;
MID_LINE_X := (MAX_LINE_X + MIN_LINE_X) div 2;
MID_LINE_Y := (MAX_LINE_Y + MIN_LINE_Y) div 2;
cleardevice;
(setcolor(BORDER_COLOR);
setlinestyle(solidln, 0, thickwidth);
rectangle(0, 0, getmaxx, getmaxy);
setfillstyle(solidfill, BACKGROUND_COLOR);
floodfill(10, 10, BORDER_COLOR);
setcolor(BORDER_COLOR);
lineto(0, MAX_Y);
lineto(MAX_X, MAX_Y);
lineto(MAX_X, 0);
lineto(0, 0);
line(0, TITLE_OFFSET, MAX_X, TITLE_OFFSET);
line(0, MAX_Y - TITLE_OFFSET, MAX_X, MAX_Y - TITLE_OFFSET);
settextstyle(SIMPLEX_FONT, horzdir, 1);
settextjustify(centertext, centertext);
setcolor(TITLE_COLOR);
outtextxy
(
MAX_X div 2,
TITLE_OFFSET div 3,
'TASKING Performance Benchmark Display (Q=Quit; G=Grab; N=Next)'
);
(* Create Legend *)
setcolor(LEGEND_COLOR);
moveto(MAX_X div 2, MAX_Y - 2 * (TITLE_OFFSET div 3));
outtext(LEGENDS[GRAPH]);
moverel(length(LEGENDS[GRAPH]) * textwidth('h') div 2, 4);
setcolor(GRAPH_COLOR[GRAPH]);
setlinestyle(solidln, 0, thickwidth);
linereel(LEGEND_LENGTH, 0);
setcolor(SCALE_COLOR);
settextjustify(centertext, centertext);
settextstyle(defaultfont, verdir, 1);
outtextxy(12, MAX_Y div 2, LABELS[GRAPH]);
settextstyle(defaultfont, horzdir, 1);
outtextxy
(
(MAX_LINE_X - 2 * MIN_LINE_X) div 2,
(MAX_Y div 2) + 18,
'Actual Timeslice (mSec)'
);
setlinestyle(solidln, 0, normwidth);
setcolor(AXIS_COLOR);
setviewport
(
MIN_LINE_X - 1, TITLE_OFFSET,
MAX_X,
MAX_Y - TITLE_OFFSET,
false
);
line(MIN_LINE_X, MIN_LINE_Y, MIN_LINE_X, MAX_LINE_Y);
line(MAX_LINE_X, MIN_LINE_Y, MAX_LINE_X, MAX_LINE_Y);
line(MIN_LINE_X, MID_LINE_Y, MAX_LINE_X, MID_LINE_Y);
MAX_Y_SCALE := -maxint;
MIN_Y_SCALE := +maxint;
MAX_X_SCALE := 0;
for LCV := 1 to MAX_DATA do
with DATA[LCV] do
with DATA[LCV] do
if (CONTEXT_SWITCHES < 0) then
begin
if ACTUAL_TIMESLICE > MAX_X_SCALE then
MAX_X_SCALE := round(ACTUAL_TIMESLICE + 0.5);
if POINTS[GRAPH] < MIN_Y_SCALE then
MIN_Y_SCALE := round(POINTS[GRAPH] - 0.5);
if POINTS[GRAPH] > MAX_Y_SCALE then
MAX_Y_SCALE := round(POINTS[GRAPH] + 0.5);
) * Make sure everything is even *)
if (MAX_X_SCALE mod 5) < 0 then
MAX_X_SCALE := ((MAX_X_SCALE div 5) + 1) * 5;
if odd(round(MIN_Y_SCALE)) then MIN_Y_SCALE := MIN_Y_SCALE - 1;
if odd(round(MAX_Y_SCALE)) then MAX_Y_SCALE := MAX_Y_SCALE + 1;
end; (* with...do *)
MID_Y_SCALE := round(MAX_Y_SCALE + MIN_Y_SCALE) div 2;
settextjustify(centertext, centertext);
) * Vertical axis *)
setcolor(AXIS_COLOR);
line
(
MIN_LINE_X - 5,
MID_LINE_Y - round(MID_Y_SCALE * MAX_LINE_Y /
(MAX_Y_SCALE - MIN_Y_SCALE)),
MIN_LINE_X + 5,
MID_LINE_Y - round(MID_Y_SCALE * MAX_LINE_Y /
(MAX_Y_SCALE - MIN_Y_SCALE))
);
(* 'Y' Ticks *)
for LCV := 0 to 19 do
begin
setcolor(GRID_COLOR);
setlinestyle(dottedln, 0, normwidth);
line
(
MIN_LINE_X - 5,
MAX_LINE_Y - round(LCV * MAX_LINE_Y / 20),
MAX_LINE_X + 5,
MAX_LINE_Y - round(LCV * MAX_LINE_Y / 20)
);
setcolor(SCALE_COLOR);
outtextxy
(
MIN_LINE_X - 4 - 2 * textwidth('H'),
MAX_LINE_Y - round(LCV * MAX_LINE_Y / 20),
REAL_STR(MIN_Y_SCALE + (MAX_Y_SCALE - MIN_Y_SCALE) / 20 *
LCV, 2, 1)
);
end; (* for...do *)
setcolor(AXIS_COLOR);
setlinestyle(solidln, 0, normwidth);
line
(
MIN_LINE_X + round(1 * MAX_LINE_X div (MAX_X_SCALE div 20)),
MID_LINE_Y + 5,
MIN_LINE_X + round(1 * MAX_LINE_X div (MAX_X_SCALE div 20)),
MID_LINE_Y - 5
);
(* 'X' Ticks *)
for LCV := 1 to (MAX_X_SCALE div 20) - 1 do
begin
setcolor(GRID_COLOR);

```

```

setlinestyle(dottedln, 0, normwidth);
line
  |
  | MIN_LINE_X + round(LCV * MAX_LINE_X div (MAX_X_SCALE div 20)),
  | MIN_LINE_Y + 5,
  | MIN_LINE_X + round(LCV * MAX_LINE_X div (MAX_X_SCALE div 20)),
  | MAX_LINE_Y - 5
  |;
setcolor(SCALE_COLOR);
str(LCV * 20, TEMP_STR);
outtextxy
  |
  | MIN_LINE_X + round(LCV * MAX_LINE_X div (MAX_X_SCALE div 20)),
  | MID_LINE_Y + 8,
  | TEMP_STR
  |;
end; /* for...to...do */

setcolor(GRAPH_COLOR[GRAPH]);
setlinestyle(solidln, 0, thickwidth);
STARTED := false;
for LCV := 1 to MAX_DATA do
  if DATA[LCV].CONTEXT_SWITCHES <> 0 then
    with DATA[LCV] do
      if STARTED
        then
          lineto
            |
            | MIN_LINE_X + round(ACTUAL_TIMESLICE / MAX_X_SCALE *
            | (MAX_LINE_X - MIN_LINE_X)),
            | MID_LINE_Y - round((POINTS[GRAPH] - MID_Y_SCALE) /
            | (MAX_Y_SCALE - MIN_Y_SCALE) *
            | (MAX_LINE_Y - MIN_LINE_Y))
            |;
        else
          begin
            moveto
              |
              | MIN_LINE_X + round(ACTUAL_TIMESLICE / MAX_X_SCALE *
              | (MAX_LINE_X - MIN_LINE_X)),
              | MID_LINE_Y - round((POINTS[GRAPH] -
              | MID_Y_SCALE) /
              | (MAX_Y_SCALE - MIN_Y_SCALE) *
              | (MAX_LINE_Y - MIN_LINE_Y))
              |;
            STARTED := true;
          end; /* if...then...else */
        end; /* DRAW_GRAPH */
      end;
    end;
  end;
end; /* GRAPH_DATA */

installuserdriver('VESAL6', @DETECT_VESA_16);
registerbgdriver(addr(VESAL6_DRIVER);
DRIVER := detect;
initgraph(DRIVER, MODE, '');

THE_GRAPH := RELATIVE_PERFORMANCE;
DRAW_GRAPH(THE_GRAPH);
USER_QUIT := false;
repeat
  if keypressed then
    case upcase(readkey) of
      'N' : begin
        if THE_GRAPH = high(GRAPHS)
          then
            THE_GRAPH := low(GRAPHS)
          else
            THE_GRAPH := succ(THE_GRAPH);
        end; /* Next Graph */
      'G' : begin
        setviewport(0, 0, getmaxx, getmaxy, true);
        SAVE_IMAGE_AS_16_COLOR_EMP_FILE
          |
          | FILENAME(THE_GRAPH),
          | 0, 0, getmaxx, getmaxy
          |;
        setviewport
          |
          | MIN_LINE_X - 1, TITLE_OFFSET,
          | MAX_X, MAX_Y - TITLE_OFFSET,
          | true
          |;
        end; /* Grab picture of graph */
      'Q' : USER_QUIT := true;
    end; /* case...of */
  until USER_QUIT;
closegraph;
end; /* GRAPH_DATA */

function UP_STRING(S : string) : string;
.....
/* This routine returns a string corresponding to the upper case value of */
/* the string passed to it. */
.....

var LCV : integer;

begin /* UP_STRING */
  UP_STRING[0] := S[0];
  for LCV := 1 to length(S) do
    UP_STRING[LCV] := upcase(S[LCV]);
  end; /* UP_STRING */

begin /* TASKING_BENCHMARK */
  if (paramcount = 1) and (UP_STRING(paramatr[1]) = '-AUTO') then
    AUTOMATIC_MODE := true;

  directvideo := false;
  /* Allow baseline info to be redirected from a file */
  assign(input, '');
  reset(input);

  writeln('TASKING Benchmarks');
  write(CR, 'Duration = ? ');clreol;
  readln(COMMAND_LINE);
  val(COMMAND_LINE, DURATION, CODE);
  if (CODE <> 0) or (DURATION < MIN_DURATION) then
    begin
      writeln(CR, LF, 'TSK-BNCH: Invalid Duration: ', COMMAND_LINE, '');
      halt();
    end; /* if...then */
  writeln(CR, 'Duration = ', DURATION;clreol;

  /* Perform benchmarking */
  GENERATE_RAW_DATA(DURATION);
  PARSE_RAW_DATA;
  CALCULATE_PERFORMANCE_CURVES;
  GRAPH_DATA;

  /* Echo results to the user */
  writeln;
  writeln('TASKING Benchmarks (Duration = ', DURATION, ')');
  writeln('Base Time = ', BASE_ABSOLUTE_TIME : 6 : 2, ' sec');
  writeln('Minimum Absolute Time = ', MINIMUM_ABSOLUTE_TIME : 6 : 2,
  ' sec');
  writeln('Minimum Effective Time = ', MINIMUM_EFFECTIVE_TIME : 6 : 3,
  ' sec');
  writeln('Average Context Switch Time = ', AVERAGE_SWITCH_TIME : 5 : 1,

```



## 10.3 TASKS-A.PAS

```

program TASKS(input, output);
(*.....*)
(* This program is merely a sample tasking program to illustrate some of *)
(* the capabilities and benefits of the TASKING unit. *)
(*.....*)
(* Compiler Options (Ver. 7.0) *)
(*$A+ Word Alignment *)
(*$B- Short Circuit Boolean Evaluation *)
(*$D+ Debug Code Generation ON (Sort of) *)
(*$L+ Requires /V option to TPC to activate *)
(*$M+ Local Debug symbols ON (Sort of) *)
(*$F- Requires /V option to TPC to activate *)
(*$I- Far calls only as needed *)
(*$J- I/O checking OFF *)
(*$K SM $8000,20000,50000 Memory (Stack, Minheap, Maxheap) *)
(*$N- Software Emulation of 80x87 *)
(*$O+ No 80x87 run-time emulation *)
(*$P- Overlays NOT allowed *)
(*$R- Range Checking OFF *)
(*$S- Stack Checking OFF *)
(*$T- Force Typed '8' references *)
(*$V- Var-string Checking OFF *)
(*.....*)

uses crt, dos, TASKING;

type
(*.....*)
(* This application is so simple that the commands can be as simple *)
(* as an enumerated data type. In a 'real' application the commands *)
(* passed around would probably be a record to allow command parameters *)
(* to be passed. The TASKING mechanisms would be exactly the same in *)
(* either case. There is no reason that message passing be restricted *)
(* to strict data types, strings (or any data type) can be passed just *)
(* as an easily. *)
(*.....*)
APPLICATION_COMMANDS =
(
TERMINATE_APPLICATION,
START_COMPUTATIONS
);

var
(*.....*)
(* These are the task identifiers used in the application. *)
SUPERVISOR, KYBD_TASK : TASK_IDS;

(*.....*)
(* Generic task attribute variable used to create all tasks. *)
TASK_ATTR : TASK_ATTRIBUTES;

(*.....*)
(* For command line timeslice string conversion. *)
(*.....*)
DURATION, CODE : integer; VALUE : longint;

(*.....*)
LCV : integer;

(*.....*)
BUSY_WORK : TASK_IDS;

procedure BUSY_TASK(TASK_ID : TASK_IDS; PRIORITY : USER_PRIORITIES); far;
(*.....*)
(* Busy work, no real purpose except for demonstration. *)
(*.....*)

var LCV : integer;
    DELAY_COUNT : longint;
    CPU_LOAD : longint;
    CH : char;

begin (* BUSY_TASK *)
    LCV := 0;
    repeat
        inc(LCV);
        (* Do some non-blocking work *)
        for DELAY_COUNT := 1 to TASK_ID do
            for CPU_LOAD := 1 to 500000 do (* Nothing really *) CH := 'A';
        until LCV = DURATION;
    end; (* BUSY_TASK *)

function UP_STRING(S : string) : string;
(*.....*)
(* This routine returns a string corresponding to the upper case value of *)
(* the string passed to it. *)
(*.....*)

var LCV : integer;

begin (* UP_STRING *)
    UP_STRING[0] := S[0];
    for LCV := 1 to length(S) do
        UP_STRING[LCV] := upcase(S[LCV]);
    end; (* UP_STRING *)

begin (* TASKS *)
    if paramcount = 0
    then
        begin
            writeln('TASKS-A: Duration must be specified');
            halt(1);
        end
    else
        val(paramstr(1), DURATION, CODE);

    if paramcount > 1 then
        if UP_STRING(paramstr(2)) = '-PREEMPT'
        then
            TASKING_CONFIGURATION.TASKING_MODEL := PREEMPTIVE
        else
            begin
                writeln('TASKS: Invalid parameter: {', paramstr(1), '}');
                halt(1);
            end; (* if..then..else *)
        if paramcount > 3 then
            if UP_STRING(paramstr(3)) = '-TIMESLICE'
            then
                begin
                    val(paramstr(4), VALUE, CODE);
                    if CODE < 0 then
                        begin
                            writeln('TASKS: Invalid number: {', paramstr(3), '}');
                            halt(1);
                        end; (* if..then *)
                    TASKING_CONFIGURATION.TARGET_TIMESLICE := VALUE * 1000;
                end (* if..then *)
            else
                begin
                    writeln('TASKS: Invalid parameter: {', paramstr(3), '}');

```

```

        halt(1);
    end; (* if..then..else *)

(* A good idea during application development, bad idea after that *)
TASKING_CONFIGURATION.STATISTICS := TASK_STATISTICS;

with TASK_ATTR do
begin
    PRIORITY := 10;
    STACK_WORDS_NEEDED := 250;
    (* Use default error handlers, carried through all task creations *)
    ERROR_HANDLERS(TASK_ALREADY_ACTIVE) := nil;
    ERROR_HANDLERS(INSUFFICIENT_RESOURCES) := nil;
    ERROR_HANDLERS(TASK_IS_NOT_ACTIVE) := nil;
    ERROR_HANDLERS(TASK_ALREADY_SUSPENDED) := nil;
    ERROR_HANDLERS(ILLEGAL_TASK_ID) := nil;
    ERROR_HANDLERS(ILLEGAL_OPERATION) := nil;
end; (* with...do *)

for LCV := 1 to 10 do
    CREATE(BUSY_WORK, TASK_ATTR, BUSY_TASK);
end. (* TASKS *)

```

## 10.4 TASKS-B.PAS

```

program TASKS(input, output);
{.....}
(* This program is merely a sample tasking program to illustrate some of
 * the capabilities and benefits of the TASKING unit.
 *.....)
{
  Compiler Options (Ver. 7.0)
  (*SA+ Word Alignment *)
  (*SB- Short Circuit Boolean Evaluation *)
  (*SD+ Debug Code Generation ON (Sort of) *)
  (*SL+ Requires /V option to TPC to activate Local Debug symbols ON (Sort of) *)
  (*SF- Requires /V option to TPC to activate Far calls only as needed *)
  (*SI- I/O Checking OFF *)
  (*SM $8000,20000,50000 Memory (Stack, Minheap, Maxheap) *)
  (*SN- Software Emulation of 80x87 *)
  (*SE- No 80x87 run-time emulation *)
  (*SO- Overlays NOT allowed *)
  (*SR- Range Checking OFF *)
  (*SS- Stack Checking OFF *)
  (*ST- Force Typed 'g' references *)
  (*SV- Var-string Checking OFF *)
}
uses crt, dos;

var {.....}
(* Loop Control Variable to create all 'tasks'. *)
LCV integer;
{.....}
(* Used to convert command line arguments to numbers. *)
DURATION, CODE integer; VALUE longint;

procedure BUSY_TASK(TASK_ID : word; PRIORITY : byte); far;
(* Busy work, no real purpose except for demonstration. *)
{.....}

var LCV : integer;
    DELAY_COUNT : longint;
    CPU_LOAD : longint;
    CH : char;

begin (* BUSY_TASK *)
  LCV := 0;
  repeat
    inc(LCV);
    (* Do some non-blocking work *)
    for DELAY_COUNT := 1 to TASK_ID do
      for CPU_LOAD := 1 to 50000 do (* Nothing really *) CH := 'A';
    until LCV = DURATION;
  end; (* BUSY_TASK *)

begin (* TASKS *)
  if paramcount = 0
  then
    begin
      writeln('TASKS-A: Duration must be specified');
      halt(1);
    end
  else
    val(paramstr(1), DURATION, CODE);

  for LCV := 1 to 10 do
    BUSY_TASK(LCV, 10);
  end. (* TASKS *)

```

## 10.5 GRAPHICS.PAS

```

unit GRAPHICS;
{.....}
(* This unit concentrates the code for VGA (640x480), VESA-16 (1024x768) *)
(* graphics drivers and useful fonts into a single code segment. *)
{.....}
{
  Compiler Options (Ver. 7.0)
  (*$A+ Word Alignment *)
  (*$B- Short Circuit Boolean Evaluation *)
  (*$D+ Debug Code Generation ON (Sort of) *)
  (*$L+ Requires /V option to BPC to activate Local Debug symbols ON (Sort of) *)
  (*$F- Requires /V option to BPC to activate Far calls only as needed *)
  (*$I- I/O Checking OFF *)
  (*$N- Software Emulation of 80x87 *)
  (*$O- Overlays NOT allowed *)
  (*$P- Standard 'string' parameters *)
  (*$Q- Overflow Checking OFF *)
  (*$R- Range Checking OFF *)
  (*$S- Stack Checking OFF *)
  (*$V- Var-string Checking OFF *)
  (*$T+ Force Typed 'g' references *)
  (*$X+ Enable Extended syntax *)
}

interface

var {.....}
(* Stroked font 'handles' - Bold is outlined, block type. Simplex is *)
(* a very simple font (similar to the default graphics font) and *)
(* triplex is a more elegant font. *)
BOLD_FONT, SIMPLEX_FONT, TRIPLEX_FONT, SMALL_FONT : integer;

const {.....}
(* This variable is used to force the VESA driver into a mode which *)
(* is not its highest capability. *)
VESA_16_MODE : -1..2 = -1;

function DETECT_VESA_16 : integer;
{.....}
(* Used to detect the presence of a VESA-16 graphics capable controller. *)
(* Called as: InstallUserDriver('VESAL6', @DETECT_VESA_16); *)
{.....}

procedure VESAL6_DRIVER;
{.....}
(* The actual VESA-16 video driver. *)
(* Called as: registerbdriver(addr(VESAL6_DRIVER)); *)
{.....}

procedure EGAVGA_DRIVER;
{.....}
(* The actual VGA video driver. *)
(* Called as: registerbdriver(addr(EGAVGA_DRIVER)); *)
{.....}

implementation

uses graph;

function VESA_CAPABILITY
{.....}
(TABLE : pointer; MODES : word; SIZE : integer) : integer; near; assembler;
{.....}
(* Determines the highest graphics capability of the VESA-16 controller. *)
{.....}

asm (* VESA_CAPABILITY *)
xor ax,ax
les di, TABLE
@E1: mov si, MODES
add si, SIZE
add si, SIZE
mov bx, es: [di]
cmp bx, 0FFFFh
je @E4
inc di
inc di
mov cx, SIZE
@E2: cmp bx, [si]
jz @E3
dec si
dec si
@E3: loop @E2
@E4: cmp VESA_16_MODE, -1
je @E5
mov si, VESA_16_MODE
mov ah, 0
@E5: end; (* VESA_CAPABILITY *)

function DETECT_VESA_16 : integer; assembler;
{.....}
(* Determines if the video controller is VESA-16 compatible. *)
{.....}

const {.....}
(* The video modes supported by VESA-16. *)
VESA_16_MODES : array[0..2] of word = ($0102, $0104, $0106);

type {.....}
(* Data structure that defines the video controller. *)
VGA_INFO_BLOCK = record
  VESA_SIGNATURE : array[0..3] of byte;
  VESA_VERSION : word;
  OEM_STRING_PTR : pointer;
  CAPABILITIES : array[0..3] of byte;
  VIDEO_MODE_PTR : pointer;
end; (* VGA_INFO_BLOCK *)

var {.....}
(* Data structure used to determine VESA-16 capabilities. *)
VESA_INFO : array[0..255] of byte;

asm (* DETECT_VESA_16 *)
mov ax, 53
mov es, ax
lea di, VESA_INFO

```

## 10.6 BMP\_UTIL.PAS

```

mov     ax,4F00h
int     10h
cmp     ax,004Fh
mov     ax,gzError
jnz     @@Exit
cmp     es:(di).VGA_INFO_BLOCK.VESA_SIGNATURE.word(0),'EV'
jnz     @@Exit
cmp     es:(di).VGA_INFO_BLOCK.VESA_SIGNATURE.word(2),'AS'
jnz     @@Exit
les     di,es:(di).VGA_INFO_BLOCK.VIDEO_MODE_PTR
push   es
push   di
mov     ax,offset VESA_16_MODES
push   ax
mov     ax,3
push   ax
call    VESA_CAPABILITY
@@Exit:
end; (* DETECT_VESA_16 *)

(* Borland supplied drivers and fonts (derived from BGI files). *)
procedure VESA16_DRIVER; (* BINOBJ VESA16.BGI VESA16.OBJ VESA16_DRIVER *)
external: (*$L BINARYVESA16.OBJ *)
procedure EGAVGA_DRIVER; (* BINOBJ EGAVGA.BGI EGAVGA.OBJ EGAVGA_DRIVER *)
external: (*$L BINARYEGAVGA.OBJ *)
procedure TRIP_FONT_PROC; (* BINOBJ TRIP.CHR TRIP.OBJ TRIP_FONT_PROC *)
external: (*$L BINARYTRIP.OBJ *)
procedure SIMP_FONT_PROC; (* BINOBJ SIMP.CHR SIMP.OBJ SIMP_FONT_PROC *)
external: (*$L BINARYSIMP.OBJ *)
procedure BOLD_FONT_PROC; (* BINOBJ BOLD.CHR BOLD.OBJ BOLD_FONT_PROC *)
external: (*$L BINARYBOLD.OBJ *)
procedure LITT_FONT_PROC; (* BINOBJ SANS.CHR SANS.OBJ LITT_FONT_PROC *)
external: (*$L BINARYLITT.OBJ *)

begin (* GRAPHICS *)
(* Register Fonts *)
TRIPLEX_FONT := registerbfont(@TRIP_FONT_PROC);
SIMPLEX_FONT := registerbfont(@SIMP_FONT_PROC);
BOLD_FONT := registerbfont(@BOLD_FONT_PROC);
SMALL_FONT := registerbfont(@LITT_FONT_PROC);
end. (* GRAPHICS *)

```

```

unit BMP_UTIL;
(* This unit allows portions of the graphic screen to be saved to and *)
(* written from Windows v3.1 compatible 16 color, Bit-Map (BMP) files. *)
(* Compiler Options (Ver. 7.0) *)
(*$A+ Word Alignment *)
(*$B- Short Circuit Boolean Evaluation *)
(*$D+ Debug Code Generation ON (Sort of) *)
(* Requires /V option to BPC to activate *)
(*$L+ Local Debug symbols ON (Sort of) *)
(* Requires /V option to BPC to activate *)
(*$F- Far calls only as needed *)
(*$I- I/O Checking OFF *)
(*$N- Software Emulation of 80x87 *)
(*$O- Overlays NOT allowed *)
(*$P- Standard 'string' parameters *)
(*$Q- Overflow Checking OFF *)
(*$R- Range Checking OFF *)
(*$S- Stack Checking OFF *)
(*$V- Far-string Checking OFF *)
(*$T+ Force Typed 'g' references *)
(*$X+ Enable Extended syntax *)

```

interface

uses graph;

```

procedure DISPLAY_16_COLOR_BMP_FILE_AS_IMAGE
(FILENAME : string; X, Y : integer);
(* This routine reads the BMP file (must be 16 color) and displays it on *)
(* the screen (viewport relative). *)
procedure SAVE_IMAGE_AS_16_COLOR_BMP_FILE
(FILENAME : string; UL_X, UL_Y, LR_X, LR_Y : integer);
(* This routine takes the pixels on the screen and creates a BMP file (16 *)
(* color) from them (viewport relative). *)

```

implementation

```

type
(* Windows Bit Map File (BMP) header. *)
BIT_MAP_FILE_HEADER = record
(* File type must be 0x4D42 ('BM' for Bit Map). *)
FILE_TYPE : word;
(* Definition varies, ignored when read, written 0. *)
SIZE : longint;
(* Pointer 'Hot Spot', ignored when read, written 0. *)
X_HOT_SPOT : word;
Y_HOT_SPOT : word;
(* Offset within file to beginning of bit map, in BYTES. *)
OFFSET_TO_BITS : longint;
end; (* BIT_MAP_FILE_HEADER *)

(* Windows Bit Map (BMP) header. *)
BIT_MAP_HEADER = record
(* Size of this structure. *)
SIZE : longint;
(* Width of image (in pixels). *)
WIDTH : longint;
(* Height of image (in pixels), origin is lower left corner, if the *)
(* image origin is upper left corner then this value is negative. *)
HEIGHT : longint;
(* Must be 1 for windows bit maps. *)
NUMBER_OF_BIT_PLANES : integer;
(* Indicates color depth, allowable values: 1, 4 or 8. *)
NUMBER_OF_BITS_PER_PLANE : integer;
(* Zero for no compression. *)
COMPRESSION_SCHEME : longint;
(* Number of bytes the image data consumes, normally 0 if there is *)
(* no compression. The size is then computed from width, height and *)
(* color depth. *)
SIZE_OF_IMAGE_DATA : longint;
(* If non-zero this is used to scale the image. *)
X_RESOLUTION : longint;
(* If non-zero this is used to scale the image. *)
Y_RESOLUTION : longint;
(* Defines the number of valid entries in the color table, 0 means *)
(* 'all of them'. *)
NUMBER_OF_COLORS_USED : longint;
(* Defines the number of necessary entries in the color table, 0 *)
(* means 'all of them'. *)
NUMBER_OF_IMPORTANT_COLORE : longint;
end; (* BIT_MAP_HEADER *)

(* Windows color palette structure, Red, Green, Blue and padding to *)
(* make it 32-bits long. *)
RGB = record
BLUE : byte;

```

```

GREEN : byte;
RED   : byte;
padding : byte;
end; (* RGB *)

(*****
(* The BMP color palette is stored in the BMP file, since only 16 *)
(* color images are supported this structure can be a fixed length. *)
*****)
COLOR_TABLES = array(0..15) of RGB;

(*****
(* Images are processed on a per line basis, this allows images up *)
(* to 1280 pixels wide to be supported (two pixels per byte). *)
*****)
SCAN_LINES = array(0..639) of byte;

(*****
(* Borland and Microsoft have different color orderings, the color *)
(* translation is done using this structure to create look-up tables. *)
*****)
COLORS = record
  BMP : RGB;
  VGA : word;
end; (* COLORS *)

const (*****
(* This is the Borland/microsoft color translation look-up table. *)
*****)
TRANSLATE_COLOR_To : array(0..15) of COLORS =
(
  (BMP : (BLUE : 0; GREEN : 0; RED : 0; padding : 0); VGA : ($0)),
  (BMP : (BLUE : 192; GREEN : 0; RED : 0; padding : 0); VGA : ($4)),
  (BMP : (BLUE : 0; GREEN : 192; RED : 0; padding : 0); VGA : ($2)),
  (BMP : (BLUE : 192; GREEN : 192; RED : 0; padding : 0); VGA : ($6)),
  (BMP : (BLUE : 0; GREEN : 0; RED : 192; padding : 0); VGA : ($1)),
  (BMP : (BLUE : 192; GREEN : 0; RED : 192; padding : 0); VGA : ($5)),
  (BMP : (BLUE : 0; GREEN : 192; RED : 192; padding : 0); VGA : ($3)),
  (BMP : (BLUE : 192; GREEN : 192; RED : 192; padding : 0); VGA : ($7)),
  (BMP : (BLUE : 128; GREEN : 128; RED : 128; padding : 0); VGA : ($8)),
  (BMP : (BLUE : 255; GREEN : 0; RED : 0; padding : 0); VGA : ($C)),
  (BMP : (BLUE : 0; GREEN : 255; RED : 0; padding : 0); VGA : ($A)),
  (BMP : (BLUE : 255; GREEN : 255; RED : 0; padding : 0); VGA : ($E)),
  (BMP : (BLUE : 0; GREEN : 0; RED : 255; padding : 0); VGA : ($9)),
  (BMP : (BLUE : 255; GREEN : 0; RED : 255; padding : 0); VGA : ($D)),
  (BMP : (BLUE : 0; GREEN : 255; RED : 255; padding : 0); VGA : ($B)),
  (BMP : (BLUE : 255; GREEN : 255; RED : 255; padding : 0); VGA : ($F))
);

procedure DISPLAY_16_COLOR_BMP_FILE_AS_IMAGE
(*****
(FILENAME : string; X, Y : integer);
*****
(* This routine reads the BMP file (must be 16 color) and displays it on *)
(* the screen (viewport relative). *)
*****)
var (*****
(* The file being displayed. *)
*****)
BMP : file;

(*****
(* BMP file data structures. *)
*****)
FILE_HEADER : BIT_MAP_FILE_HEADER;
BMP_HEADER : BIT_MAP_HEADER;
COLOR_DEPTH : integer;
COLOR_TABLE : COLOR_TABLES;

(*****
(* Loop control variables to process lines and pixels. *)
*****)
LCV, PIXEL_LCV : integer;

(*****
(* Pixels are processed in lines. *)
*****)
SCAN_LINE : SCAN_LINES;
SCAN_LINE_SIZE : integer;
PIXEL : word;

begin (* DISPLAY_16_COLOR_BMP_FILE_AS_IMAGE *)
  assign(BMP, FILENAME);
  (*SI-*) reset(BMP, 1); (*SI+*)

  (* Get headers *)
  blockread(BMP, FILE_HEADER, sizeof(BIT_MAP_FILE_HEADER));
  blockread(BMP, BMP_HEADER, sizeof(BIT_MAP_HEADER));

  (* Get color table *)
  COLOR_DEPTH := 1;
  with BMP_HEADER do
    for LCV := 1 to NUMBER_OF_BIT_PLANES * NUMBER_OF_BITS_PER_PLANE do
      COLOR_DEPTH := COLOR_DEPTH * 2;
    if COLOR_DEPTH = 16 then
      begin
        for LCV := 0 to (COLOR_DEPTH - 1) do
          blockread(BMP, COLOR_TABLE[LCV], sizeof(RGB));

          (* Get pixel data *)
          SCAN_LINE_SIZE := ((BMP_HEADER.WIDTH + 3) div 4) * 2;
          seek(BMP, FILE_HEADER.OFFSET_TO_BITS);
          for LCV := 0 to BMP_HEADER.HEIGHT-1 do
            begin
              blockread(BMP, SCAN_LINE, SCAN_LINE_SIZE);

              (* Display this line of pixels *)
              for PIXEL_LCV := 0 to SCAN_LINE_SIZE-1 do
                begin
                  PIXEL := SCAN_LINE[PIXEL_LCV] shr 4;
                  putpixel
                    (
                      X + 2 * PIXEL_LCV,
                      Y + BMP_HEADER.HEIGHT - LCV - 1,
                      PIXEL
                    );
                  PIXEL := SCAN_LINE[PIXEL_LCV] and $F;
                  putpixel
                    (
                      X + 2 * PIXEL_LCV + 1,
                      Y + BMP_HEADER.HEIGHT - LCV - 1,
                      PIXEL
                    );
                end; (* for...to...do *)
              end; (* if...then *)
            close(BMP);
          end; (* DISPLAY_16_COLOR_BMP_FILE_AS_IMAGE *)
        end;
      end;
    end;
  end;

procedure SAVE_IMAGE_AS_16_COLOR_BMP_FILE
(*****
(FILENAME : string; UL_X, UL_Y, LR_X, LR_Y : integer);
*****
(* This routine takes the pixels on the screen and creates a BMP file (16 *)
(* color) from them (viewport relative). *)
*****)
var (*****
(* The file being displayed. *)
*****)
BMP : file;

(*****
(* BMP file data structures. *)
*****)
FILE_HEADER : BIT_MAP_FILE_HEADER;
BMP_HEADER : BIT_MAP_HEADER;
COLOR_DEPTH : integer;
COLOR_TABLE : COLOR_TABLES;

(*****
(* Loop control variables to process lines and pixels. *)
*****)
LCV, PIXEL_LCV : integer;

(*****
(* Pixels are processed in lines. *)
*****)
SCAN_LINE : SCAN_LINES;
SCAN_LINE_SIZE : integer;
PIXEL_1, PIXEL_2 : word;

begin (* SAVE_IMAGE_AS_16_COLOR_BMP_FILE *)
  assign(BMP, FILENAME);
  (*SI-*) rewrite(BMP, 1); (*SI+*)

  (* Put headers *)
  with FILE_HEADER do
    begin
      FILE_TYPE := $4D42;
      SIZE := sizeof(BIT_MAP_FILE_HEADER);
      X_HOT_SPOT := 0;
      Y_HOT_SPOT := 0;
      OFFSET_TO_BITS := sizeof(BIT_MAP_FILE_HEADER) +
        sizeof(BIT_MAP_HEADER) + 16 * sizeof(RGB);
    end; (* with...do *)
  blockwrite(BMP, FILE_HEADER, sizeof(BIT_MAP_FILE_HEADER));

  with BMP_HEADER do
    begin
      SIZE := sizeof(BIT_MAP_HEADER);
      WIDTH := LR_X - UL_X + 1;
      HEIGHT := LR_Y - UL_Y + 1;
      NUMBER_OF_BIT_PLANES := 1;
      NUMBER_OF_BITS_PER_PLANE := 4;
      COMPRESSION_SCHEME := 0;
      SIZE_OF_IMAGE_DATA := 0;
      X_RESOLUTION := 0;
      Y_RESOLUTION := 0;
      NUMBER_OF_COLORS_USED := 16;
      NUMBER_OF_IMPORTANT_COLORS := 16;
    end; (* with...do *)
  blockwrite(BMP, BMP_HEADER, sizeof(BIT_MAP_HEADER));

  (* Put color table *)
  COLOR_DEPTH := 16;
  for LCV := 0 to COLOR_DEPTH-1 do
    blockwrite(BMP, TRANSLATE_COLOR_To[LCV].BMP, sizeof(RGB));

    (* Put pixel data *)
    SCAN_LINE_SIZE := ((BMP_HEADER.WIDTH + 3) div 4) * 2;
    seek(BMP, FILE_HEADER.OFFSET_TO_BITS);
    for LCV := 0 to BMP_HEADER.HEIGHT do
      begin
        (* Get pixels from screen into this line *)
        for PIXEL_LCV := 0 to SCAN_LINE_SIZE-1 do
          begin
            PIXEL_1 := getpixel(2 * PIXEL_LCV, BMP_HEADER.HEIGHT - LCV - 1);
            PIXEL_2 := getpixel(2 * PIXEL_LCV + 1, BMP_HEADER.HEIGHT - LCV - 1);
            SCAN_LINE[PIXEL_LCV] := (PIXEL_1 shl 4) or (PIXEL_2 and $F);
          end; (* for...to...do *)
        blockwrite(BMP, SCAN_LINE, SCAN_LINE_SIZE);
      end; (* for...to...do *)
    close(BMP);
  end; (* SAVE_IMAGE_AS_16_COLOR_BMP_FILE *)

  (* No initialization *)
end. (* BIT_MAP_UTILITIES *)

```

# 10.7 MS\_MOUSE.PAS

```

unit MS_MOUSE;
(* This unit provides access to a Microsoft Mouse Driver for text and
  graphics mode Turbo Pascal application. *)
(*-----*)
(* Compiler Options (Ver. 7.0) *)
(*$A+ Word Alignment *)
(*$B- Short Circuit Boolean Evaluation *)
(*$D- Debug Code Generation ON (Sort of) *)
(*$E- Requires /V option to TPC to activate *)
(*$L+ Local Debug symbols ON (Sort of) *)
(*$M- Requires /V option to TPC to activate *)
(*$P- Far calls only as needed *)
(*$I- I/O Checking OFF *)
(*$N- Software Emulation of 80x87 *)
(*$O- Overlays NOT allowed *)
(*$R- Range Checking OFF *)
(*$S- Stack Checking OFF *)
(*$V- Var-string Checking OFF *)
(*-----*)
(* NOTE: This code is very sensitive to compiler options, i.e. the event
  handler won't work if checks which change code at the entry point
  of a routine are added ($S, $V). *)
(*-----*)

interface

const (
  (* This is the application visible flag which indicates whether an
    MS-Mouse driver is installed on the PC or not *)
  MOUSE_INSTALLED : boolean = false;
);

type (
  (* This type is used to allow communication between the application
    mouse handler and the mouse driver. *)
  MOUSE_PARAMETERS = record
    (* Indicates why the application was activated. See allowable
    * events defined below. Note that it is possible for more than one
    * event to be active at the same time. *)
    ACTIVITY_MASK : word;
    (* Contains status of both mouse buttons at the time of activation. *)
    BUTTON_STATES : word;
    (* Absolute screen position of the mouse cursor, relative to upper
    * left corner which is at (1, 1). *)
    VERTICAL_TEXT_POSITION, HORIZONTAL_TEXT_POSITION : word;
    (* Absolute screen position of the mouse cursor, relative to upper
    * left corner which is at (0, 0). *)
    VERTICAL_MICKEY_POSITION, HORIZONTAL_MICKEY_POSITION : word;
  end; (* MOUSE_PARAMETERS *)
);

(* The procedure used to handle the mouse events must be declared of
  this type. Note that the actual procedure declared must be 'far'. *)
HANDLER_PROC = procedure(PARAMS : MOUSE_PARAMETERS);

(*-----*)
(* Maximum values for mouse mickey positions. *)
(*-----*)
var MAX_HORIZONTAL_MICKEY_POSITION, MAX_VERTICAL_MICKEY_POSITION : integer;

procedure GET_MOUSE_CURSOR_POSITION(var HORIZONTAL, VERTICAL : word);
procedure SET_MOUSE_CURSOR_POSITION(HORIZONTAL, VERTICAL : word);
(* These procedures manipulate the MS Mouse cursor on the text mode video
  screen. MS Mouse cursor positions are independent of the keyboard cursor
  position but are otherwise positioned similarly. The maximum vertical
  position is a function of the video mode selected. *)

(*-----*)
(* Horizontal cursor range diagram *)
(*-----*)
(* This routine restricts the mouse to movement within a column of the text
  screen. When combined with row restrictions this can be used to
  arbitrarily restrict mouse movement. *)
(*-----*)
(* Vertical cursor range diagram *)
(*-----*)
(* This routine restricts the mouse to movement within a row of the text
  screen. When combined with column restrictions this can be used to
  arbitrarily restrict mouse movement. *)
(*-----*)

```

```

(* This routine restricts the mouse to movement within a row of the text
  screen. When combined with column restrictions this can be used to
  arbitrarily restrict mouse movement. *)
(*-----*)
(* Horizontal cursor range diagram *)
(*-----*)
(* This routine restricts the mouse to movement within a row of the text
  screen. When combined with column restrictions this can be used to
  arbitrarily restrict mouse movement. *)
(*-----*)
(* Vertical cursor range diagram *)
(*-----*)
(* This routine restricts the mouse to movement within a row of the text
  screen. When combined with column restrictions this can be used to
  arbitrarily restrict mouse movement. *)
(*-----*)

procedure SHOW_MOUSE_CURSOR;
(* If the cursor is off then this routine turns the mouse cursor on. If
  the cursor is on then it becomes 'more on', i.e. the number of times that
  it was turned on is retained such that the same number of 'off' commands
  must be issued before the cursor will actually disappear. *)

procedure HIDE_MOUSE_CURSOR;
(* If the cursor is on then this routine turns the mouse cursor off. If
  the cursor is off then it becomes 'more off', i.e. the number of times
  that it was turned off is retained such that the same number of 'on'
  commands must be issued before the cursor will actually appear. *)

const (
  (* The following events can be handled by a Microsoft compatible
    mouse (and hence this unit). Note that the MS Mouse has only two
    buttons (left and right). *)
  CURSOR_MOVEMENT = $01;
  LEFT_BUTTON_PRESSED = $02;
  LEFT_BUTTON_RELEASED = $04;
  LEFT_BUTTON_ACTIVITY = LEFT_BUTTON_PRESSED or LEFT_BUTTON_RELEASED;
  RIGHT_BUTTON_PRESSED = $08;
  RIGHT_BUTTON_RELEASED = $10;
  RIGHT_BUTTON_ACTIVITY = RIGHT_BUTTON_PRESSED or RIGHT_BUTTON_RELEASED;
  BUTTON_PRESSED = LEFT_BUTTON_PRESSED or RIGHT_BUTTON_PRESSED;
  BUTTON_RELEASED = LEFT_BUTTON_RELEASED or RIGHT_BUTTON_RELEASED;
  BUTTON_ACTIVITY = LEFT_BUTTON_ACTIVITY or RIGHT_BUTTON_ACTIVITY;
  ANY_ACTIVITY = CURSOR_MOVEMENT or BUTTON_ACTIVITY;
);

procedure PUSH_MOUSE_EVENT_HANDLER(MASK : word; HANDLER : HANDLER_PROC);
(* This routine establishes the mouse activities which will cause the user
  application handler to be called. The user application should consider
  the mouse handler call to be an interrupt in so far as not returning from
  the handler call will cause system interrupts to be disabled and may
  result in erratic system behavior. The best approach to mouse activity
  handling is for the mouse handler to manipulate global data structures
  and for the application to periodically check the structure in order to
  take appropriate action. *)
(* The activity mask bits are defined as follows (see above definitions): *)
(*-----*)
(* Activity mask bit diagram *)
(*-----*)
(* NOTE: Using the handler 'push' concept for more than a single installation
  requires that the application provide adequate heap-space (each
  mouse handler push requires -10 bytes of heap). *)

procedure POP_MOUSE_EVENT_HANDLER;
(* This routine restores a previously installed handler and activity mask.
  the present handler is removed. This is useful for layered applications
  where mouse actions are different based on entry/exit of procedures or
  displaying of windows (i.e. handler must change as windows are brought
  to foreground). *)

function LEFT_BUTTON_HELD : boolean;
(* This routine merely returns the status of the left mouse button, true if
  the button is pressed and false otherwise. *)

function RIGHT_BUTTON_HELD : boolean;
(* This routine merely returns the status of the right mouse button, true
  if the button is pressed and false otherwise. *)

procedure DISABLE_MOUSE_DRIVER;
(* This routine shuts down the mouse driver so that events will not be
  reported and status will not be available. *)

procedure ENABLE_MOUSE_DRIVER;
(* This routine turns on the mouse driver so that events will be reported
  and status will be available. *)

procedure RESET_MOUSE_DRIVER;
(* This routine causes all mouse driver parameters to be reset to default
  values. This includes activity mask and mouse cursor position. *)

const (
  (* The following constants are used to control the characteristics of
    the mouse cursor. Only one from each group can be specified at any
    given time using SET/CLEAR_CURSOR_SCREEN functions. *)
  (* Controls mouse cursor blinking. *)
  BLINK_MOUSE = $80;
  (* Controls mouse cursor background color (in a manner similar to
    standard text/background colors). *)

```

```

(*****)
BACKGROUND_BLACK = $00;
BACKGROUND_BLUE = $10;
BACKGROUND_GREEN = $20;
BACKGROUND_CYAN = $30;
BACKGROUND_RED = $40;
BACKGROUND_MAGENTA = $50;
BACKGROUND_BROWN = $60;
BACKGROUND_LIGHT_GRAY = $70;

(*****)
(* Controls mouse cursor foreground color (in a manner similar to
 * standard text/background colors).
(*****)
FOREGROUND_BLACK = $00;
FOREGROUND_BLUE = $01;
FOREGROUND_GREEN = $02;
FOREGROUND_CYAN = $03;
FOREGROUND_RED = $04;
FOREGROUND_MAGENTA = $05;
FOREGROUND_BROWN = $06;
FOREGROUND_LIGHT_GRAY = $07;
FOREGROUND_DARK_GRAY = $08;
FOREGROUND_LIGHT_BLUE = $09;
FOREGROUND_LIGHT_GREEN = $0A;
FOREGROUND_LIGHT_CYAN = $0B;
FOREGROUND_LIGHT_RED = $0C;
FOREGROUND_LIGHT_MAGENTA = $0D;
FOREGROUND_YELLOW = $0E;
FOREGROUND_WHITE = $0F;

procedure SET_CURSOR_SCREEN_BITS(SCREEN BITS : byte);
(*****)
(* Causes the specified attribute(s) of the mouse cursor to become active,
 * all unspecified attributes remain the same.
(*****)
procedure CLEAR_CURSOR_SCREEN_BITS(SCREEN BITS : byte);
(*****)
(* Causes the specified attribute(s) of the mouse cursor to become dormant,
 * all unspecified attributes remain the same.
(*****)
procedure SET_CURSOR_CHARACTER(CURSOR : char);
(*****)
(* Changes the mouse cursor to the specified character, if the character
 * requested is the null character, i.e., chr(0), then the default mouse
 * cursor is used. The default mouse cursor is the character which is
 * beneath the cursor.
(*****)
type
(* Graphics cursor description.
(*****)
GRAPHIC_CURSOR = record
IMAGE : array [1..32] of word;
X_HOT, Y_HOT : word;
end; (* GRAPHIC_CURSOR *)

const
(* Cursor in the shape of a check mark.
(*****)
CHECK_MARK_CURSOR : GRAPHIC_CURSOR =
(
IMAGE :
(
$FFFD, $FFFD, $FFFD, $FFFD, ( screen mask )
$FF03, $0607, $000E, $001F,
$003B, $007E, $FFFE, $FFFE,
$FFFE, $FFFE, $FFFE, $FFFE,
$0000, $0006, $000C, $0018, ( cursor mask )
$0030, $0060, $070C, $1D80,
$0700, $0000, $0000, $0000,
$0000, $0000, $0000, $0000
);
X_HOT : 6;
Y_HOT : 7;
);

(* Cursor in the shape of an arrow pointing left.
(*****)
LEFT_ARROW_CURSOR : GRAPHIC_CURSOR =
(
IMAGE :
(
$FE1F, $FD1F, $0000, $0000, ( screen mask )
$0000, $001F, $E1FF, $E1FF,
$FFFE, $FFFE, $FFFE, $FFFE,
$FFFE, $FFFE, $FFFE, $FFFE,
$0000, $000C, $070C, $7FEF, ( cursor mask )
$070C, $000C, $0000, $0000,
$0000, $0000, $0000, $0000,
$0000, $0000, $0000, $0000
);
X_HOT : 0;
Y_HOT : 3;
);

(* Cursor in the shape of a cross.
(*****)
CROSS_CURSOR : GRAPHIC_CURSOR =
(
IMAGE :
(
$FC3F, $FC3F, $FC3F, $0000, ( screen mask )
$0000, $0000, $FC3F, $FC3F,
$FC3F, $FFFE, $FFFE, $FFFE,
$FFFE, $FFFE, $FFFE, $FFFE,
$0000, $0180, $0180, $0180, ( cursor mask )
$7FEF, $0180, $0180, $0180,
$0000, $0000, $0000, $0000,
$0000, $0000, $0000, $0000
);
X_HOT : 7;
Y_HOT : 4;
);

(* Cursor in the shape of a hand with a pointing finger.
(*****)
POINTING_HAND_CURSOR : GRAPHIC_CURSOR =
(
IMAGE :
(
$E1FF, $E1FF, $E1FF, $E1FF, ( screen mask )
$E1FF, $E000, $E000, $E000,
$0000, $0000, $0000, $0000,
$0000, $0000, $0000, $0000,
$1E00, $1200, $1200, $1200, ( cursor mask )
$1200, $13FF, $1249, $1249,
$1249, $9001, $9001, $9001,
$8001, $8001, $8001, $FFFE
);
X_HOT : 5;
Y_HOT : 0;
);

)
(* Cursor in the shape of an 'I'-Beam (for text editing).
(*****)
I_BEAM_CURSOR : GRAPHIC_CURSOR =
(
IMAGE :
(
$FFFF, $FFFF, $FFFF, $FFFF, ( screen mask )
$FFFF, $FFFF, $FFFF, $FFFF,
$FFFF, $FFFF, $FFFF, $FFFF,
$FFFF, $FFFF, $FFFF, $FFFF,
$F00E, $0C30, $0240, $0240, ( cursor mask )
$0180, $0180, $0180, $0180,
$0180, $0180, $0180, $0180,
$0240, $0240, $0C30, $F00E
);
X_HOT : 7;
Y_HOT : 7;
);

procedure SET_GRAPHIC_CURSOR(var CURSOR : GRAPHIC_CURSOR);
(*****)
(* Causes the specified graphic cursor to be used for display. The video
 * mode must be graphics "before" the mouse is reset for the graphic cursor
 * to be visible.
(*****)
implementation
uses dos;

const MOUSE_SERVICE_NUMBER = $33;

(* Mouse functions, always passed to the driver in AX *)
MOUSE_RESET_AND_STATUS = $00;
SHOW_CURSOR = $01;
HIDE_CURSOR = $02;
GET_BUTTON_STATUS_AND_POSITION = $03;
SET_CURSOR_POSITION = $04;
GET_BUTTON_PRESS_INFORMATION = $05;
GET_BUTTON_RELEASE_INFORMATION = $06;
SET_MIN_MAX_HORIZONTAL_POSITION = $07;
SET_MIN_MAX_VERTICAL_POSITION = $08;
SET_GRAPHICS_CURSOR_BLOCK = $09;
SET_TEXT_CURSOR = $0A;
READ_MOTION_COUNTERS = $0B;
SET_CALL_MASK_AND_ADDRESS = $0C;
SET_LIGHT_PEN_EMULATION_ON = $0D;
SET_LIGHT_PEN_EMULATION_OFF = $0E;
SET_MICKY_TO_PIXEL_RATIO = $0F;
CONDITIONAL_OFF = $10;
SET_DOUBLE_SPEED_THRESHOLD = $13;
SWAP_INTERRUPT_ROUTINES = $14;
GET_DRIVER_STATE_STORAGE_ROMTS = $15;
SAVE_DRIVER_STATE = $16;
RESTORE_DRIVER_STATE = $17;
SET_ALT_CALL_MASK_AND_ADDRESS = $18;
GET_USER_ALT_INTERRUPT_ADDRESS = $19;
SET_SENSITIVITY = $1A;
GET_SENSITIVITY = $1B;
SET_INTERRUPT_RATE = $1C;
SET_CRT_PAGE_NUMBER = $1D;
GET_CRT_PAGE_NUMBER = $1E;
DISABLE_DRIVER = $1F;
ENABLE_DRIVER = $20;
SOFTWARE_RESET = $21;
SET_LANGUAGE_FOR_MESSAGES = $22;
GET_LANGUAGE_NUMBER = $23;
GET_DRIVER_VER_TYPE_AND_IRQ_NUM = $24;
GET_GENERAL_DRIVER_INFORMATION = $25;
GET_MAXIMUM_VIRTUAL_CDCOORDINATES = $26;
GET_CURSOR_MASKS_AND_MICKY_CNTR = $27;
SET_VIDEO_MODE = $28;
EMULATE_VIDEO_MODES = $29;
GET_CURSOR_HOTSPOT = $30;
LOAD_ACCELERATION_CURVES = $31;
READ_ACCELERATION_CURVES = $32;
SET_GET_ACCELERATION_CURVES = $33;
MOUSE_HARDWARE_RESET = $35;
SET_GET_BALLPOINT_INFORMATION = $36;
GET_MIN_MAX_VIRTUAL_CDCOORDINATES = $37;
GET_ACTIVE_ADVANCED_FUNCTIONS = $38;
GET_SWITCH_SETTINGS = $39;
GET_MOUSE_INIT_LOCATION = $40;

(*****)
(* Address of exit procedure for previous member of exit chain.
(*****)
var MOUSE_SAVE_EXIT : pointer;

procedure GET_MOUSE_CURSOR_POSITION(var HORIZONTAL, VERTICAL : word);
(*****)
(* This routine merely asks the mouse driver what the current cursor
 * coordinates are and returns those values to the caller.
 * NOTE: An installed MS-Mouse driver is assumed.
(*****)
var REGS registers;

begin (* GET_MOUSE_CURSOR_POSITION *)
(
table
Parameter(s) Return Values
AX GET_BUTTON_STATUS_AND_POSITION Button Status
BX 0=Button up
1=Button held down
Bit 0 used for left button
Bit 1 used for right button
CX Horizontal position (mickeys)
DX Vertical position (mickeys)
end;
with REGS do
begin
AX := GET_BUTTON_STATUS_AND_POSITION;
CX := 0;
DX := 0;
intr(MOUSE_SERVICE_NUMBER, REGS);
HORIZONTAL := CX shr 3;
VERTICAL := DX shr 3;
end; (* with...do *)
end; (* GET_MOUSE_CURSOR_POSITION *)

procedure SET_MOUSE_CURSOR_POSITION(HORIZONTAL, VERTICAL : word);
(*****)
(* Mouse equivalent of gotoxy(x, y).
 * NOTE: An installed MS-Mouse driver is assumed.
(*****)
var REGS registers;

begin (* SET_MOUSE_CURSOR_POSITION *)
(
table
Parameter(s) Return Values
AX SET_CURSOR_POSITION
end;

```

```

(*
(* CX Horizontal position (mickeys)
(*
(* DX Vertical position (mickeys)
(*
with REGS do
begin
  AX := SET_CURSOR_POSITION;
  BX := HORIZONTAL shl 3;
  DX := VERTICAL shl 3;
end; (* with...do *)
intr(MOUSE_SERVICE_NUMBER, REGS);
end; (* SET_MOUSE_CURSOR_POSITION *)

procedure SET_HORIZONTAL_CURSOR_RANGE(MIN, MAX : word);
(* Sets the limit on the mouse movement in the horizontal direction. *)
(* NOTE: An installed MS-Mouse driver is assumed. *)
var REGS : registers;
begin (* SET_HORIZONTAL_CURSOR_RANGE *)
  (*
  (* Parameter(s) Return Values
  (*
  (* AX SET_MIN_MAX_HORIZONTAL_POSITION
  (* CX Minimum horizontal position (mickeys)
  (* DX Maximum horizontal position (mickeys)
  (*
  (* NOTE: if minimum is greater than maximum then the values are swapped *)
  with REGS do
  begin
    AX := SET_MIN_MAX_HORIZONTAL_POSITION;
    CX := MIN shl 3 - 1;
    DX := MAX shl 3 - 1;
  end; (* with...do *)
  intr(MOUSE_SERVICE_NUMBER, REGS);
  end; (* SET_HORIZONTAL_CURSOR_RANGE *)

procedure SET_VERTICAL_CURSOR_RANGE(MIN, MAX : word);
(* Sets the limit on the mouse movement in the vertical direction. *)
(* NOTE: An installed MS-Mouse driver is assumed. *)
var REGS : registers;
begin (* SET_VERTICAL_CURSOR_RANGE *)
  (*
  (* Parameter(s) Return Values
  (*
  (* AX SET_MIN_MAX_VERTICAL_POSITION
  (* CX Minimum vertical position (mickeys)
  (* DX Maximum vertical position (mickeys)
  (*
  (* NOTE: if minimum is greater than maximum then the values are swapped *)
  with REGS do
  begin
    AX := SET_MIN_MAX_VERTICAL_POSITION;
    CX := MIN shl 3 - 1;
    DX := MAX shl 3 - 1;
  end; (* with...do *)
  intr(MOUSE_SERVICE_NUMBER, REGS);
  end; (* SET_VERTICAL_CURSOR_RANGE *)

procedure SHOW_MOUSE_CURSOR;
(* Turns on the mouse cursor. *)
(* NOTE: An installed MS-Mouse driver is assumed. *)
var REGS : registers;
begin (* SHOW_MOUSE_CURSOR *)
  (*
  (* Parameter(s) Return Values
  (*
  (* AX SHOW_CURSOR
  (*
  REGS.AX := SHOW_CURSOR;
  intr(MOUSE_SERVICE_NUMBER, REGS);
  end; (* SHOW_MOUSE_CURSOR *)

procedure HIDE_MOUSE_CURSOR;
(* Turns off the mouse cursor. *)
(* NOTE: An installed MS-Mouse driver is assumed. *)
var REGS : registers;
begin (* HIDE_MOUSE_CURSOR *)
  (*
  (* Parameter(s) Return Values
  (*
  (* AX HIDE_CURSOR
  (*
  REGS.AX := HIDE_CURSOR;
  intr(MOUSE_SERVICE_NUMBER, REGS);
  end; (* HIDE_MOUSE_CURSOR *)

(* These declarations ensure that the user application won't crash if the
(* mouse is enabled before the application has installed a handler. And *)
(* allow handlers to be pushed/popped to accommodate layered applications. *)
procedure NULL_MOUSE_HANDLER(PARAMS : MOUSE_PARAMETERS); far;
begin (* No Op *) end;

type HANDLER_NODE_PTR = ^HANDLER_NODE;
HANDLER_NODE = record
  HANDLER : HANDLER_PROC;
  MASK : word;
  NEXT : HANDLER_NODE_PTR;
end; (* HANDLER_NODE *)

const DEFAULT_HANDLER : HANDLER_NODE =
  (
    HANDLER : NULL_MOUSE_HANDLER;
    MASK : 0000;
    NEXT : nil
  ); (* DEFAULT_HANDLER *)

TOP_OF_HANDLER_STACK : HANDLER_NODE_PTR := @DEFAULT_HANDLER;

procedure GENERIC_MOUSE_EVENT_HANDLER; far;
(* This routine is installed to interface the Mouse Driver call to the
(* Pascal application. The Pascal handler must be defined to be of type

```

```

(* HANDLER_PROC. *)
(* NOTE: An installed MS-Mouse driver is assumed. *)
var PARAMS : MOUSE_PARAMETERS;
MAX_X, MAX_Y : integer;

begin (* GENERIC_MOUSE_EVENT_HANDLER *)
asm
  (* NOTE: Subroutine is passed information as follows: *)
  (*
  (* AX Mask with condition bit(s) set that triggered call
  (*
  (* BX Button State
  (* 0=Button up
  (* 1=Button held down
  (* Bit 0 used for left button
  (* Bit 1 used for right button
  (*
  (* CX Horizontal cursor position (mickeys)
  (*
  (* DX Vertical cursor position (mickeys)
  (*
  (* SI Horizontal mouse counts since last reset (mickeys)
  (*
  (* DI Vertical mouse counts since last reset (mickeys)
  (*
  (* Move mouse parameters from registers into variable *)
  mov PARAMS.ACTIVITY_MASK, ax
  mov PARAMS.BUTTON_STATES, bx
  mov PARAMS.HORIZONTAL_MICKEY_POSITION, cx
  mov PARAMS.VERTICAL_MICKEY_POSITION, dx
  (* Convert CX from mickeys to character cursor coordinates *)
  mov bx, cx
  shr bx, cl
  inc bx
  mov PARAMS.HORIZONTAL_TEXT_POSITION, bx
  (* Convert DX from mickeys to character cursor coordinates *)
  shr dx, cl
  inc dx
  mov PARAMS.VERTICAL_TEXT_POSITION, dx
  (* Save Mouse Driver's registers *)
  push es
  push ds
  (* Setup application data segment *)
  mov ax, seg @DATA
  mov ds, ax
end; (* asm *)

(* Let the user application see the parameters *)
TOP_OF_HANDLER_STACK^.HANDLER(PARAMS);

(* Restore Mouse Driver's registers *)
asm
  pop ds
  pop es
end; (* asm *)
end; (* GENERIC_MOUSE_EVENT_HANDLER *)

procedure PUSH_MOUSE_EVENT_HANDLER(MASK : word; HANDLER : HANDLER_PROC);
(* Must be called to install the application mouse handler procedure *)
(* before the application can receive events. *)
(* NOTE: An installed MS-Mouse driver is assumed. *)
var TEMP_PTR : HANDLER_NODE_PTR;
REGS : registers;
begin (* PUSH_MOUSE_EVENT_HANDLER *)
  if @TOP_OF_HANDLER_STACK^.HANDLER <> @NULL_MOUSE_HANDLER then
  begin
    new(TEMP_PTR);
    TEMP_PTR^.NEXT := TOP_OF_HANDLER_STACK;
    TOP_OF_HANDLER_STACK := TEMP_PTR;
    end; (* If...then *)
    TOP_OF_HANDLER_STACK^.HANDLER := HANDLER;
    TOP_OF_HANDLER_STACK^.MASK := MASK;
  (*
  (* Parameter(s) Return Values
  (*
  (* AX SET_CALL_MASK_AND_ADDRESS
  (*
  (* CX Call mask
  (* Bit 0=Cursor position changes
  (* Bit 1=Left button pressed
  (* Bit 2=Right button pressed
  (* Bit 3=Left button released
  (* Bit 4=Right button released
  (* Bit 5-15=Unused
  (*
  (* DX Offset of subroutine
  (*
  (* ES Segment of subroutine
  (*
  with REGS do
  begin
    AX := SET_CALL_MASK_AND_ADDRESS;
    CX := MASK;
    DX := ofs(GENERIC_MOUSE_EVENT_HANDLER);
    ES := seg(GENERIC_MOUSE_EVENT_HANDLER);
  end; (* with...do *)
  intr(MOUSE_SERVICE_NUMBER, REGS);
  end; (* PUSH_MOUSE_EVENT_HANDLER *)

procedure POP_MOUSE_EVENT_HANDLER;
(* This routine restores a previously installed handler and activity *)
(* mask. *)
var TEMP_PTR : HANDLER_NODE_PTR;
MASK : word;
REGS : registers;
begin (* POP_MOUSE_EVENT_HANDLER *)
  TEMP_PTR := TOP_OF_HANDLER_STACK^.NEXT;
  if TEMP_PTR = nil
  then
  begin
    TOP_OF_HANDLER_STACK^.HANDLER := NULL_MOUSE_HANDLER;
    TOP_OF_HANDLER_STACK^.MASK := 0000;
    end; (* If...then *)
  else
  begin
    dispose(TOP_OF_HANDLER_STACK);
    TOP_OF_HANDLER_STACK := TEMP_PTR;
    end; (* If...then...else *)
  (* Install new handler *)
  MASK := TOP_OF_HANDLER_STACK^.MASK;
  (*
  (* Parameter(s) Return Values
  (*
  (* AX SET_CALL_MASK_AND_ADDRESS

```

```

[*] CX Call mask
[*] Bit 0=Cursor position changes
[*] Bit 1=Left button pressed
[*] Bit 2=Right button pressed
[*] Bit 3=Left button released
[*] Bit 4=Right button released
[*] Bit 5-15=Unused
[*] DX Offset of subroutine
[*] ES Segment of subroutine
with REGS do
begin
AX := SET_CALL_MASK_AND_ADDRESS;
CX := MASK;
DX := ofs|GENERIC_MOUSE_EVENT_HANDLER|;
ES := seg|GENERIC_MOUSE_EVENT_HANDLER|;
end; [*] with...do [*]
intr|MOUSE_SERVICE_NUMBER, REGS|;
end; [*] POP_MOUSE_EVENT_HANDLER [*]

function LEFT_BUTTON_HELD : boolean;
.....
[*] Determines whether the Left button is being held down or not.
[*] NOTE: An installed MS-Mouse driver is assumed.
.....

var REGS : registers;

begin [*] LEFT_BUTTON_HELD [*]
.....
Parameter|s| Return Values
AX GET_BUTTON_STATUS_AND_POSITION
BX Button Status
0=Button up
1=Button held down
Bit 0 used for left button
Bit 1 used for right button
CX Horizontal position [mickeys]
DX Vertical position [mickeys]
with REGS do
begin
AX := GET_BUTTON_STATUS_AND_POSITION;
BX := 0;
intr|MOUSE_SERVICE_NUMBER, REGS|;
LEFT_BUTTON_HELD := |BX and $01| = $01;
end; [*] with...do [*]
end; [*] LEFT_BUTTON_HELD [*]

function RIGHT_BUTTON_HELD : boolean;
.....
[*] Determines whether the Right button is being held down or not.
[*] NOTE: An installed MS-Mouse driver is assumed.
.....

var REGS : registers;

begin [*] RIGHT_BUTTON_HELD [*]
.....
Parameter|s| Return Values
AX GET_BUTTON_STATUS_AND_POSITION
BX Button Status
0=Button up
1=Button held down
Bit 0 used for left button
Bit 1 used for right button
CX Horizontal position [mickeys]
DX Vertical position [mickeys]
with REGS do
begin
AX := GET_BUTTON_STATUS_AND_POSITION;
BX := 0;
intr|MOUSE_SERVICE_NUMBER, REGS|;
RIGHT_BUTTON_HELD := |BX and $02| = $02;
end; [*] with...do [*]
end; [*] RIGHT_BUTTON_HELD [*]

procedure DISABLE_MOUSE_DRIVER;
.....
[*] Simply disables mouse events by turning off the driver.
[*] NOTE: An installed MS-Mouse driver is assumed.
.....

var REGS : registers;

begin [*] DISABLE_MOUSE_DRIVER [*]
.....
Parameter|s| Return Values
AX DISABLE_DRIVER Status
-1=Error occurred
BX Offset of old interrupt handler
[MOUSE_SERVICE_NUMBER]
ES Segment of old interrupt
handler [MOUSE_SERVICE_NUMBER]
REGS.AX := DISABLE_DRIVER;
intr|MOUSE_SERVICE_NUMBER, REGS|;
end; [*] DISABLE_MOUSE_DRIVER [*]

procedure ENABLE_MOUSE_DRIVER;
.....
[*] Simply enables mouse events by turning on the driver.
[*] NOTE: An installed MS-Mouse driver is assumed.
.....

var REGS : registers;

begin [*] ENABLE_MOUSE_DRIVER [*]
.....
Parameter|s| Return Values
AX ENABLE_DRIVER
REGS.AX := ENABLE_DRIVER;
intr|MOUSE_SERVICE_NUMBER, REGS|;
end; [*] ENABLE_MOUSE_DRIVER [*]

procedure RESET_MOUSE_DRIVER;
.....
[*] Simply resets all mouse driver variables to initialized state.
[*] NOTE: This will also determine if a mouse driver is installed or not.
.....

var REGS : registers;

```

```
begin [*] RESET_MOUSE_DRIVER [*]
```

Parameter s	Return Values
AX SET_SENSITIVITY	
BX Horizontal mickey sensitivity	
CX Vertical mickey sensitivity	
DX Threshold for double speed	

```

with REGS do
begin
AX := SET_SENSITIVITY;
BX := 200;
CX := 200;
DX := 1;
intr|MOUSE_SERVICE_NUMBER, REGS|;
end; [*] with...do [*]

```

Parameter s	Return Values
AX GET_MAXIMUM_VIRTUAL_COORDINATES	
BX	Mouse disabled flag 0=Enabled -1=Disabled
CX	Absolute virtual X maximum
DX	absolute virtual Y maximum

```

with REGS do
begin
AX := GET_MAXIMUM_VIRTUAL_COORDINATES;
intr|MOUSE_SERVICE_NUMBER, REGS|;
MAX_HORIZONTAL_MICKY_POSITION := CX;
MAX_VERTICAL_MICKY_POSITION := DX;
end; [*] with...do [*]

```

Parameter s	Return Values
AX SOFTWARE_RESET	Status -1=Mouse driver installed, SOFTWARE_RESET otherwise
BX	2 [if Status=-1]

```

with REGS do
begin
AX := SOFTWARE_RESET;
intr|MOUSE_SERVICE_NUMBER, REGS|;
MOUSE_INSTALLED := |AX = $FFFF| and |BX = 2|;
end; [*] with...do [*]
end; [*] RESET_MOUSE_DRIVER [*]

```

```

procedure GET_SCREEN_AND_CURSOR_MASKS|var SCREEN_MASK, CURSOR_MASK : word|;
.....
[*] This routine in conjunction with SET_SCREEN_AND_CURSOR_MASKS are
[*] used to modify the mouse cursor given in the following tables.
.....

```

Masks			Screen Bits		
Screen Mask	Cursor Mask	Screen Bit	Bit Number	Description	Comments
0	0	0	15	Blink Ctrl	1=Blink
0	1	1	14-12	Background	Color
1	0	Unchanged	11	Intensity	1=High
1	1	Inverted	10-8	Foreground	Color
			7-0	Character	ASCII

NOTE: Not a user callable procedure.

```

var REGS : registers;

begin [*] GET_SCREEN_AND_CURSOR_MASKS [*]
.....
Parameter|s| Return Values
AX GET_CURSOR_MASKS_AND_MICKY_CNTS Screen mask value [or scan line
start]
BX Cursor mask value [or scan line
stop]
CX Raw horizontal mickey count
DX Raw vertical mickey count
with REGS do
begin
AX := GET_CURSOR_MASKS_AND_MICKY_CNTS;
intr|MOUSE_SERVICE_NUMBER, REGS|;
SCREEN_MASK := AX;
CURSOR_MASK := BX;
end; [*] with...do [*]
end; [*] GET_SCREEN_AND_CURSOR_MASKS [*]

```

Masks			Screen Bits		
Screen Mask	Cursor Mask	Screen Bit	Bit Number	Description	Comments
0	0	0	15	Blink Ctrl	1=Blink
0	1	1	14-12	Background	Color
1	0	Unchanged	11	Intensity	1=High
1	1	Inverted	10-8	Foreground	Color
			7-0	Character	ASCII

NOTE: Not a user callable procedure.

```

var REGS : registers;

begin [*] SET_SCREEN_AND_CURSOR_MASKS [*]
.....
Parameter|s| Return Values
AX SET_TEXT_CURSOR
BX Cursor Type
0=Software Cursor
1=Hardware Cursor

```





Evolution of Solutions to Real-Time Problems

by  
Greg P. Semeraro

A Thesis Submitted  
in  
Partial Fulfillment of the  
Requirements for the Degree of  
MASTER OF SCIENCE  
in  
Computer Engineering  
April 1997

R · I · T

ROCHESTER INSTITUTE OF TECHNOLOGY

